# Assembler ROM Manual

Owner's Manual

HP-87





Assembler ROM and HP 82928A System Monitor Reference Manual

HP-87

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#### INTRODUCTION

#### 1.1 General Information

This manual outlines the commands, statements, instructions, and use of both the HP-87 Assembler ROM and the HP 82928A System Monitor. The manual assumes you have some knowledge of programming in assembly language. If you are not familiar with the HP-87 Personal Computer, you should read the owner's manual.

The HP-87 contains both read only memory (ROM) and read-write or random access memory (RAM). The RAM contains the user's BASIC language programs and data, and can also contain up to five binary (machine language) programs. The ROM contains the machine language program that recognizes and executes the statements provided by the BASIC language. Thus, the operating system ROM provides such statements as PRINT, DISP, and INPUT.

When external peripheral devices are added, their wider range of capabilities requires more extensive BASIC language statements to fully use these capabilities. Additional external ROMs enrich the BASIC language by increasing the number of statements and functions that can be recognized and executed. Similarly, a binary program also extends the BASIC language.

#### 1.2 The Assembler ROM

Using the Assembler ROM, you can write assembly language binary programs for residence and execution within the computer or for creation of a plug-in EPROM for the computer. A binary program can:

- Extend the BASIC language.
- Give increased execution speed.
- · Redefine the system.

The Assembler ROM permits you to enter and edit source code for binary programs on the computer's CRT screen. Automatic line numbering and cursor movement are active, and the source code can be stored on a mass storage device, listed, and edited. As source statements are entered, they are automatically checked for syntax errors and duplicate labels.

At assembly time the resulting object code (machine language) is stored on the mass storage device. The object code can also be loaded automatically or on command, and it is then ready to run.

To aid in programming, a disc is supplied with the Assembler ROM. This disc contains a global file of the system labels and their memory addresses for use during assembly. The disc also contains the sample programs from section 7 to help illustrate how binary programs are created and run.

The Assembler ROM gives you the ability to tailor statements for your own applications, to speed up program execution, and to perform sophisticated graphics. But with all the power and system accessibility provided by the Assembler ROM, it is also possible to defeat the computer's internal safeguards and even seriously damage the computer. For this reason, you should understand assembly language programming before attempting to use the Assembler ROM.

## 1.3 The HP 82928A System Monitor

The system monitor is an optional plug-in module that is designed for use only in conjunction with the Assembler ROM. The system monitor is not required, but it makes the debugging and modification of binary programs much easier.

With the system monitor module attached, you can set breakpoints that interrupt the execution of a program. After program execution has been interrupted, you can examine or change the contents of memory, execute one instruction at a time (single-step), or you can trace the operation of a machine language program, printing the status of the CPU after each instruction.

System monitor instructions are discussed in detail in section 5 and the use of these instructions is demonstrated in section 7.

## 1.4 Using HP-83/85 Binary Programs on the HP-87

The HP-87 uses the same CPU as the HP-83/85. The programs are entered, stored, listed, and run in the same manner. There are some differences on the HP-87 which include:

- BASIC programs are stored in reverse order (executing from the higher addresses and progressing to the lower addresses).
- The extended memory controller makes it possible to access more memory.
- Five binary programs can be resident in the computer at a time.
- · PTR1 is used as the BASIC program execution pointer at run time.
- · PTR2 is used as the output stack pointer at parse time.
- Entire programs are no longer allocated before execution begins.
- The BASIC program control block in the HP-87 is 40 bytes long.
- The operating stack is of fixed length in the system RAM.
- String values are passed on the operating stack as a two-byte length and a three-byte address.
- Inverse video, more display modes, eight bit CRT addresses, and access during horizontal retrace periods are a few of the changes affecting the CRT.

Because the differences are only highlighted in this section, you should refer to individual sections in this manual to become familiar with the HP-87 Assembler ROM before writing programs.

To modify an existing HP-83/85 binary program for use on the HP-87:

 Pick a binary program number and put it in the NAM statement. This should be a value between 200 and 377 (octal). Numbers from 0 to 177 are reserved for use by Hewlett-Packard.

Two different binary programs may have the same binary program number, but they cannot be loaded and used at the same time. Attempting to do so will cause a BAD BIN-LOAD error.

2. Modify any ABS statements. In the HP-83/85, all binary programs were loaded so that the absolute base address could be calculated by the Assembler ROM at assembly time based upon the length of the binary program. In the HP-87, this is not true. A change must be made in the ABS pseudo-opcode. You must have an absolute base address.

This only applies to binary programs that were written as absolute code. Most binary programs are relative and not affected by this change.

- Modify all parse routines to use PTR2 as the output pointer rather than R12-R13.
- 4. Modify all parse routines to push the binary program token out as:

TOK# BPGM# 371

rather than:

#### 371 GARBAGE-BYTE TOK#

- 5. Change all RUNTIME references to R10 into references to PTR1.
- Modify all code that uses string parameters that are passed on the R12 stack. These strings use three-byte addresses on the HP-87, rather than the two-byte addresses used by the HP-83/85.
- Check all references to any system routines to see if any changes have occurred to the input/output conditions of the routine. Make any necessary changes.
- 8. Change all system/global address definitions.
- Any routine that gets control through a RAM hook (such as CHIDLE, KYIDLE, IOTRFC) must calculate the base address of the binary program rather than loading it from BINTAB. Use the code:

LABEL LDM R20,R4 BIN SBM R20,=LABEL

This will leave R20-R21 with the absolute base address of the binary program. This change is necessary only in relative binary programs.

10. If the binary program uses its own error messages, ERRBP# (a RAM location in the system global addresses) must be set to the binary program number before calling ERROR or ERROR+.

#### 1.5 Assembler Commands, Statements, and Functions

The commands and the statements and functions provided by the Assembler ROM are added those which are already part of the instruction set. They are executed exactly as the rest of the instruction set, and have been created to help the programmer control and use the assembler.

Assembly language elements are used as the actual instructions in writing binary programs. The format and use of these elements are discussed in section 6, and complete list may be found in sections 6 and 8.

#### Assembler Commands

A command is nonprogrammable, and can be executed only from the keyboard. The assembler commands permit the user to transfer between assembler and BASIC system modes, to assemble, store and load binary program source code, and to find labels within the source code in memory.

## ALOAD file name Assembler Command

Legal only in assembler mode. Loads source code that was previously stored with the ASTORE command into computer memory from the file specified on the currently selected mass storage device. The file must be of the type known as extended \*\*\*\* or ASSM.

Note: The extended type of file, denoted by \*\*\*\* on the directory of a mass storage device, does not necessarily mean that the file contains source code. In fact, other HP firmware and software may generate extended type files.

ASSEMBLE file name [,numeric value]
Assembler Command

Legal only in assembler mode. Assembles source code currently in the computer memory and stores it in the file specified by file name on the currently selected mass storage device. The assembled source code is stored as either a binary program or, if the file has been declared a ROM or global file, as a series of strings in a data file.

If at assembly numeric value is evaluated as zero, the binary program currently in the computer memory is scratched, and the object code of the newly assembled binary program is loaded from the mass storage device into memory. Default numeric value is evaluated as zero.

If at assembly numeric value is other than zero, any binary program currently in memory remains inviolate, and the object code of the newly assembled binary program is stored only on the current mass storage device.

Note: If a program contains an error or if programs are linked at assembly, this command can destroy the source code; if the source code is to be saved on a mass storage device, it should be stored there before typing ASSEMBLE.

ASSEMBLER Assembler Command

Legal only when the computer is in normal system mode, this command scratches memory and puts the computer into assembler mode. In assembler mode, most normal BASIC statements will still operate, but only as calculator mode statements; they are not programmable. Source code for a binary program can then be typed in with line numbers, just as a BASIC program is typed in while in normal system mode (but with only one instruction per line). Unlike its operation in normal system mode, the computer is somewhat sensitive to character spacing while in assembler mode. Auto line numbering, screen editing, listing, etc., are all function. The [CONT], [STEP], and [INIT] keys are inoperative in assembler mode. Displays READY when executed.

ASTORE file name Assembler Command

Legal only in assembler mode. Stores the source code currently in the computer memory into the specified file on the currently selected mass storage device. File is of the type known as extended, shown in the directory as extended (\*\*\*\*) or ASSM.

BASIC Assembler Command

Legal only when in assembler mode, this command scratches memory and puts the computer back into BASIC mode. Display READY when executed.

FLABEL label Assembler Command

Legal only in assembler mode. This command searches through the source code in memory for the label specified. For each occurrence of the label the line is listed. After an FLABEL command has been executed, pressing the [LIST] key causes the source code to be listed, beginning with the last line where the label occurs.

FREFS string Assembler Command

Legal only in BASIC or assembler mode. Searches through the source code in memory for all occurrences of the specified string. After an FREFS command has been executed, pressing the [LIST] key causes the source code to be listed, beginning with the first line where the string occurred. Pressing any key will cause the FREFS command to halt prematurely.

Assembler Statements and Functions

Statements and functions are programmable BASIC language elements. The statements and functions provided by the Assembler ROM are simply additions to the BASIC language of the computer. As with all BASIC statements and functions, they may be used either in calculator mode or as part of a BASIC program when in BASIC mode. When the computer is in assembler mode, all BASIC statements and functions may be executed only from the keyboard.

DEC

Assembler Provided BASIC Function

Returns the decimal equivalent of the specified octal value.

MEM address [:ROM#]],# of bytes][=#,#,..]
Assembler Provided BASIC Function

Dumps the contents of computer RAM or ROM memory to the current CRT IS device beginning with the octal address given. Continues dumping for the specified octal [,# of bytes]. At power-on, default # of bytes is 100 octal; otherwise, default is the last # of bytes specified.

The [:ROM #], if included, is an octal value that selects the plug-in ROM from which memory is dumped. At power-on, default value for ROM # is Ø; otherwise, default is the last ROM # specified.

If =#,# is included in the statement, memory is not dumped, but instead the contents of memory locations beginning at the address given are changed to the octal values specified after the = sign. The memory locations must be in RAM. The contents of one succeeding memory location are changed for each value specified after the = sign. The # of bytes, if included in the statement, is disregarded in this case. Pressing any key will cause the memory dump to halt.

MEMD address [:ROM#][,# of bytes][=#,#,...]
Assembler Provided BASIC Statement

Same as MEM except reads the contents of three bytes of memory beginning with the address given and uses those contents as the address.

OCT decimal numeric value Assembler Provided BASIC Statement

Returns the equivalent of the specified decimal value.

REL octal address Assembler Provided BASIC Statement

Returns the absolute address of a relative address. Takes the relative octal address and adds to it the address (called BINTAB) of the beginning of the last binary program that was accessed to yield the octal absolute addres. May be used alone or with the MEM command. May also be used with the command BKP if HP 82928A System Monitor is installed.

SCRATCHBIN Assembler Provided BASIC Statement

Scratches all current binary programs from computer memory, without affecting anything else.

#### CPU STRUCTURE AND OPERATION

### 2.1 CPU Register Bank

The central processing unit (CPU) consists of 64 eight-bit registers, an address register pointer (ARP), a data register pointer (DRP), an arithmetic-logic unit (ALU), a shifter, and a set of status indicators.

The 64 eight-bit registers are grouped into two sections. The first 40 (octal) registers have two-byte boundaries and are used principally for addresses. Many of these bytes are reserved by the CPU for use as special purpose registers, and direct access to these should be avoided. The next 40 (octal) registers are separated by eight-byte boundaries. Floating-point numbers, 64 bits long, are stored here. The programmer must be aware of what is destroyed when the system uses these registers. The effects of system routines on register contents are found in section 8.

Any register in the CPU may be used as an accumulator when performing an operation. To distinguish between the registers, the CPU uses the DRP to designate the accumulator and the ARP to designate the operand. The DRP directs the results of arithmetic operations to the register it points to, and the ARP supplies the second operand when it is needed. Both the ARP and the DRP can be used to address any of the bytes in the CPU register bank. The CPU register addressed by the ARP is called the address register, or AR. The register addressed by the DRP is called the data register, or DR.

Hardware-Dedicated Registers

Registers	Description		
0,1	Register Bank Pointer: RØ points to the remainder of the CPU register bank. R1 is only accessable through RØ.		
2,3	Index Scratch: R2-R3 are used for address calculation for indexed addressing.		
4,5	Program Counter (PC): R4-R5 hold the absolute address of the next instruction location.		
6,7	Return Stack Pointer: R6-R7 contain the pointer for the subroutine return stack. When a "JSB=" subroutine jump is execute the CPU pushes the PC (R4-R5) on the stawhen the RTN is executed, the CPU pops to bytes from this stack and places them in R4-R5 (program counter).		

## Software-Dedicated Registers and EMC Pointers

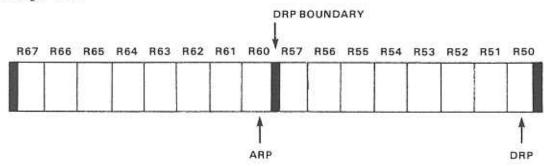
Registers/Pointers	Description		
PTR1	At run time, contains the program counter (PCR), a pointer for executing BASIC programs.		
PTR2	At parse time, used to point to the parse output stack.		
16,11	Not software dedicated at run time. When parsing, R10-R11 point to the next character of the input ASCII stream.		
12,13	Operation Stack: Parameters and results are passed on the stack pointed to by this register pair. Contains expressions when the BASIC program is decompiling.		
14	When parsing or decompiling, R14 contains the current token being processed.		

Software-Dedicated	Registers	and	EMC	Pointers
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Registers/Pointers	Description
16	Current Status (CSTAT): R16 contains the code that indicates the current mode of operation. The table of CSTAT codes is found in paragraph 3.4.
17	External Communication Status (XCOM): When an external interrupt takes place the status is stored in R17. The table of XCOM status codes is found in section 3.

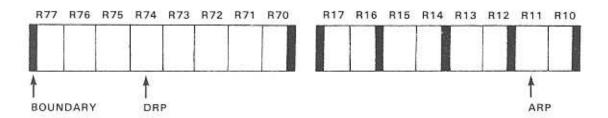
Multi-byte operations can be performed with the help of the register boundaries. The number of consecutive registers that will be used in the operation is determined by the distance between the DRP and the next boundary.

Example: In a multi-byte addition a 64-bit quantity contained in registers 50 through 57 will be added to a 64-bit quantity in registers 60 through 67.



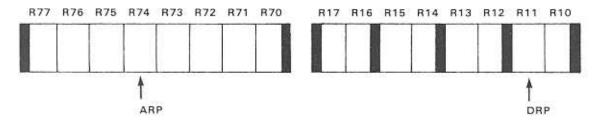
The operation begins with the registers pointed to by the DRP and the ARP, processing the registers within the boundary. The result is stored as a multi-byte quantity in the registers pointed to by the DRP.

Example: A multi-byte load with the DRP set to R74 and the ARP set to R11 will load the the four registers R74-R77 with the contents of R11-R14.



The boundary is determined by the DRP and is ignored by the ARP. In the previous example, the load terminates when the DRP reaches the next boundary.

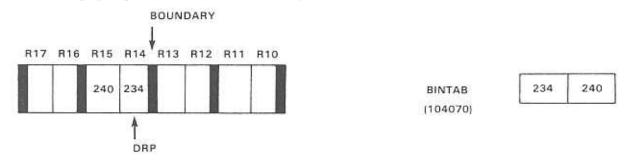
Example: The multi-byte store recognizes the boundaries in exactly the same way as the multi-byte load. Attempting to store with the DRP set to R11 and ARP set to R74 would result in the loss of several bytes due to the boundary after R11.



The boundary after R11 stops the multi-byte operation. Only one register is transferred to its destination, that is, R74. The DRP always determines how many bytes will be involved in a multi-byte operation.

There are also two-operand operations where the DRP points to one operand, and the second is located in computer memory. The number of bytes used in the operation is dependent upon the boundary after the DRP. That number of bytes of memory will be used starting at the location described by the label or pointer accessing computer memory.

Example: This load will be done with the address BINTAB, which is a label pointing to an address which contains the address of the start of the binary program. The DRP will point to R14.



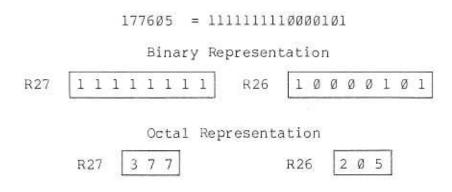
Because the boundary is two bytes from the DRP, two bytes are accessed from memory.

#### 2.2 Number Representation

The CPU can operate on numbers as octal and binary-coded decimal (BCD) quantities. All registers and register addresses are represented as octal numbers, and all floating-point numbers are represented in BCD notation, that is, each decimal digit is stored as a four-bit binary number, with two digits per register. Since the CPU cannot tell one representation from another, it is important to keep track of the way numbers are stored when doing arithmetic operations.

An address is always an octal value that occupies 16 bits or, for the extended memory pointers, 24 bits. The highest-numbered byte contains the first, or most significant, part of the address, and the lowest-numbered byte contains the last, or least significant, part of the address.

Example: The octal address 177605 is stored in two registers, R13 and R12. An address is always an octal value that occupies 16 bits and is contained in two registeres.



The ARP and the DRP will always point to the least significant byte of a multi-byte operation.

With BCD numbers, decimal one is represented with four bits,  $\emptyset$   $\emptyset$   $\emptyset$  1, decimal two is  $\emptyset$   $\emptyset$  1  $\emptyset$ , on up to decimal nine, which is 1  $\emptyset$   $\emptyset$  1. When the decimal number has more than one digit, each digit is represented by four bits.

Example: The decimal number 3738 is represented by 16 bits.

Octal Representation

Each byte can contain two four-bit BCD digits. Each register can represent numbers in the range 00 to 99.

The ten's complement is used to simulate subtraction exactly like the two's complement is used in binary arithmetic. The ten's complement is formed by subtracting each binary-coded digit from nine (nine's complement arithmetic),  $1\ \emptyset\ \emptyset\ 1$ , then putting the digits back together to form the number again and finally incrementing the entire quantity by one (one's complement arithmetic).

The negative of a number in BCD representation, for subtraction purposes or in special cases to show the sign of an exponent, is found by taking the ten's complement.

Example: To find the negative of 19, each digit, 1 and 9, is subtracted from 9, or, another way of looking at it, 19 is subtracted from 99.

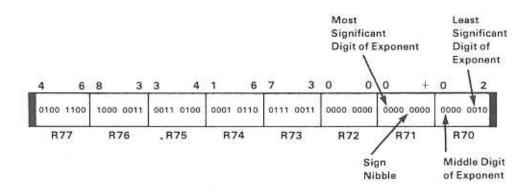
Add one to the combined result:

In effect, when 81 is added to 19 the result is 00 in BCD notation.

Numeric quantities may be represented as real floating-point, short, and integer formats. The real and short forms are expressed as BCD digits, and the integer form is a five-digit number with a sign digit at the end of the quantity. The system represents all numeric quantities in BCD notation.

Real numbers have a mantissa of 12 digits, and exponent and sign information, all stored in eight bytes. The mantissa fills the 12 most significant nibbles of the number, the sign takes one nibble, and the exponent is contained in the last three nibbles. The most significant digit of the number is stored in the most significant byte, and the decimal point is assumed to be immediately after the most significant digit. The sign of the number follows the least significant digit of the mantissa, and the exponent, expressed in ten's complement notation, is found in the three least significant nibbles of the quantity.

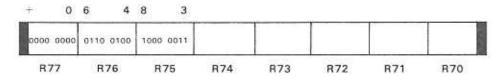
Example: The real number 468.3341673 (in scientific notation 4.683341673 x 10^2), would be represented in BCD as:



The radix is assumed to be in R77 between the four and the six.

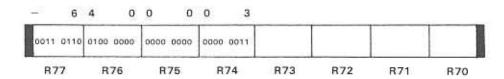
Integers are stored in three bytes, with five digits and a sign. The most significant digit of the integer is stored in the least significant byte. Real and short number representations are not right justified like integer representation.

Example: The integer 6483 would be represented in BCD notation from the least significant digit, 3, to the most significant digit, 6, with the sign, positive or 0 0 0 0, in the most significant four bits of the quantity.



Short numbers have a mantissa of only five digits and an exponent of two digits. Both the mantissa and the exponent have sign bits, found in the most significant digit. The representation of the mantissa begins immediately following the sign bits with the most significant digit of the mantissa found in the second digit of the most significant byte. The assumed decimal point is directly after the first digit, then the rest of the mantissa is represented. The two least significant digits hold the exponent, which is not in complement form because the sign of the exponent is in the most significant digit.

Example: The short number -.0064 need not be represented as a 12-digit real number. In BCD short form it is represented as:



The radix is assumed to be between R77 and R76.

#### 2.3 Status Indicators

The CPU contains eight flags and a four-bit register for program status. The flags signal the present condition of the data, while the four-bit register serves as an extended register for counting and data manipulation.

Status can affect or be affected by CPU instructions. The instruction set has data movement instructions of both the arithmetic and nonarithmetic types. These instructions include:

- Arithmetic: Add, substract, compare, increment, decrement, and complement.
- Nonarithmetic: Load, store, "and", "or", "exclusive or", shift, clear, and test.

The CPU contains the following one-bit status flags and four-bit extend register:

DCM Decimal Mode Flag: This flag determines whether the system is using binary numbers or BCD numbers in arithmetic operations. In BCD mode, each decimal digit is converted to BCD, and all arithmetic operations are done with the resulting four-bit digits. This is the way floating-point real numbers and integers from BASIC programs are represented.

The system uses this flag to determine the correct mode, so the user must make sure it is set properly for arithmetic operations. All shifts and all arithmetic operations are affected by the DCM flag.

Two instructions affect the status of the DCM flag: BCD sets it to 1, and BIN clears it.

- Extend Register: In BCD mode, this four-bit register will accept the displaced digit resulting from a shift. Once in the register, a BCD digit may be incremented, decremented, or cleared, and, if needed, the digit may be returned to the register it came from using the extended shift instructions.
- CY Carry Flag: In binary mode, this flag will indicate the result of a bit shift. A bit may be shifted into the CY flag, tested, and then shifted back into a register, using the extended shifts. It functions similar to the extend register in BCD mode.

During all arithmetic operations, the CY flag will be set with the carry out of the most significant part of the operation. In addition between two numbers where the result is too large for the register to hold, or in subtraction where the result is positive, the CY flag is set to 1. The CY flag may be thought of as the "borrow" if needed for subtraction.

When two quantities are added, the CY flag is set with the carry, if any, resulting from the addition of the most significant bits.

## Examples:

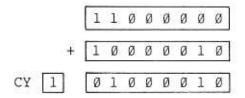
If two positive numbers, both with a most significant bit of 0, are added, then the carry will always be 0.

If a positive number is added to a negative number, in reality a subtraction, then two possibilities could occur:

 The result could be negative, in which case no carry would be made.

The result could be positive, causing a carry out.

If two negative numbers are added then the CY flag is set to 1.



The carry flag is set by comparisons in the same manner as additions.

An increment sets the CY flag if the data register is all 1's.

OV Overflow: The overflow status is determined by taking the "exclusive or" of the CY flag and the most significant bit of the data register. It is set to I when the addition of two positive numbers yields a negative result, when the addition of two negative numbers yields a positive result, and when the result of a left shift changes the sign of the data register.

OD Least Significant Bit: After any data movement instruction, the least significant bit is shown as the OD flag. If the OD flag is set to 1, then the number is odd. If the OD flag is 0, the number is even. The right-most bit in the data register is always the least significant bit.

## Section 2: CPU Structure and Operation

NG Most Significant Bit: This flag displays the most significant bit in the data register. If this flag is set to 1 then the quantity is negative, and if the NG flag is clear then the quantity is positive.

ZR Zero: If the data register is Ø or if a comparison is made between two equal numbers then this flag is set to 1.

LDZ Left Digit Zero: This flag is set if the left-most four bits are  $\emptyset$   $\emptyset$   $\emptyset$ . In BCD mode this would indicate the most significant digit.

RDZ Right Digit Zero: If the least significant four bits are 0 0 0 0 then this flag is set to 1. In BCD mode this would indicate the least significant digit.

Example: Status information is based on the entire multi-byte quantity that is being processed. All multi-byte operations, except right shift, start execution with the least significant byte. The right shift starts with the most significant byte. All status flags, except OD, RDZ, and DCM, are updated after each byte of execution and will be correct as the register boundary is met. The OD and RDZ flags are set for the first byte and never changed. The E, CY, and OVF flags are only affected by arithmetic operations.

## Section 2: CPU Structure and Operation

After the multi-byte addition of the two system addresses, OFFSET (000100) and the label VARIABLE (000365), the status indicators will be set as follows:

OFFSET 1 0 0 0 1 0 0 0 00111000 11110101 + VARIABLE ØØØØØØØØ RESULT 1 0 0 0 1 0 0 1 00101101 DCM E CY OV NG OD LDZ RDZ ZR Ø Ø 1 1 Ø 0000 Ø Ø Ø

#### OPERATING SYSTEM

## 3.1 Introduction

This section explains how system memory is allocated, how extended memory is accessed, and how a statement is parsed and becomes part of a BASIC program. It also explains the sequence of operations that occurs when a BASIC program is run.

BASIC programs are executed by an interpreter. However, the code that is interpreted is vastly different from the BASIC statements as they were originally entered. As the statements are entered, they are parsed and compiled into a form of RPN (Reverse Polish Notation), which can be interpreted more efficiently. The BASIC reserved words are converted to single-byte tokens (refer to Execution by Tokens). This makes the internal form of the code somewhat more compact than the original form, and also makes interpretation easier and faster.

Also during the process of parsing and compiling, variables are placed in a variable storage area, with only their addresses and names remaining in the area containing the tokens. The BASIC program is held in memory as a series of tokens and addresses of variables and associated data bytes. To execute the program, the computer processes these token and variable addresses in order. As each token is processed, it causes the machine to access a table of routine addresses and execute a specific routine corresponding to the token. If the token indicates a variable, the machine uses the next three bytes as the variable address.

## 3.2 System Memory

Several distinctly different regions comprise the system memory. They are (all numbers are octal unless indicated otherwise):

- Six system ROMs, each containing 8192 decimal bytes. A subset of the ROM area is the address range from 60000 to 77777. This range is shared by system ROMs 0, 1, 320, and all of the external plug-in ROMs. Each of the ROMs in this area can be selected or deselected for talking on the bus, but only one of them can be selected at a time. Each of these ROMs has a bank-select address which is its ROM number, ranging from 0 to 376. To select a particular ROM you store the desired ROM number to an I/O address called RSELEC. The chosen ROM will be selected and all other bank-selectable ROMs will be deselected.
- · RAM, 32768 (decimal) bytes in the basic machine.
- Memory addressable directly by the CPU (addresses 0 to 177777).
- Memory addressable through the extended memory controller (EMC), 32K-544K.
- The block of 400 addresses (177400 to 177777) which act as I/O addresses, when accessed directly by the CPU. The same addresses accessed through the extended memory controller will act as RAM memory, not as I/O addresses.

# Section 3: Operating System

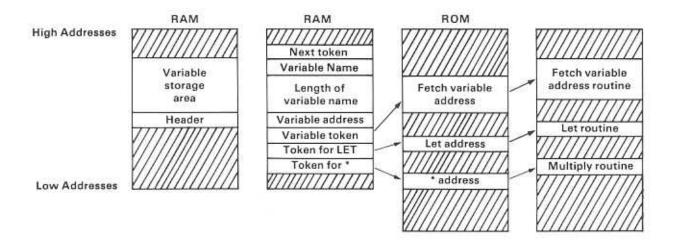
DECIMAL ADDRESS	OCTAL ADDRESS	8				
0	0	SYSTEM ROM				
8K	20000	SYSTEM ROM				
16K	40000	SYSTEM ROM				
24K	60000	ROM 0 SYSTEM	ROM 1 GRAPHICS	ROM 320 MASS	****	ROM x PLUG-IN
32K	100000	ROM	ROM	STORAGE		ROM
52.11		RAM (If accessed either directly by the CPU or through the EMC)				(OPTIONAL)
64K-256	177400	RAM	I/O ADDRES	and the contract of the contra		
64K	200000	(If accessed through the EMC)	(If acce directly the CPL	by		
1 M	4000000					

## Computer Operation

The basic machine is controlled by system routines that are permanently resident at fixed addresses in memory. The addresses and names of many of these system routines may be found in the global file in section 8.

In addition to the system routines, control can also pass to one of the plug-in bank-selectable ROMs, or to a binary program in memory. At certain times in the operation of the system, the resident binary programs and ROMs are polled by the main system. In addition, there are a number of entry points (hooks) that allow operation to be intercepted and modified by a binary program or ROM. These hooks are normally idle, but they can be used to take over the system at certain key times.

Execution by Tokens (Run Time)



Tokens are used to represent the keyword, such as LET, FOR, BEEP, etc., that make up each BASIC statement. Each token is a one-byte quantity that the machines uses to find the addresses of routines associated with that token. Each token must have an associated entry in a table of routines for execution at run time, another entry in an ASCII keyword table, and a third entry in a table of parse routines. A list of system tokens may be found in section 8.

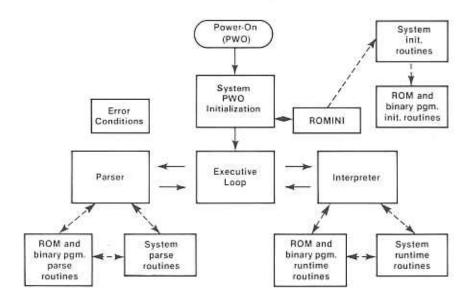
The computer is a token-driven machine. A program is held in memory as a series of tokens and variable addresses which the machine processes.

For example, at run time as the system executes a program, it processes a token by fetching the address of an associated run time routine from a table of addresses. The run time table may exist in a binary program and/or an external ROM as well as in the main system. The system jumps to the specified address to execute the routine, then fetches the next token and searches for its run time routine in the tables, etc.

Some tokens indicate to the system that the three bytes following the token contain a variable address. In this case, the system attempts to find the variable in the variable storage area and, if not found, creates a place for it. Other tokens indicate that the bytes following the token are constants to be pushed onto the R12 stack.

Two tokens, 370 (octal) and 371 (octal), are used to expand the token tables. Token 370 indicates to the system that the next byte is the number of a ROM, and that the byte after the ROM number is the token within the ROMs table that is to be executed. Token 371 directs the system to a binary program in the same way.

## 3.3 Overall System Flow



System flow is shown by the chart above. In general, loading and running a program, or executing a calculator mode statement, will require execution within the following areas:

Power-on Initialization: When the computer initially powers-up, it performs a sequence of operations: performs a self-test, accesses and resets any interface modules, reserves memory for later use, allows any ROMs to reserve memory, and returns to the system.

Executive Loop: External stimulus (such as a keyboard interrupt) and changes within the computer (such as an error condition) will cause the executive loop to call the appropriate routines to take control at the right time.

Parser: Parsing occurs when [END LINE] is pressed after a program line or calculator mode statement has been typed. Parsing is the changing of ASCII code into tokens. The parser first searches the ASCII tables in any resident binary programs for a keyword match, then the ASCII tables in any external ROMs, and finally the system tables. This makes it possible to redefine system keywords.

Interpreter: The interpreter actually runs a program or executes a calculator mode statement by fetching tokens and calling the run time routines to execute them.

In addition, there are two other areas which may be called:

Initialization: At many times, including power-on, RESET, SCRATCH, etc., the system calls routines for initialization. Initialization routines are called through the ROMINI routine; the system polls system initialization routines first, ROM routines second, and the routines in the resident binary programs last. ROMFL is a RAM location that initialization routines called by ROMINI can look at to see why they were called.

Initialization routines are called before, during, or after a condition occurs, depending upon the following conditions:

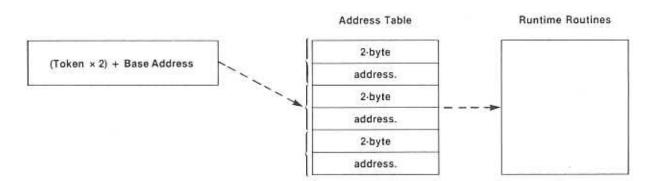
ROMFL	Meaning	Initialization Routines Called
Ø	Power-On	After system initialization.
1	RESET	After system reset.
Ø 1 2 3 4	SCRATCH	Before scratch.
3	LOADBIN	After loadbin.
4	RUN	Before execution begins.
	INIT	After allocation done.
5	LOAD	Before load.
6	STOP, PAUSE	During.
5 6 7	CHAIN	After.
10	Allocate token class>56	During.
	Deallocate token class>56	During.
11 12	Decompile token classs>56	During.
13	Program halt on error	During.

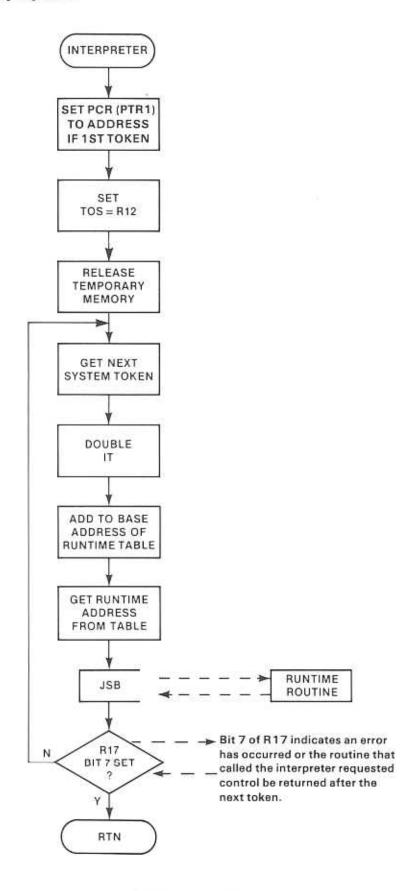
When errors occur, the system generates the proper warning or error message.

## Interpreter Loop

The interpreter loop fetches the next token, processes it, and passes control to the respective run time code. When the run time code has been executed, control returns and the interpreter continues with another token.

A token is an ordinal into a table of addresses. The address table is made up of two-byte addresses. To find the actual address, the token is doubled, then added to the base address. This changes the ordinal into an offset pointing to the current address.





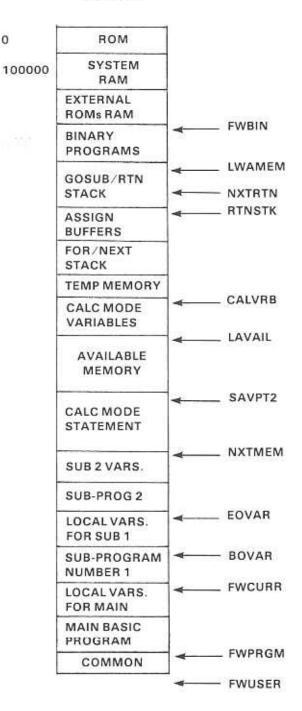
## 3.4 Allocation and Deallocation

Allocation is the process of reserving and assigning memory for program variables. The three modes of allocation are:

- If the program line is a dimension statement, the entire line is allocated before execution continues.
- If the current token being allocated is the start of a user defined function (DEF FN), then allocation will continue for the duration of the definition (until the FN END) before execution will resume.
- All other tokens are allocated one at a time as they're encountered.

The class of a token determines if the token needs to be allocated. The following diagram shows how memory looks in a program that has been allocated:

#### RUNTIME



The GOSUB/RTN stack is for the BASIC program not for assembly language, NXTRTN points to the next RETURN address on the stack.

Each ASSIGN buffer takes 284 bytes.

TEMP memory is released by the system at the end of each line of a program and when an "@" token (statement concatenation) is encountered.

Both the RUN and the CONT commands set LAVAIL equal to CALVRB, so during the running of a BASIC program, they will always be equal, and there will be no CALC mode variable.

If a CALC mode statement has been entered and is executing, it will begin at NXTMEM — 1 and end at SAVPT2. Otherwise, SAVPT2 will be equal to NXTMEM.

NXTMEM points to the last byte of the BASIC programs (or subprograms).

EOVAR points to the last byte of the variable space of the current active BASIC program and BOVAR points to the [first byte] + 1.

FWCURR points to the first byte of the currently active program.

FWPRGM points to the first byte of the MAIN BASIC program.

FWUSER points to one higher than the highest address in memory. Unless you have created a token whose class is greater than 56, allocation is handled by the computer. Because a program is allocated in segments, a memory overflow could occur after you are well into your program.

The system executes the BASIC program starting at the highest address. The line of BASIC code

### 10 A=B

would be parsed into this stream of bytes:

016	END OF LINE TOKEN
010	STORE SIMPLE VARIABLE TOKEN
102	B, THE VARIABLE NAME
001	LEN OF VARIABLE NAME
0001	
000	3 BYTES RESERVED FOR ALLOCATION TIME ADDRESS
000)	
001	FETCH SIMPLE NUMERIC VARIABLE TOKEN
101	A, THE VARIABLE NAME
001	LEN OF VARIABLE NAME
000)	
000}	3 BYTES RESERVED FOR ALLOCATION TIME ADDRESS
0001	
Ø21	STORE SIMPLE NUMERIC VARIABLE TOKEN
Ø16	LEN OF LINE (16 OCTAL BYTES FOLLOW)
020)	
000	BCD LINE NUMBER 10
000)	

Token number 21 (and tokens such as 1, 2, 3, 22, and 23) is immediately followed by three bytes which are used to contain the relative address from the first byte of the currently active program (FWCURR). Since the variable storage area for BASIC programs is at a lower address than the program, this relative address will always be negative (that is, the most significant bit of the address will always be set). Therefore, if the most significant bit is Ø then the system knows that the current token has not yet been allocated. In this case, the allocator would be called, which would search through the variable storage area for the current variable. If found, the allocator would calculate the correct relative address and place it where the three Ø's are following the 21 token. If not found, the allocator would create a storage area for it in the space allocated for variable storage (EOVAR), then calculate and store the relative address after the number 21 token. Execution would then continue.

Since the variable name could be long or short, the length of the name and the ASCII characters for the name immediately follow the address. The length of the name and the ASCII characters will be skipped at execution time, similar to the way a comment is skipped.

Line numbers are handled in a similar manner. The BASIC code

10 GOSUB 100

would be parsed like this:

016	END OF LINE TOKEN
0001	
001}	BCD LINE NUMBER OF DESTINATION
000)	
133	GOSUB LINE NUMBER TOKEN
005	LEN OF LINE
0201	
000}	BCD LINE NUMBER 10
000)	

Since line numbers are only five digits long, the most significant bit will be Ø if the line is not allocated. If the line is allocated the address will always be negative and the most significant bit will be set. All line numbers are converted into addresses relative to FWCURR (the first byte of the currently active program) at allocation time.

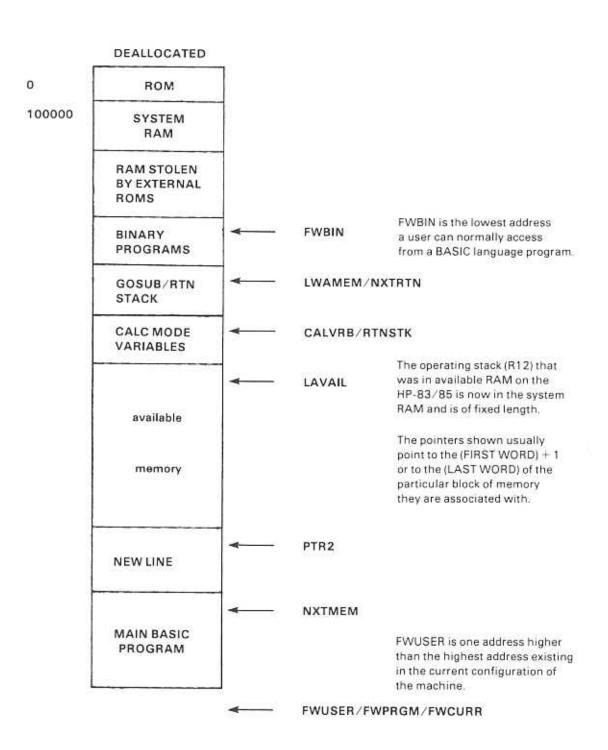
Line labels are handled in a way similar to variable names. The BASIC code

10 GOSUB [linelabel]

## is parsed as:

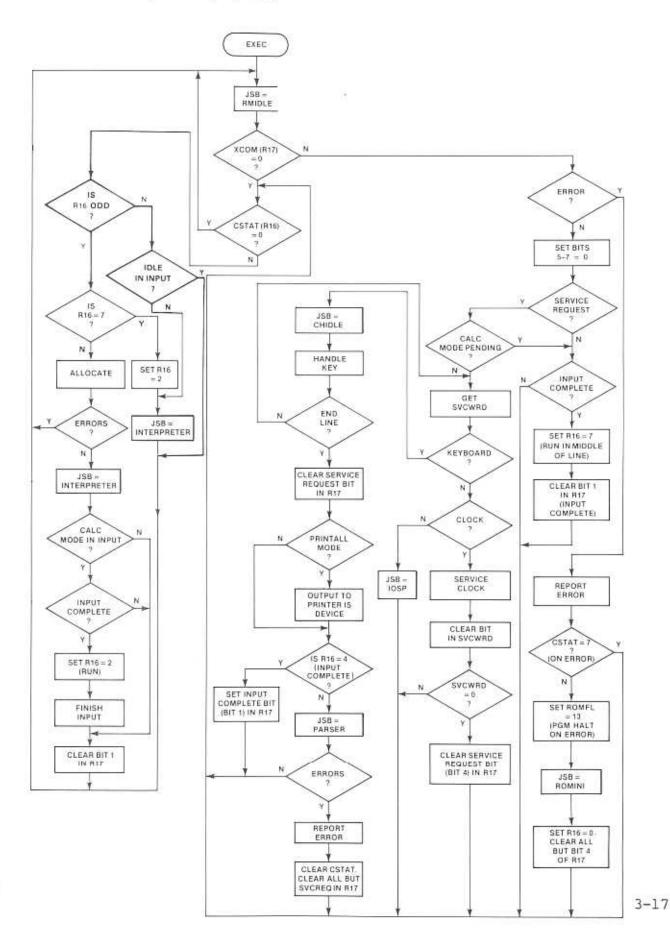
016	END OF LINE TOKEN
154	1
145	e
142	Ь
141	a
154	1
145	e
156	n
151	n i
114	L
011	LEN OF LINE LABEL
0001	
000	3 BYTES RESERVED FOR ALLOCATION ADDRESS
000)	
270	GOSUB line label TOKEN
Ø17	LEN OF LINE
0201	
000	BCD LINE NUMBER 10
000)	

A program is deallocated while you are typing in lines or while it is being edited. When a BASIC statement is typed and [END LINE] is pressed, the computer deallocates the program if it has not already been done. Program variables are held as names rather than addresses. This diagram shows memory when a program is deallocated:



## 3.5 Executive Loop

After power-on initialization, the executive loop portion of the system takes control. The executive loop examines CPU registers R16 and R17 for changes in the status of the computer, listens for external communications, and takes the appropriate actions based upon the information received. The current status information (CSTAT) is kept in register R16 and the external communication flags (XCOM) are kept in register R17. As long as registers R16 and R17 are both zero, the system is idle. The executive loop flowchart is shown on the next page.



## 3.6 Interrupts

When there is a change in status and the system is no longer idle, CSTAT (R16) indicates the computer mode of operation, according to the value stored there.

CSTAT

Value	Current Status		
Ø			
1	Calculator mode execution.		
2	Program is running.		
3	Not used.		
4	Idle during input statement.		
5	Calculating during input statement		
6	Not used.		
7	RUN in middle of a line.		
8-255	Not used.		

If execution halts, the system needs to know what caused it to halt and how to handle it. Each of the eight bits in XCOM (R17) have a different meaning associated with it. The service request bit is the only bit directly affecting interrupts.

XCOM

Bit Set	Execution Halt
7 6 5 4 3 2 1 Ø	
X	End of calculator mode.
X	Input complete.
X	Step mode.
X	Trace mode.
X	Service request (any interrupt).
X	Immediate set.
X	Error set.
X	Break ("or" of bits 5 and 6).

One of the controlling devices on the internal communications bus will generate an interrupt to begin execution. An interrupt will set bit 4 in R17 (XCOM) and a bit in a memory location which is used to keep track of the cause of an interrupt (SVCWRD). The executive loop knows that the interrupt has occurred (from XCOM) and which device caused the interrupt (from SVCWRD).

SVCWRD Bit Set	Type of Interrupt
7 6 5 4 3 2 1 0	
x x x	Keyboard interrupt. I/O interrupt. Timer 1 interrupt. Timer 2 interrupt.
X X X	Timer 3 interrupt. Special interrupt. Not used. Not used.

Whenever an interrupt occurs, the CPU expects the interrupting device to send a pointer to an interrupt handling routine in a table of addresses. This pointer is a one-byte quantity and the two bytes that it points to in memory indicate the starting address of the service routine. If multiple interrupts occur then the first interrupt is handled and the rest are disabled.

The service routine pointers are located at addresses  $\emptyset$  thru 25 in memory.

Table of System Interrupt Pointers

ADDRESS	CODE	FUNCTION
000000	DEF STARTX	Power-on vector.
000002	DEF SPARØ	Spare hook 0.
000004	DEF KEYSRV	Keyboard.
000006	DEF SPARØ	Spare hook 0.
000010	DEF CLKSRØ	Clock Ø.
000012	DEF CLKSR1	Clock 1.
000014	DEF CLKSR2	Clock 2.
000016	DEF CLKSR3	Clock 3.
000020	DEF IRQ20	I/O modules.
000022	DEF SPAR1	Spare hook 1.
000024	DEF SPAR1	Spare hook 1.

An interrupt may be caused by the keyboard, a timer, an I/O module, or a special device. Keyboard interrupts are handled using KEYSRV and the character editor (CHEDIT). If the clock causes an interrupt, an ON TIMER routine is called. An interrupt from an I/O module is handled by the IRQ2Ø and IOSP hooks, and special interrupts must be handled by the spare interrupt routines SPARØ and SPARI from other hardware.

Programmer created interrupt routines may be handled by taking control of certain memory locations accessed by the executive loop or by taking control of the interrupt service hooks SPARØ, SPARI, KYIDLE, or IRQ2Ø. The interrupt service hooks are accessed prior to the executive loop. Therefore, these locations may bypass the executive loop. Jumps to these locations (hooks), cause the instructions located there to be executed. Initially, only a RTN instruction is stored at each of these locations, so control immediately passes back to the executive loop.

When an I/O device interrupts the system, a jump is made to IRQ20 before control passes to the executive loop. This gives the I/O interrupt routine the chance to bypass the operation of the executive loop, taking control more efficiently.

The executive loop always performs these functions: tests CSTAT, tests XCOM, and jumps to RMIDLE. If an interrupt has occurred from the keyboard, a jump is made to CHIDLE. When an I/O interrupt occurs a jump is made to IOSP, provided that the proper bits in XCOM and SVCWRD are set.

When an interrupt occurs during the execution of a program, the CPU finishes the current instruction, saves the program counter (R4-R5) on the R6-R7 stack, and acknowledges the interrupt. The device puts a pointer to the address of the service routine on the bus, and the CPU loads the service routine address into the program counter (R4-R5). This is effectively a subroutine jump to the service routine, because the return address has been saved. The status of the CPU and the contents of any registers that will be used in the service routine must be saved and restored from within the routine. This is important because an interrupt could occur between the execution of an instruction which sets the status indicators and an instruction that depends on that status.

### 3.7 Hooks

A binary program or a ROM can gain control of the system using RAM hooks. Some are accessed directly by the executive loop and some by routines that branch from the executive loop. The four types of hooks are:

- Language hooks: Allow you to create new BASIC keywords or redefine existing ones.
- General hooks: Allow you to take over various parts of the operating system by storing subroutine jumps to a binary program or ROM routine at specified RAM locations.
- Initialization routines: Called by the system, external ROMs, and binary programs at initialization time. An initialization routine can steal RAM, change flag status, or gain control of the operating system.
- Error message table: Allows a binary program or ROM to flag specialized error conditions with custom error messages.

### Language Hooks

With language hooks the binary program or ROM can define new keywords, functions, or auxiliary tokens. Because the system first polls the resident binary program and then all external ROMs, a binary or ROM program can take over or supersede the system tables.

#### General Hooks

To provide for each general hook, the system at certain times executes a subroutine jump to a specific RAM location. During normal operation each of these RAM locations contains a RTN or is otherwise idle. By placing a jump to a binary program or ROM at the hook location, the program or ROM gains access to the operating system. It is the responsibility of the external program writer to determine how to use the hook and how to avoid conflict with other usages of the hook. No support is supplied by the system.

Because support is not supplied by the system before calling any of the RAM hooks, any binary program base address might be in BINTAB when the system calls a hook. You must ensure that the correct base address is loaded into BINTAB before a hook is taken.

The following code stores a copy of the binary base address for future use:

\*\*Initialization Routine\*\* LDBD R34,=ROMFL See why INIT routine was called. BIN Binary mode necessary for CMB and ADM instructions. CMB R34,=3 Is a binary program being loaded? If no, return. JNZ INITRTN LDMD R34,=BINTAB If yes, save the binary base address. STMD R34, X34, OURBAS Store it in the program. ADM R34,=HOOK Make hook routine address absolute. Make a copy of the address to store STM R34, R45 in HOOK. LDB R47,=236 Load the opcode for return instruction. LDB R44,=316 Load the opcode for a JSB instruction. Store R44-R47 into CHIDLE. STMD R44,=CHIDLE \*\*Hook Routine\*\* INIRTN RTN Done. HOOK BIN Entry to hook routine. DRP R34 Set the DRP to R34. BYT 251 Do a LDM R34, = instruction. \*\*Store Base Address Here\*\* OURBAS BSZ 2 Base address is stored here. STMD R34,=BINTAB Load the base address into BINTAB.

Unless otherwise noted, each general hook is seven bytes long. Flowcharts are provided for selected hooks in section 8 of this manual. General hooks are supplied at the following points:

RAM Name	Location	Function
CHIDLE	103670	Character editor intercept.
DCIDLE	104035	System decompiler hook called at entry time. If you take this hook and don't want to let the system have a chance at decompiling, then you need to discard a couple of return addresses.
DGHOOK	104044	If the PLOTTER IS select code is one or two and a DIGITIZE command is executed, this hook will be called so software that has been loaded can digitize off of the CRT.
IMERR	103724	Used to expand the IMAGE statement. This hook is called when there is something in an IMAGE statement that the system doesn't recognize.

RAM Name	Location	Function			
IOSP	103652	I/O service pointer. Used by I/O and mass storage ROMs.			
IOTRFC	103643	General output hook. If the select code of the CRT or PRINTER IS device is not 1 or 2, the DISP or PRINT will go to IOTRFC.			
IRQ2Ø	103742	The CPU vectors to IRQ20 when an I/O module interrupts.			
KYIDLE	102425	Keyboard intercept. Polled whenever a key is pressed.			
MSHIGH	103764	High level hook that allows modification of mass storage commands.			
MSLOW	103773	Low level hook to allow driving of mass storage devices not already supported by the system mass storage ROM.			
MSTIME	104002	TIMEOUT hook in the mass storage ROM.			
PLHOOK	103661	If the PLOTTER IS select code is other than one or two, PLHOOK gets called. The contents of R30-R31 determine what routine is executed.			
PRSIDL	1Ø3733	Parser intercept. Should be taken anytime you want to alter the way something is parsed by the system or if the system can't parse something.			
DEF SPARØ	104011	One of the two spare hardware hooks (currently used by the system monitor).			
SPAR1	104022	Second spare hardware hook.			
STRANGE	103715	Parameters for parsing functions are usually numeric, array, or string types. When the system encounters a parameter not of one of these types, it is of type strange. The STRANGE hook is called and parsing this parameter is up to the programmer.			

The hooks RMIDLE, CHIDLE, and IOSP are directly accessed by the executive loop. The following code shows how to take control at these hooks.

#### RMIDLE

Starting at the RAM location 103706, room is allowed to store the following 7 bytes of code:

JSB =ROMJSB 3 bytes - used to select external ROM (if needed).

DEF LABEL 2 bytes - the address of the routine that will be written by the programmer.

VAL ROM# 1 byte - the number of the external ROM that will be accessed using ROMJSB.

RTN 1 byte - return to the executive loop.

Since ROM 0 is usually selected when the system is in the executive loop, external ROMs must go through ROMJSB in order to be selected. Binary programs need only store the following 4 bytes:

JSB =LABEL 3 bytes - subroutine jump to programmer's routine.

RTN 1 byte - return to the executive loop.

The following two pieces of code are examples of how to take control of a hook from a ROM and from a binary program.

### From a ROM:

LDM R41,=316 Opcode for 'JSB ='
DEF ROMJSB Address of the ROMJSB routine
DEF LABEL Address of the hook routine
VAL ROM# Number of the external ROM
RTN Return to the executive loop.
STMD R41,=RMIDLE Store the subroutine jump to the hook routine LABEL at the RMIDLE location.

Since the DRP is set to R41 in the first instruction, seven bytes will be loaded, which will include the 316 (JSB =), the DEF ROMJSB, the DEF LABEL, the VAL ROM#, and the RTN. The code itself will not be executed until the executive loop accesses RMIDLE.

### From a binary program:

LDM R44,=316 DEF LABEL RTN ADMD R45,=BINTAB Opcode for "JSB ="
Address of the hook routine
Return to the executive loop.
Finds the absolute address of the
label LABEL.

STMD R44,=RMIDLE

Store the subroutine jump to the hook routine LABEL at the RMIDLE location.

Here, the 316 opcode, the DEF LABEL, and the RTN are loaded into R44-R47. BINTAB can be safely added to the address LABEL, even though LABEL is a two-byte address and BINTAB is a three-byte address. This is because the most significant byte will be added to the RTN and the most significant byte of BINTAB is always zero and will not affect the RTN opcode. The absolute address of LABEL will always be less than 177400, the limit of binary program memory.

The normal method of returning to the system from RMIDLE is to execute a RTN instruction. Nothing will be on the R6-R7 stack except the return addresses from RMIDLE.

#### CHIDLE

When a key is pressed on the keyboard, the keyboard controller will generate an interrupt request which causes control to pass to the key-service routine. The key-service routine will immediately execute a reset when the [RESET] key is pressed. If no other key is being processed at the same time, the keycode is stored in the location called KEYHIT. The flags are set in XCOM and SVCWRD that indicate that the keyboard is awaiting service for its interrupt. The keyboard controller is reset, and the key service routine returns to whatever it was doing. The next time execution returns to the executive loop XCOM is checked for any pending service requests. If there are any pending requests, the executive loop checks SVCWRD to see which device needs servicing.

In this case the keyboard is the interrupt device, and the executive loop will call the character editor (CHEDIT). CHEDIT will do three things before processing the character input:

- Set binary arithmetic mode.
- 2. Clear the E register.
- 3. Jump control to the location CHIDLE.

At this point you can check the contents of KEYHIT to determine if you want to return to the system or handle the key.

In order for a binary program to handle the key you must pop two return addresses off the R6 stack to insure returning to the executive loop and not to return to CHIDLE or CHEDIT. You must also execute a JSB =EOJ2. This routine clears the bit in SVCWRD which indicates the keyboard needs servicing, and if no other devices have requested service, clears the service request bit in R17.

The status of the E register should also be checked before returning to the executive loop. The E register is cleared by CHEDIT before calling CHIDLE and expects it to be cleared before returning back to CHEDIT. If the E register is nonzero when you return, it assumes that the key pressed was [END LINE] and tries to parse whatever is in the input buffer (INPBUF).

The following section of code illustrates how to take over CHIDLE:

LDM R36, =KEYCHK	Load address of routine to handle CHIDLE.
ADMD R36, =BINTAB	Add value of BINTAB for an absolute address.
STM R36, R45	Store desired address in R45 and R46.
LDB R47, =236	Load the opcode for RTN.
LDB R44, =316	Load the opcode for JSB.
STMD R44, =CHIDLE	Store it all (multi-byte store) to CHIDLE hook.

#### IOSP

When an interface module generates an interrupt, the CPU jumps control to location IRQ20, which is usually taken by the I/O ROM. If IRQ20 has not been taken, the interrupt is ignored. The IOSP interrupt hook is accessed through the executive loop. The I/O ROM IRQ20 routine does minimal interrupt processing and sets the CSTAT and XCOM flags to indicate that an interrupt has occurred. This causes the executive loop to jump to IOSP, where the I/O ROM finishes processing the interrupt. If you take IOSP you must clear the service bit in CSTAT before returning.

### Initialization Hooks

A routine called ROMINI is called on several occasions to perform initialization in external programs. Power-on, allocation, reset, deallocation, and executive loop hooks are times when the binary program may need to initialize special values. When this occurs, the initialization routines in binary programs and ROMs are given control.

A parameter is passed to the ROMINI routine through ROMFL. The occasions and corresponding ROMFL values are:

ROMFL Value	Function		
Ø	Power on		
1	RESET key		
2	SCRATCH		
3	LOADBIN		
4 5	RUN, INIT		
5	LOAD		
6	STOP, PAUSE		
7 CHAIN			
10	Allocate token with class greater than 56.		
11	Deallocate token with class greater than 5		
12	Decompile token with class greater than 56		
13	Program halt because of an error.		

These calls to the ROMs and binary programs allow these programs to initialize or otherwise keep track of operation. For instance, if a ROM needs to reserve or steal memory permanently, it would check for ROMFL =  $\emptyset$ , and reserve memory only when that is true. Another example is that during RESET the I/O ROM might want to deallocate buffers.

During initialization, a binary program or ROM should never destroy any CPU registers below R20. Similarly, no initialization routine should use CPU registers other than R34-R37 until it is verified that the value of ROMFL is not 10, 11, or 12. Once this is verified, all CPU registers numbered higher than 20 may be used.

### Error Handling

When an error is detected inside the executive loop, a system routine immediately reports the error and waits for the error to be corrected. The first 10 (octal) error numbers are default math errors which do not stop execution after the warning is reported. The routine which has found the error supplies a default value, and the processing continues. The defaults must be turned off in order to stop the execution.

The routine that displays the warning message, or sets the error flags if no other errors have occurred begins at location ERROR. When setting an error, the subroutine will use the next byte after the return address as the error number.

The subroutine ERROR has three basic parts to its operation:

- Initializing the error information.
- Interpreting error status.
- · Carrying out the appropriate action.

ERROR saves the address that it will return to in R36-R37, increments it, and stores it on the R6-R7 stack. Then it finds the error number which is stored at the return address and puts it into R20-R21 after saving the previous contents. Checks are made to determine the proper action for the routine. If an error has already been found, then the routine restores the previous contents of the registers and returns immediately. If the error number is less than 10 (octal) or greater than 366 (octal), then the warning for the error is immediately displayed, and the contents of the registers restored before returning. If 'ON ERROR' has been declared and a program is running or if error defaults are off, then the error number, line number, and ROM number (if any) are stored, bits 6 and 7 in XCOM are set to 1, and the previous contents of the registers are restored before returning.

A subroutine jump to ERROR+ is equivalent to a subroutine jump to ERROR followed by a return.

An error condition tested by an assembly language program would go through the following steps:

- The assembly language program finds an error and calls the system routine ERROR.
- 2. ERROR checks to see that no other errors have occurred which haven't been reported yet, in which case ERROR returns without doing anything (because only one error can be in process at a time). Otherwise, ERROR sets the error flags in XCOM and in other RAM locations such as ERRORS, ERLIN#, and ERNUM#.
- Control returns to the assembly language program which returns to the system interpreter.
- 4. The interpreter will check the error flag in XCOM and, noting that it is set, will exit from the interpreter loop back to the main body of the executive loop.
- The executive loop will see that XCOM is not Ø and will see that an error has occurred and will jump to the error-reporting routine REPORT.

- 6. REPORT checks to see if ON ERROR has been declared and a program is running. If so, it sets CSTAT to 'run in middle of line', changes the BASIC program counter to the next line and returns to the executive loop without printing the error message. If a program is not running or ON ERROR has not been declared, then REPORT prints the error message and returns to the executive loop.
- 7. The executive loop checks CSTAT to see if 'run in middle of line' is set. If so, control returns to the interpreter, and the program continues running. Otherwise, ROMFL is set to 13 and ROMINI is called, which is the routine that calls initialization routines in all the external ROMs and the binary programs. When ROMINI returns to the executive loop, CSTAT is set to idle mode.

### 3.8 Extended Memory Controller

Addresses 0 to 177777 (octal) can be directly accessed using 16-bit addressing. The extended memory controller (EMC) is used to access memory locations above 177777. Communication with the EMC, as with the CRT and keyboard controllers, is through the I/O addresses 177400 through 177777. Access to these locations above 177777 is through two pointers, PTR1 and PTR2.

The pointers determine where in memory an access will occur, and since they must access memory locations greater than 177777, they are three-byte quantities. To set the contents of the pointers, a direct store must be performed. For example, STMD R55,=PTR2 will take the three bytes in R55-R57 and move them to PTR2 in extended memory. To store data at the desired location in memory, an indirect store must be performed. For example, STMI R32,=PTR2 will put the two bytes contained in R32, R33 at the address stored in PTR2.

The EMC pointers may be used to create stacks, with the special I/O addresses provided for each pointer. The two pointers are entirely independent of each other. Although PTR2 is used in the following examples, PTR1 and PTR2 function the same.

Each pointer has four I/O addresses: PTR1, PTR1-, PTR1+, PTR1-+, PTR2, PTR2-, PTR2-+. PTR1 and PTR2 act as pointers to memory and must be given a value in order to use the other functions. If data is stored at PTR2, it fills the memory starting at the address stored in PTR2, moving toward the higher numbered addresses.

PTR2- acts as a decreasing stack pointer. A LOAD or STORE through PTR2will first decrement the pointer by the appropriate number of bytes The LOAD or STORE operation will then be performed, leaving the pointer at the new location.

LDM R45,=102,233,114 STMI R45,=PTR2-

BEFORE		AFTER	
1		102	◆PTR2
2		233	
3		114	
4	←PTR2	4	7
5		5	7
6		6	7
7		7	LDM R45, = 102, 233, 114
8		8	STMI R45, = PTR2-
	<b></b>		

PTR2+ is an increasing stack pointer which will perform the load or store operation at the location pointed to by the pointer, and then will increment the pointer after the load or store operation by the appropriate number of bytes.

LDM R45,=102,233,114 STMI R45,=PTR2+

BEFORE		AFTER	
1		1	LDM R45, = 102, 233, 114
2		2	STMI R45, = PTR2+
3		3	
4	◆PTR2	102	
5		233	
6		114	
7		7	←PTR2
8		8	7

When the CPU accesses an I/O address directly, it causes the controller to respond to the address. Each of the controllers is linked to the bus and monitors the information that is being passed from memory to the CPU. For example, the direct access instruction LDBD R32,=CRTDAT will fetch an address from memory. If this address is one which the controller must use for an operation, the controller will send an information byte to the CPU to tell it what to do. In this case the CRT controller will send the CPU the current status of the CRT.

The EMC must constantly monitor the machine code instructions being fetched by the CPU, since the DRP setting determines how many bytes are to be used in a given operation. Whenever a DRP instruction appears, it must store that information to keep track of the current DRP setting.

This can be done with PAD (restore status) and SAD (save status) instructions. SAD pushes three bytes onto the R6 stack containing information about the ARP, the DRP, and the status flags. PAD restores this information using these bytes.

Because of this method of keeping track of the DRP setting, there are cases where the EMC cannot know the DRP setting which include:

 After a PAD instruction: Since the PAD instruction restores status and the ARP and DRP settings, the EMC is not aware of what the DRP setting is until another DRP instruction is executed. Therefore you should avoid using the following or similar code:

> PAD LDMI R#,=PTR2

Restores status, the ARP, and DRP Fetches bytes from extended memory. The CPU assumes the number of bytes is determined by the PADs DRP. whereas the EMC is using the last DRP instruction.

 When the DRP is set indirectly by the contents of CPU register RØ, as in the following case:

LDMI R\*,=PTR2- This sets the DRP according to the least significant six bits of R0, which the EMC knows nothing about.

Because of the first situation, all interrupt service routines must be written to save and restore the contents of the registers used before returning to the routine that was interrupted. Interrupt service routines are those that are called immediately when a hardware interrupt occurs, such as a key being pressed or an I/O module needing attention. Because the interrupt is usually granted almost immediately by the CPU, interrupts can occur between any two instructions (as long as interrupts are enabled). Before restoring everything, you must do the following to solve this problem:

- Pop the SAD status information off of the R6 stack to get a copy.
- Push a copy back on for the eventual PAD.
- · Figure out what the actual DRP needs to be.
- Put the appropriate DRP instruction into RAM along with a RTN.
- Restore all the registers and status (PAD).
- Jump to the DRP and RTN instruction so the EMC will get its DRP pointer back to the right value.

There is another I/O address that the EMC listens to. If you store a l to RULITE (177704), the power light will start blinking. If you store a 0 to that address, the light will stop blinking. This light blinks when a BASIC program is running, or when an HP-85 BASIC program is being translated, or when a program is temporarily halted waiting for input. It normally stops blinking when program execution is complete, if program execution is halted by an error, or if the program is paused.

## 3.9 Parsing

When you type in a BASIC program as a series of ASCII characters it is translated (parsed) and stored internally as a stream of tokens and associated data and addresses. The tokens represent the BASIC reserved words, functions, operators, and punctuation. The data bytes represent the constants, variables, and line number references.

Parsing begins with the line number or the first character of the statement and moves to the right, processing each character and space. Multiple nonquoted spaces are ignored during parsing except those occurring at the beginning of a program line. As a line is parsed, it is checked for syntax errors, changed to RPN (Reverse Polish Notation), and converted into tokens which are stored internally.

Each token consists of a single byte, and can represent a single keyword, such as LET or PRINT. Tokens 370 (ROM token) and 371 (binary program token) are used to allow extensions of the system by means of external ROMs and binary programs. A table of system tokens can be found in section 8. ASCII codes can be found in the HP-87 owner's manual.

Example: In parsing the line

10 LET A = B \* SIN (45),

the system produces the following tokens in the order shown.

Tokens (Octal Value)	Comments
16	End of statement.
1Ø	Store numeric value token.
52	Multiply token.
33Ø	Sine token.
1Ø5	BCD 45 in integer format.
Ø	(Refer to paragraph 3.9, Numeric
Ø	Formats.)
32	Integer constant token.
102	ASCIÍ "B", variable name.
1	Length of variable name.
Ø	Variable address space for allocation.
0 0 0 1	(Refer to Format of BASIC Programs and
Ø	Variables, paragraph 3.12.)
1	Fetch simple numeric variable
101	ASCII "A"
1	Length.
Ø	Variable Address Space
Ø Ø Ø	The state of the s
Ø	
21	Store simple numeric variable token.
142	Let token.
25	Length of line in bytes.
20	Line number in BCD (two digits per
	byte except for most significant
Ø	byte which contains only one).

The extended memory pointer, PTR2, is used as the output pointer during parsing. Tokens are stored indirectly to PTR2-. At the beginning of the parsing process PTR2 is set equal to NXTMEM, so the parsed line will be built up in available memory at the end of the last BASIC program. Parsing begins with the line number. This is loaded in BCD form; 20 is loaded first, since it is the least significant byte.

Next is the size or length of the statement. During parsing this is a blank place holder byte; STSIZE is a pointer to the place holder byte. In order to find a match for the keyword LET, the system looks first in keyword tables in the resident binary programs, then in any external ROMs, and finally in the internal system keyword table. For this reason, a binary program or external ROM can take over any keyword (that is, a binary program can implement a custom version of PRINT, while the preprogrammed PRINT is ignored). The extend register indicates if the token searched for has been found. Refer to the section on status indicators in paragraph 2.3.

After parsing, if the statement was a program line, its tokens and addresses are inserted into the program space at the correct locations. If it was an expression or calculator mode statement, the parsed code remains at the end of the BASIC program and is executed immediately, being discarded when execution is finished.

For further details of parsing operations and register conventions at parse time, along with specific parse routines, refer to the system routines which are listed in alphabetical order in section 8.

### 3.10 Decompiling

Programs or statements are decompiled as they are listed. This is the reverse process of parsing and compiling. Internally, it requires the reconstruction of code as it was entered. The tokens which have been parsed into RPN and distributed in the system are reassembled.

PTRl points to the input stream, which is accessed by loading indirect through PTRl-. Input is then decompiled to an expression stack or an output stack. The expression stack (R12) is used to reconstruct expressions from RPN to their original form, and an output stack (pointed to by R30) is used to buffer the output.

Since the tokens are arranged in RPN internally, the system decompiles the tokens as it pushes missing operator tokens (016) onto the expression stack. These missing operator tokens are merely "place holders" until the arithmetic operators can be inserted at a later step.

Unlike parsing, decompiling is not an operation to which a binary program or ROM normally has access, since these programs are seldom required to perform any unique operations during decompiling. In some special cases the parse routines for a binary program or ROM may require modification if a statement is to be decompiled correctly. But for the most part, decompiling will not be a problem for the writer of binary or ROM programs.

The system processes each token and uses its class (a component of the token's primary attributes) to determine how the token is to be decompiled. Following are some common classes and how they are decompiled:

Class	Type of Token	Action
Ø	End-of-line	Unstack.
1	Fetch variable	To expression stack.
2	Integer	To expression stack.
1 2 3 4	Store variable	To expression stack.
4	Numeric constant	To expression stack.
5	String constant	To expression stack.
32	Subscript, such as, A(3)	<ul><li>() to expression stack if token odd; otherwise (,) to expression stack.</li></ul>
34	Dimension subscript like, A\$[ ]	[] to expression stack if toker odd; otherwise [,] to expression stack.
36	Prints	Unstack and push to output.
41	Other reserved words	If : then unstack, output reserved word, then unstack.
42	Miscellaneous output	If , then push to expression stack and unstack; otherwise output.
44	Miscellaneous ignore	Ignore.
50	Unary operator	Insert after most recent missing operator in expression stack.
51	Binary operator	Replace most recent missing operator in expression stack.
52	String unary operator	Same as class 50.
53	String binary operator	Same as class 51.
55	Numeric function	Replace the most recent missing operator with "," for each parameter. Then insert function name (at most recent missing operator) and push onto expression stack.
56	String system function	Same as class 55.

The following example illustrates how decompiling occurs:

## 10 LET A=B\*SIN(45)

After being parsed as shown, these tokens are decompiled into the output stack and the expression stack as illustrated.

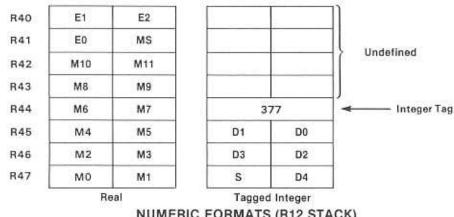
STEP	TOKENS	R12 EXPRESSION STACK	OUTPUT BUFFER	STEP	TOKENS		R12 EXPRESSION STACK	OUTPUT BUFFER
9.	16 EOL		10 LET A = B * SIN(45)					
						16		
		16					A	
		A		12	105	ve:	16	40.00
		( E		5.	105 0 0 32	45	8	10 LET
8.	10=	В	10 LET		32 )		16	
		*					45	
		SINI						
		45						
		).			102)		16	
				4.	1 0 0	В	A	
					0		16	
		16			1 )		В	
		A						
		16			101		16	
7.	52 *	В	10 LET	3.	0 0	Α	A	
		3000			0			
		SIN(			21 7			
		45						
		1		2.	142 LET		6	10 LET
		16		1.	25 20	LEN		
		A				LINE	#	10
		16			5616			
6. 33	330 SIN	B	10 LET					
		16						
		SIN(						
		45						
		0						

### 3.11 Operating Stack

The stack to which R12 points is used for passing values in many internal system routines. The formats of values that are fetched and stored during run time execution of certain specific tokens, as well as the formats of numeric quantities are in this section.

### Numeric Formats

In internal routines, numbers popped off the R12 stack are eight bytes long, so integer values are tagged with octal 377.



NUMERIC FORMATS (R12 STACK)

In the illustration, the byte above the number contains the octal quantity 377. This 377 acts as a tag for the number, specifying the quantity as an integer value that is only three bytes in length. next four bytes popped off the stack are then undefined and are ignored by the system. The numbers are shown as they would be if they were taken off of the stack by the instruction POMD R40,-R12. The tagged integer is right justified so that the most significant digits (starting with D4) are Ø if unused. For tagged integers, the decimal point is to the right of D0, the least significant digit. The real number decimal point is between MØ and Ml.

A short numeric variable is formatted as follows:

R44	E0 .	E1
R45	мз	M4
R46	M1	М2
R47	0 0 SM SE	мо

EØ	Most significant four-bit BCD digit of the exponent.
El	Least significant four-bit BCD digit of the exponent.
MØ	Most significant four-bit BCD digit of the mantissa.
M4	Least significant four-bit BCD digit of the mantissa.
SM	Sign of the mantissa (0=positive, 1=negative).
SE	Sign of the exponent.

The decimal point is assumed to be between digit MØ and digit Ml. The most significant nibble (four bits) contains the signs of the mantissa and the exponent. The two most significant bits are zeroes.

Strings on the R12 Stack

String values are passed on the operating stack as a two-byte length and a three-byte address of the next character higher than the first character of the string. The first character is at the highest address of any characters of the string. To fetch successive characters of the string, the following code could be used:

```
POMD R45,-R12 ! Get the address of $
STMD R45,=PTR2- ! Set PTR2 pointing to first character
POMD R36,-R12 ! Get the length of $
LOOP LDBI R32,=PTR2- ! Get the next character

.

DCM R36 ! Decrement length count
JNZ LOOP ! Loop until done
```

Operating Stack Routines

There are several system routines available to help you in parsing various kinds of parameters for BASIC statements. These routines will parse your BASIC statement into tokens that, at run time, will load the R12 stack with the appropriate variable or parameter.

Following is a list of the routines that can be used and what they leave on the stack:

NUMCON (8 bytes) Real or tagged integer.

NUMVAL (8 bytes) Real or tagged integer.

REFNUM (3 bytes) Absolute address of variable value.

(3 bytes) Absolute address of name of variable.

(1 byte) Head of variable.

STRCON (2 bytes) Length of string.

(3 bytes) Absolute address of string.

STREXP (2 bytes) Length of string.

(3 bytes) Absolute address of string.

STRREF Will parse both a normal string variable and a string array variable reference. There will be slightly different information on the stack depending on which of these it is. String arrays will have everything that nonarray strings will have but string arrays may also have row, column, and dimension information if the variable is being traced. You can tell if that information is there by checking the trace bit in the header byte which will come off the stack before the tracing information would. You also tell whether you have a string array or normal string by inspecting the appropriate bit in the header byte.

### Nonarray Strings

(3 bytes) Absolute address of name of variable.

(1 byte) Header of variable.

(2 bytes) Maximum length of string variable.

(3 bytes) Absolute address of first byte of string address.

(2 bytes) Maximum length available to store into. This will be different from the maximum length if subscripts were used.

(3 bytes) Absolute address of first byte to store into. This will also be different from the address of the first byte of the string if subscripts were used.

### Array Strings

The first three will only be on the stack if the variable is being traced.

- (2 bytes) Row of element.
- (2 bytes) Column of element.
- (1 byte) Dimension flag (Ø=2 dim., 1=1 dim.).
- (3 bytes) Absolute address of name of variable.
- (1 byte) Header of variable.
- (2 bytes) Maximum length of string variable.
- (3 bytes) Absolute address of first byte of string variable.
- (2 bytes) Maximum length available to store into. Different than maximum length of variable if subscripts were used.
- (3 bytes) Absolute address of first byte to store into. Different from the address of the first byte if subscripts were used.

NARREF

Used when you wish to use a simple numeric variable name to refer to an array variable. An example would be:

### MAT C=ZER

In this example 'C' refers to an array C, not to a simple numeric variable.

(3 bytes) Address of variable header. This address is a relative address. The easiest way to make it an absolute address is:

POMD R65,-R12 JSB =FETSVA

FORMAR Used when you wish to refer to an entire array.

### PRINT# 1; C(),D(,)

- (3 bytes) Absolute address of the first element of the array.
- (3 bytes) Absolute address of the array name.
- (1 byte) Array header.

In all of the above examples of stack contents, the bottom of the page represents the direction of higher addresses. As you popped things off the stack you would be removing things from the bottom first.

#### 3.12 Format of BASIC Programs and Variables

The following figure shows how a BASIC program line is formatted:

END OF LINE TOKEN MISCELLANEOUS BYTES (ACTUAL TOKENS OF CODE) 1-BYTE LENGTH OF LINE 3-BYTE BCD LINE NUMBER (5 DIGITS)

The BASIC line

15160 END

would be parsed as:

016 'END OF LINE' TOKEN 212 'END' TOKEN 002 LENGTH OF LINE (212 AND 016 MAKES TWO BYTES) 140 121 --> 3-BYTE BCD LINE #

001

Let's take a look at how a line number of 15160 generates the three bytes 140, 121, and 001. Since a BCD digit takes four bits, two digits can be packed into one byte. So, let's split the line number into three digit groups:

> 51 1 60

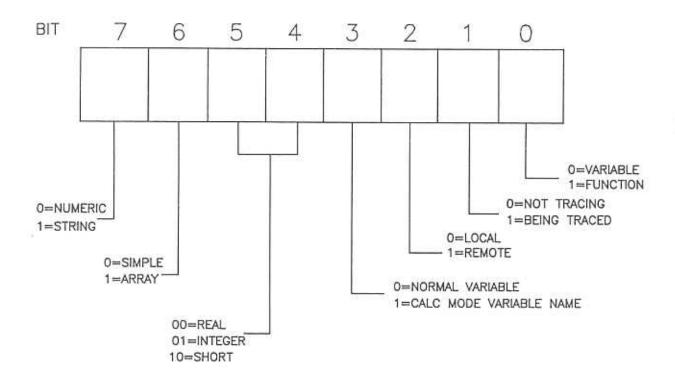
Now we turn those groups into bits:

0000 0001 5 1 0101 0001 6 0 0110 0000 Arrange the binary representation with three to a group. Convert this form to an octal number to obtain the three bytes that represent the line number.

00 000 001 001 01 010 001 121 01 100 000 140

The values of the variables are stored at the end of the current program in one continuous block of memory. Each variable has a header which contains information about that variable. Following are the structures of different kinds of variable storage areas. All variable storage areas begin with a one byte header. The bits in that header and their meanings are:

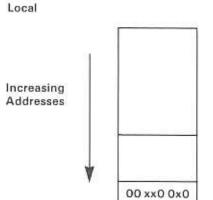
## VARIABLE HEADER BYTE LEGEND



## Section 3: Operating System

In the following diagrams in this section, an "x" will mean that that particular bit position can be occupied by a "l" or a "0."

Simple Numeric Variable



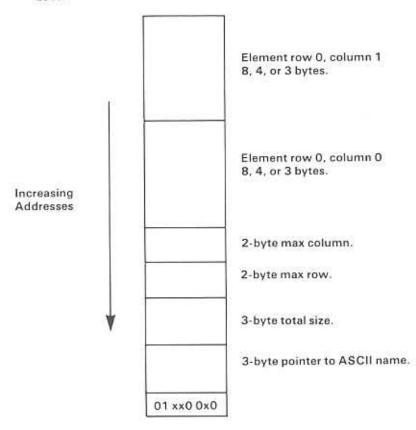
8, 4, or 3 bytes of value depending upon whether it's REAL, SHORT, or INTEGER.

3-byte pointer to ASCII name.

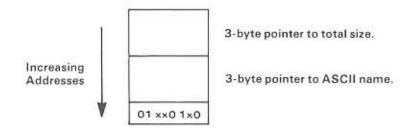
Remote		
Ť		3-byte pointer to value.
Increasing Addresses		3-byte pointer to name.
*	00 xx0 1x0	

## Numeric Array

Local

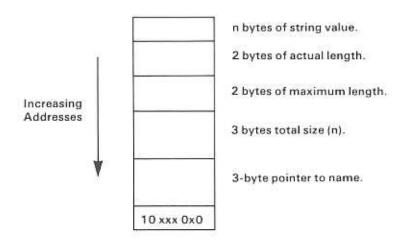


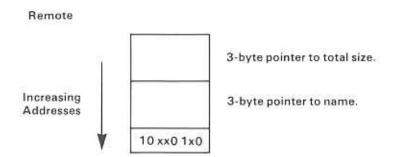
Remote



# Simple String Variable

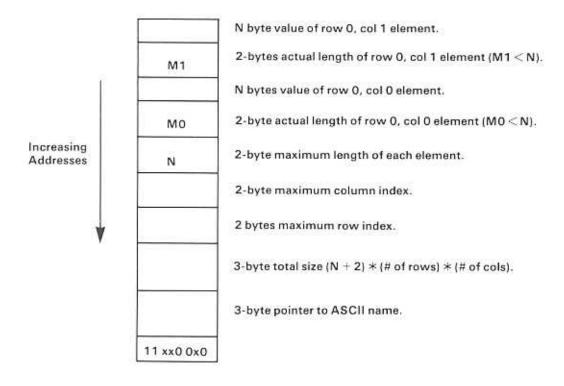
Local



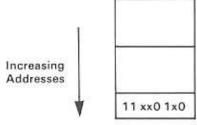


## String Array Variable

#### Local



Remote



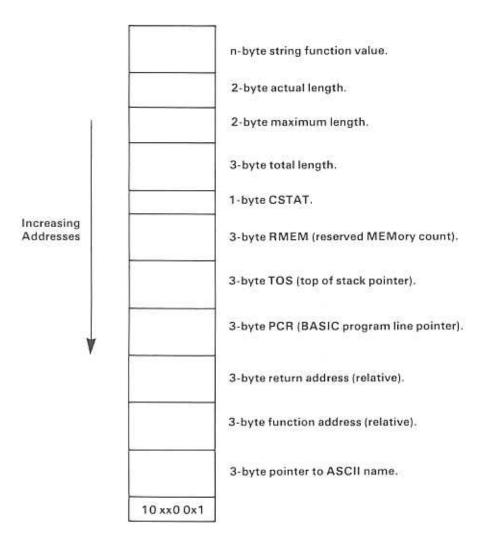
3-byte pointer to total size.

3-byte pointer to ASCII name.

Numerical User Defined Functions

A	
	8-byte function value.
	1-byte CSTAT.
	3-byte RMEM (reserved MEMory count).
	3-byte TOS (top of stack pointer).
	3-byte PCR (BASIC program line pointer).
	3-byte return address (relative).
	3-byte function address (relative).
	3-byte pointer to ASCII name.
00 xx0 0x1	
	00 xx0 0x1

# String User Defined Functions



Because calculator mode statements destroy all previous calculator mode statements but not their variables, the pointers to the ASCII names of the variables cannot point to the calculator mode statement. A dummy calculator mode simple string variable is created with the bit set in the header that indicates this is a calculator mode variable name. This dummy variable is skipped for all purposes other than searching for variable names at allocation time for calculator mode statements. When a calculator mode statement is allocated, the addresses used for the variables are relative to FWCURR.

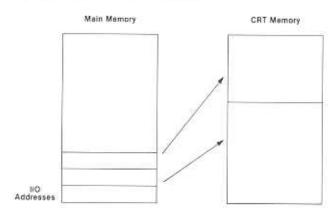
#### CONTROLLERS

#### 4.1 Introduction

The HP-87 is a multi-processor system. The keyboard, the CRT, the timers, and the interface modules are all controlled by individual microcomputers. The mainframe CPU coordinates activities between the peripherals using the I/O addresses. To communicate with these controllers, refer to the appropriate sections.

#### 4.2 CRT Controller

The CRT is an intelligent component that is controlled by an internal computer, or CRT controller. The CRT also has a memory which continuously refreshes the CRT display.



The CRT controller and the CPU communicate using four addresses in RAM. Each address requires a two-byte quantity to specify a CRT memory address. The I/O addresses are:

## CRTBAD DAD 1777Ø1

Storing a two-byte address to this location causes the CRT controller to load its byte address pointer with that address.

#### CRTSAD DAD 177700

Storing a two-byte address to this location causes the display to be started at that address. This makes the display appear to scroll up and down or side to side or to jump to a different page depending on the new start address. Storing to CRTSAD has no effect when in GRAPH NORMAL or GRAPH ALL modes.

#### CRTDAT DAD 177703

Storing a single byte to this location causes that byte to be stored to the CRT memory location currently pointed to by the controllers byte address. Loading a single byte from this location reads the byte from the CRT memory location currently pointed to by the controller's byte address.

After either a load or store operation through CRTDAT the CRT controller automatically increments by one its internal byte address pointer. If you did a series of single byte store instructions to CRTDAT without storing anything to CRTBAD in between, those bytes would be stored in successive CRT memory locations.

However, before storing to CRTDAT, you must first read CRTSTS and check the least significant bit to make sure the controller is not busy. Before loading from CRTDAT, you must store a byte to CRTSTS with the least significant bit set to tell the CRT controller that you want to read the current memory location. You must then read CRTSTS until the BUSY bit indicates the controller is not busy, at which point you can load from CRTDAT to get the byte. An easier way is to simply execute a JSB =INCHR (call the system routine) that does all the rest for you.

## CRTSTS DAD 177702

Loading a single byte from CRTSTS gets you information about the current status of the CRT controller. Each bit has a specific meaning:

Reading from CRTSTS

Bit	Ø	1
Ø	Not Busy	Busy
1	Unblank	Blank
2	Power-up	Power-down
3	16 lines	24 lines
4	Display time	Retrace time
5	Noninverse	Inverse Display
6	Normal	All
7	Alpha	Graphics

Storing a single byte to CRTSTS sets the CRT controller to a specific mode and/or requests a read:

Storing to CRTSTS

Bit	Ø	1
Ø	No read	Read Request
1	Unblank	Blank
2	Power-up	Power-down
3	16 lines	24 lines
4	-	<del>-</del>
5	Noninverse	Inverse Display
6	Normal	A11
7	Alpha	Graphics

When the CRT is blank, the controller has disabled the electron beam, causing the display to go blank. When this is the case, the controller does not have to refresh the display, causing it to transfer data to and from the CRT memory much faster. When you switch from alpha to graphics or graphics to alpha there will be a flash on the display unless it has been blanked. To avoid this, you must set the blank bit during a retrace. There is a pair of system routines that will blank and unblank the CRT for you. They are CRTWPO and CRTUNW.

When the CRT is powered up or powered down, the controller turns the high voltage section of the CRT driver on or off. This is done to conserve power.

# 4.3 Display Modes

ALPHA NORMAL

OCTAL		
ADDRESS		
000000		
000120		
000240		
007760		
010100		
010220	last line of ALPHA NORMAL memory	
010340	start of GRAPH NORMAL memory	

ALPHA addresses in CRT memory are 000000 to 010337. In alpha mode the display shows 16 or 24 (decimal) lines of 80 (decimal) characters per line. The scrolling keys permit viewing of an additional 38 (decimal) lines of alphanumeric data.

Because each ASCII character occupies eight bits, one character can be stored at each memory location. To move the cursor to the right one position, add one to the address.

If the start address (CRTSAD) is at an address where there is not enough ALPHA NORMAL memory left for an entire display, then the CRT controller will start fetching bytes from address 000000 when it reaches the end of the ALPHA memory. Because of this a mod operation must be performed on alpha addresses when moving the cursor around.

At power-on and after a RESET the CRT start address is set to 00000. If you roll the display up one line the CRT start address will then be set to 00120. If you were to roll the display down one line, the start address would be 10220.

GRAPH NORMAL

first (top) line of GRAPH NORMAL display
last (bottom) line of GRAPH NORMAL display
last 64 bytes are unused.

In GRAPH NORMAL mode, the screen is 50 bytes (decimal) wide. The GRAPH display always starts at 10340.

The last 64 (decimal) bytes of CRT memory are unused in NORMAL mode. The contents of memory location 10340 will determine whether or not the first eight dots in the top line of the display will be on. The contents of memory location 10341 will determine the state of the next eight dots on the top line.

## ALPHA ALL

ADDRESS	E- (F (ALPHA ALL	
000000	first line of ALPHA ALL memory	
000120		
037440		
037560	last line of ALPHA ALL memory	

The ALPHA ALL memory maps 80 addresses per line of the CRT display and the last 64 (decimal) bytes of memory are unused. When the start address gets too close to the end of memory, the controller wraps around to address 000000 to finish the display page.

GRAPH ALL

OCTAL		
010340	first (top) line of GRAPH ALL	display
010444		evi-av-
037640		
037744	addresses 37744-37777	addresses 0-47
000050		
000154		*****
010134	last line of GRAPH ALL displa	ау
010240	last 64 bytes are unused	

In GRAPH ALL mode there are 68 (104 octal) bytes per line of the graphics display, giving a dot resolution of 544 dots wide by 240 dots high. The controller will again wrap back to address 000000 to continue fetching bytes when it runs out of memory at the end of the NORMAL graphics area.

# 4.4 Keyboard Controller

The keyboard controller monitors the RAM location keyboard scanner, four timers, and the beeper.

# Keyboard Scanner

All of the keys are connected to keyboard inputs. The controller monitors these connections, waiting for a key to be pressed. When a key is pressed, the controller generates a service request to the CPU. When the request is granted execution vectors to the service routine KEYSRV. The keyboard service routine saves the CPU status then does a JSB=KYIDLE instruction (refer to Hooks, paragraph 3.5). If the KYIDLE hook has not been taken, control will return to KEYSRV. It will then disable interrupts, save registers, and read the key code of the key that was pressed from the keyboard controller through the I/O address KEYCOD. The key is checked by KEYSRV to see if it was RESET. If so, KEYSRV does a RESET. If not, it checks to see if any other keys have been pressed that have not been handled by the system.

If another key has been pressed, the system re-enables the keyboard scanner and restores the registers and status. The system returns to what it was doing when the CPU received the service request. As long as other keys are not pending, the key code is saved in a RAM location called KEYHIT and bits are set in R17 and SVCWRD, indicating that a key has been pressed. The routine KEYSRV then restores the registers and status.

Once a key has been pressed, no more keyboard interruptions will be seen until the previous key is released, and a 1 has been stored to I/O address KEYCOD (which restored the keyboard scanner). If the interrupt were to occur between the last DRP instruction and an extended memory access, the EMC could lose track of what the DRP setting is. Refer to paragraph 3.6.

The following define the I/O addresses associated with the keyboard scanner:

#### KEYSTS

Write:	Bit	Ø	1
	Ø	No effect	Enable keyboard
		No effect	Disable keyboard
	2	Not used	=
	3	Not used	-
	4	not used	-
	1 2 3 4 5 6 7	Speaker off	Speaker on
	6	No effect	1.2 kHz
	7	No effect	Toggle Flip FF
Read:	Bit	Ø	1
	Ø	Device disabled	Device enabled
	1	No key pressed	Key pressed
	2	Not used	-
	3	Shift key up	Shift key down
	4	Not used	78
	2 3 4 5 6 7	Not used	E70
	6	Not used	=
	7	Globals disabled	Globals enabled

Bits 0 and 1 of KEYSTS allow you to disable and enable the keyboard separately from all other devices. Bit 7 (when reading) tells you whether global interrupts are enabled or disabled.

## KEYCOD

The status of KEYCOD utilizes a byte rather than individual bits.

Write: If the value is 1, then the keyboard scanner will be re-enabled as soon as the key is released.

Read: Returns the keycode of the key that was pressed.

Following is a listing of the system key service routine, KEYSRV, presented here as an example of what you need to do if you take over KYIDLE.

5880 5890 5910 5910 5920 5930 5940 5950	KEYSRV	JSB =KYIDLE   STBD R32, =GINTBS   PUMB R32, +R6   LDBD R32, =KEYCOD   BIN   CMB R32, =213	Save the STATUS, ARP, and DRP Call the RAM hook Disable global interrupts Save register contents to recall later Get the keycode from the controller IC Force BIN mode for keycode compare Is it the RESET key? Jif no
	RSTART	LDM R6, =STACK ! JSB =RESET. !	Else reset the return stack pointer Do a RESET Need to store a 1 out to KEYCOD to restart the keyboard scanner Output cursor to CRT and fall into exec.
120001200	NORSET	LDBD R33, =SVCWRD ! J0D HAVE1 ! ICB R33 STBD R33, =SVCWRD !	Any other unserviced keys been pressed?  Jif yes, throw this key away  Else set the keyboard bit  And restore SVCWRD  Save the keycode for the system
		LDB R32,=20 ! ORB R17,R32 ! JSB =E0J1 ! LDB R#,=1 ! STBD R#,=KEYCOD !	Load the mask to set service request bit in XCDM (R17)  Make sure we're set to slow repeat speed Can't get any more keys unless we restart the keyboard scanner Restore the registers we used (R32-R33) Make the current DRP setting available to the EMC.
6140 6150 6150 6170 6180	E0J1		Get the slow repeat count Set the counter to the slow repeat
6200 6210 6220 6230 6240 6250 6260 6270		POMD R10,-R6   PUMD R10,+R6   R10,+R6   R10,+77,0   RDB R10,+100   RTBD R10,+RAID+1   LDMD R10,+S10   R10,+S10	Save R10-11 in a reserved RAM location Get the byte of SAB that contains DRP Restore so we can PAD later Isolate the DRP register bits Make it a DRP instruction Store it into RAM so we can execute it Restore R10-11 Finish

## Section 4: Controllers

At power-on, the system initialization routine has stored at RAID+1 the following code:

RAID+1 BSZ 1 STBD R#,=GINTEN PAD RTN Place holder for DRP instruction Re-enable global interrupts Restore status, the ARP, and the DRP Done

#### Timers

The timer section of the keyboard controller consists of four separate timers and four registers each containing eight BCD digits. The timers and registers are updated at a rate of 1 kHz. During this updating, no read or write operations should be performed to the CLKDAT address. Each timer that equals its register count causes a service request interrupt. It is then set to zero to begin another count sequence. The contents of the timers are transferred in four consecutive bytes each containing two BCD digits.

The keyboard scanner has the highest priority on the controller regarding interrupts. Next highest is timer  $\emptyset$ , with timer 3 being the lowest.

## CLKSTS

This address contains the following information needed when using the timers.

Write: Bit	Comments
Ø	Disable addressed timer.
1	Enable addressed timer.
2	Stop addressed timer.
3	Start addressed timer.
4	Clear addressed timer.
1 2 3 4 5	Clear interrupt service flip FF.
6	Bits 6 and 7 are the
7	timer address (0 through 3).
Read: Bit	Comments
Ø	Timer Ø enabled.
Ø 1	Timer 1 enabled.
2	Timer 2 enabled.
3	Timer 3 enabled.
4	Not used.
2 3 4 5 6 7	Not used.
6	Not used.
7	Read (timers available for
	access through CLKDAT).

## CLKDAT

When loading from CLKDAT, you must execute a four-byte load to get eight BCD digits which represent the value of the last addressed timer.

When storing to CLKDAT, you must execute a four-byte store and the four bytes must be the eight-digit value you want the last addressed timer set to.

Before executing a load or store instruction to CLKDAT you must first check the most significant bit of CLKSTS to make sure the timers are ready to be accessed (bit 7=1).

RTN

There are no hooks in the timer interrupt routines. The only way to make use of the timers from assembly language programming is to periodically check SVCWRD to see if any timers have been interrupted. This will only work if you never return to the BASIC interpreter, as the executive loop will also check for timer interrupts at the end of each BASIC statement and handle them if necessary.

The following code will read the value of timer  $\emptyset$  (the system clock). It will use that value and the base time to generate the current time and return the current time to the R12 stack.

TIME.	CLB R55	! ADDRESS TIMER Ø
	STBD R55,=GINTDS	! DISABLE INTERRUPTS
	JSB =TIMWST	! WAIT FOR READY AND STORE
		! TIMER ADDRESS
	CLM R40	! CLEAR UPPER FOUR BYTES
	JSB =TIMRDY	! WAIT FOR READY
	LDMD R44, =CLKDAT	! TIME TO R44-R47
	STBD R44,=GINTEN	! RE-ENABLE INTERRUPTS
	LDM R36,=4,0	! SET EXPONENT
	BCD	
	CLB R32	! SET SIGN TO POSITIVE
	JSB =SHRONF	! SHIFT, PACK AND PUSH
		! ON R12 STACK
	LDMD R50,=TIME	! GET BASE TIME
	POMD R40,-R12	! RECOVER INITIAL TIME
	JSB =ADD10	! ADD TO BASE TIME AND
		! PUSH ON R12 STACK
	RTN	
TIMWST	JSB =TIMRDY	! WAIT FOR READY
	STBD R55,=CLKSTS	! STORE OUT STATUS BYTE
	RTN	
TIMRDY	LDBD R37,=CLKSTS	! GET TIMER STATUS
	JPS TIMRDY	! JIF BUSY
	-4.506.00.00	500 BECHELLER BECHELLER DE

The system routine SHRONF takes a 16-digit number in R40-R47, an exponent in R36-R37 and a sign byte in R32 and normalizes it (shifts out leading zeroes and adjusts the exponent to match). It then packs the exponent and sign into R40-R41, and pushes the floating point result onto the R12 stack. The ADD10 routine is basically the same as ADDROI except it expects as inputs two real (floating point) numbers in R40-R47 and R50-R57, rather than two real or integer numbers on the R12 stack.

1 ELSE RETURN

## Section 4: Controllers

The following code sets timer  $\emptyset$  (the system time clock) the way it's set at power-on.

TIMEØ LDB R55,=32 ! SET UP STATUS BYTE. ! BITS 4, 3, 1 WILL CLEAR ! TIMER O. ! START IT, AND ENABLE IT TO ! INTERRUPT. CLM R44 ! GENERATE 86400000, THE ! NUMBER OF MILLISECONDS LDM R46,=100,206 ! IN A DAY. STBD R#,=GINTDS ! DISABLE INTERRUPTS. JSB =TIMRDY ! WAIT FOR READY. ! SEND THE TERMINAL COUNT. STMD R44,=CLKDAT STBD R44,=GINTEN ! RE-ENABLE GLOBAL INTERRUPTS. RTN ! DONE.

# Speaker

The speaker can be controlled through the I/O address KEYSTS. Bits 5 and 6 of KEYSTS allow you to either make the speaker beep at 1.2 kHz or turn it off and on at whatever frequency you wish (within the limits of the clock cycle of the CPU).

#### SYSTEM MONITOR

## 5.1 Introduction

The HP 82928A System Monitor is an optional plug-in module that permits you to set breakpoints and single step or trace through the execution of assembly code. Two breakpoints can be set in any portion of memory with an address lower than 200000. Any time either of these addresses is referenced in any manner, an interrupt is caused. The user can use this interrupt to examine CPU registers, status bits, memory locations, and extended memory pointers.

# 5.2 System Monitor Commands

The system monitor commands described in this section are demonstrated later in this manual. Refer to section 7.

BKP octal address [,select code for output]

Sets breakpoint (BKP) #1 or #2 at a specified address in memory. If no breakpoints have been set, the command sets BKP#1. If BKP#1 is already set, the command sets BKP#2. If BKP#1 and BKP#2 are both set, the command resets BKP#2; BKP#1 remains set at its original address. Breakpoints can be set at any address lower than 200000 in system RAM or ROM. They can be cleared only by using the CLR command. Using the [RESET] key will not clear breakpoints.

When a breakpoint is encountered, execution halts and a block of status information is output to the CRT IS device. The following keys are typing aids:



Key	Use
B C	Set an additional breakpoint (BKP)
C	Clear (CLR) a breakpoint.
M	Obtain a memory dump (MEM).
P	Change program counter (PC=).
R	Change contents of a register (REG).
$\mathbf{T}$	Using the TRACE command.
1	Change value of pointer #1 (PTR1=).
2	Change value of pointer #2 (PTR2=).
[STEP]	Single step execution.
[ROLL ^]	Roll up display.
[ROLL V]	Roll down display.
[RUN]	Resumes program execution.
[BACK SPACE]	Back space.
[A/G]	Alternates between graphics and alpha modes.

Most other keys on the keyboard are inactive at a breakpoint until a typing aid has been used.

									E 01							TR2 77732	R0	530		
	0		2	3	4	5	6	7	MEM	0:0										
ROO	000	012	265	230	273	044	150	204	026	000	011	210	303	030	011	210			C	
R10	242	200	350	212	371	000	001	000	153	031	305	031	266	031	247	031	ke	Ε	6	
R20	044	044	233	230	140	011	236	200	342	207	022	210	022	210	106	251	b			F)
R30	237	200	034	000	075	210	320	230								251	8	8	2	H )
R40	110	233	230	001	000	000	044	044				030						2	10	3
R50	000	051	000	000	000	000	000	000	366	012	262	065	210	261	014	140	v	25	1	
R60	000	000	000	000	000	000	357	012	036	306	000	000	316	274	011	316	F		ΝĈ	N
R70	016	316	000	000	000	000	000	000	030	030	230	316	306	207	117	220		N	F	0

Output at a breakpoint includes:

1. The following CPU status indicators:

PC: The setting of the program counter stored in registers R4 and R5. When execution is resumed, it will begin at the address specified by the PC.

DR: Contents of the current data register pointer.

- AR: Contents of the current address register pointer.
- OV: Status of the overflow flag.
- CY: Status of the carry flag.
- NG: Status of the MSB (most significant bit), used to indicate a negative quantity.
- LZ: Status of the LDZ (left significant zero) flag.
- ZR: Status of the zero flag.
- RZ: Status of the RDZ (right digit zero) flag.
- OD: Status of the LSB (least significant bit), used to indicate an odd quantity.
- DC: Setting of DCM (decimal) flag. Used to indicate decimal or BCD mode.
- E: Contents of the E (extend) register. This will be a quantity between 0 and 17 octal.
- BKPl: Indicates absolute address where breakpoint 1 is currently set.
- BKP2: Indicates absolute address where breakpoint 2 is currently set.
- PTR1: Indicates address of extended memory pointer 1.
- PTR2: Indicates address of extended memory pointer 2.
- ROM: Indicates number of ROM which was selected when the breakpoint occurred.
- 2. The contents of 100 (octal) RAM or ROM locations are output beginning with the octal address specified in the last executed MEM and will continue for 100 octal bytes. If no MEM was executed, 100 (octal) bytes of memory will be output beginning with zero. The default ROM number is zero unless previously indicated. If MEM was executed, 100 octal bytes will be output starting with the address of MEM.
- 3. Contents of CPU registers 0 through 77.
- 4. Memory contents in ASCII.

# CLR breakpoint number

After CLR is displayed (as a result of typing "C"), the user can type 1 [END LINE] to clear BPl or 2 [END LINE], to clear BP2. After CLR is displayed pressing [END LINE] or typing a number other than 1 or 2 will clear both breakpoints.

The CLR functions can be used any time execution has been halted, whether or not it has been halted by a breakpoint.

MEM address [:ROM#][,# of bytes][=#,#,...]

This command dumps the contents of computer RAM or ROM memory to the current CRT IS device beginning with the octal address selected. One-hundred octal bytes are dumped unless another parameter was input. The MEM function can be used after execution has been halted by a breakpoint.

The ROM number if included, is an octal value of selected plug-in ROMs from which memory is dumped. Default value for the ROM number is system ROM  $\emptyset$ , if no other ROM number has been selected.

The output shows the octal representation of the bytes in memory and the ASCII representation of the bytes.

If there are numeric entries after the "=" sign, memory is not dumped; the contents of memory locations beginning at the octal address specified are changed to the octal values after the "=" sign. The memory locations must be in RAM. The contents of one succeeding memory location are changed for each value specified after the "=" sign. The number of bytes, if included is disregarded in this case.

## Examples: MEM 103300

Dumps contents of 100 octal bytes of memory to the CRT IS device, beginning with memory location 103300.

MEM 103300,20

Dumps contents of 20 octal bytes of memory to the CRT IS device, beginning with memory location 103300.

MEM 60200: 40,200

Dumps contents of 200 bytes of the assembler ROM (ROM #40) to the CRT IS device, beginning with memory location 60200.

MEM 105000 = 0,0,0,15

Loads memory locations 105000, 105001, and 105002 with zeros, and loads location 105003 with 15 octal.

PC= octal address

Changes contents of program counter stored in CPU registers R4 and R5 to the specified address, and dumps CPU status and memory contents exactly as when a breakpoint (BKP) is executed. When execution is resumed, it will begin at the address now specified by the contents of the program counter (PC).

Example: PC = 3477 (Sets the PC to resume execution with byte 003477.)

REG number of CPU register = value

Changes contents of specified CPU register to the value given, and dumps CPU status and memory contents as when a breakpoint (BKP) is executed. Value may be octal, decimal, or BCD.

Example: REG 35 = 31 (Changes contents of register R34 to 31 octal.)

REG 36 = 19C (Changes contents of register R36 to BCD 19.)

REG 37 = 25D (Changes contents of register R37 to 25 decimal.)

STEP

Although STEP is not a command, it is a typing aid which executes the next complete machine code instruction (not just the next byte). Beginning with the location currently addressed by the PC, it halts and dumps CPU status and memory contents like a breakpoint.

TRACE octal, decimal, or BCD value

Resumes execution with the next machine code instruction, and continues for the number of instructions (not bytes) specified by the octal, decimal, or BCD value.

After each instruction is executed, CPU breakpoint and status is output to the current CRT IS device. When execution halts, the CPU status and memory contents are output as at a breakpoint. Because of the internal coding of the system monitor, the address of BKPl appears to increase as each instruction is traced and status is output. However, when trace execution halts, both breakpoints are reset to their original addresses (when the TRACE command was executed).

To halt execution during TRACE, press any key. Repeatedly pressing a key may be necessary to halt TRACE.

Section 5: System Monitor

Example: TRACE 10 output

BKP1 BKP2 PTR1 PC DR FIR DV CY NG LZ ZR RZ OD DC E PTR2 ROM 022274 46 36 0 0 1 0 0 1 0 0 0 0 022273 114303 0377713 0377732 000 DR AR DV CY NG LZ ZR RZ OD DC Ε BKP1 BKP2 PTR1 PTR2 ROM 022275 46 36 0 0 0 0 01 022274 114303 0377713 0377732 000 PTR1 PTR2 DR AR DV CY NG LZ ZR RZ OD DC Ε BKP1 BKP2 ROM 01 022275 114303 0377713 0377732 000 021636 46 36 0 0 0 0 0 DR AR DV CY NG LZ ZR RZ OD DC BKP1 BKP2 PTR1 PTR2 ROM 021637 46 36 0 0 0 0 01 021636 114303 0377713 0377732 000 PTR2 DR AR DV CY NG LZ ZR RZ OD DC Ε BKP1 BKP2 PTR1 ROM 021640 46 36 0 0 0 0 0 0 0 01 021637 114303 0377713 0377732 000 DR BR DV CY NG LZ ZR RZ OD DC £ BKP1 BKP2 PTR1 ROM 021641 20 36 0 0 0 0 0 0 0 01 021640 114303 0377713 0377732 000 DR AR DV CY NG LZ ZR RZ OD DC BKP1 BKP2 PTR1 Ε PTR2 ROM 021642 20 10 0 0 0 0 0 0 0 0 01 021641 114303 0377713 0377732 000 PC DR AR DV CY NG LZ ZR RZ DD DC E BKP1 BKP2 PTR2 PTR1 ROM 021643 20 10 0 0 0 0 01 114333 114303 0377713 0377732 000 MEM 0:0 ROO 000 012 265 230 243 043 155 204 026 000 011 210 303 030 011 210 R10 242 200 350 212 371 000 001 000 153 031 305 031 266 031 247 031 K E E 342 207 022 210 022 210 106 251 R20 040 044 233 230 140 011 236 200 R30 237 200 034 000 075 210 316 230 070 204 230 136 262 001 377 251 R40 110 233 230 001 000 000 030 056 340 040 262 030 377 321 000 140 R50 000 051 000 000 000 000 000 000 366 012 262 065 210 261 014 140 v 25 1 1 R60 000 000 000 000 000 000 357 012 036 306 000 000 316 274 011 316 F NX N R70 016 316 000 000 000 000 000 000 030 030 230 316 306 207 117 220 NF 0

PTR1= octal value
Changes pointer address.

PTR2= octal value
Changes pointer address.

#### WRITING BINARY PROGRAMS

# 6.1 Program Structure

An assembly language program is required to have a table of five pointers, or addresses, to tell the system where important parts of the program are. The system will use these pointers to find the table of keywords which the binary program implements and the associated routines to execute each of those keywords. This structure called the program shell is shown on the next page.

```
NAM
       DEF RUNTIM
       DEF ASCIIS
DEF PARSE
       DEF ERMSG
       DEF INIT
PARSE BYT Ø, Ø
      --Parse routine addresses go here.
RUNTIM BYT Ø, Ø
       --Runtime routine addresses go here.
       BYT 377, 377
ASCIIS BSZ Ø
       --Keyword table goes here.
       BYT 377
ERMSG BSZ Ø
       --Error message table goes here.
      BYT 377
INIT BSZ Ø
       --Initialization code goes here.
```

-- The rest of the binary program goes here.

RTN

FIN

The shell consists of the following parts:

- 1. The program control block.
- 2. Label definitions describing the locations of the tables that will allow the system to hook into the binary program. The following addresses must be included in this order:
  - 1. Run time routine table.
  - 2. ASCII keyword table.
  - 3. Parse routine table.
  - 4. Error message table.
  - 5. Initialization routine address.
- The actual tables that have been defined previously. They must contain the addresses of the routines that will be performed.
  - The parsing routines will tell the system how to check a keyword for the proper syntax and parameters, and how to convert it to the internal RPN token format.
  - The actual translation of the keywords into machine operations is done by the run time routines whose addresses are defined in the run time table.
  - A marker, two bytes containing 377, must be set directly after the run time and parsing routine tables. When a binary program is loaded, this marker tells the system to assign an absolute address to all routines. All other addresses (routine references) are relative to the beginning of the program.
  - To let the system know which character strings will be the keywords, an ASCII table must be created to specify the keywords.
  - An error message table allows assembly language programs to specify custom error messages.
  - The code for a special initialization routine that is to be executed during initialization of the system, as at power-on, reset, allocation, and deallocation times. Refer to Initialization Hooks in section 3.
- The routines that will actually do the operations required for defining and executing the new BASIC keyword must come after the tables.

The system will use the structure of the program shell to access the routines in the program. If a mistake is made in the structure, then the system cannot run the program.

The labels that are used to reference routines and routine tables can be any name as long as the names of routines in the tables correspond with the names of the routines themselves.

In addition, after the execution of a routine, control must be passed back to the system by executing a return. A return may be included after every routine.

#### Control Block

The program control block is 40 (octal) bytes long and is required to tell the system important things including:

- The first four characters in the name of the binary program.
- · The length of the program in bytes, including the control block.
- The type of file is contained in the seventh byte. The format of the bits in this byte are as follows:

Bit	Meaning
Ø	000=BASIC Main Program
1	001=BASIC Subprogram
2	002=Binary Program
3	Undefined
4	Undefined
1.00	Undefined
5	Undefined
6	Ø=Option base 1
	1=Option base 0
7	Ø=No COMMON
	1=COMMON

- · The binary program number.
- The name of the file in mass storage (up to 10 characters).
- · Six bytes required by the system.
- The base address of the first byte of the control block.

The control block is generated by the NAM instruction, which specifies the program name and number.

The program listed below is used in examples throughout this section.

1000		-NAM	167,TEST RUNTIME KEYWORDS PARSING
1010		DEF	RUNTIME
1020		DEF	KEYWORDS
1030		DEF	PARSING
1040		DEF	ERMSG
1050		DEF	INIT
1060	RUNTIME	BYT	0,0
1070		DEF	TEST.
1080	PARSING	BYT	0,0
1090		DEF	TESTPARS
1100	ERMSG	BYT	377,377
1110	KEYWORDS	ASP	"TEST"
1120		BYT	377
1130	INIT	RTN	
1140	****	1 1000	R56, #0,371
1150		LDM	R55, =0,371 R55, =PTR2- R55, =PTR2- =SCAN
1160		STMI	R55, =PTR2-
1170		JSB	=SCAN
1180		RTN	
1190		BALL	241
1200	TEST.	JSB	*STBEEP
1210	SCAN	RTN	
1220	SCAN	DAD	21110
1230	PTR2-	DAD	177715
1240	STBEEP	DAD	10441
1250		FIN	

Example: The program TEST is 107 (octal) bytes long and contains the following NAM statement and control block.

1000

NAM 167, TEST

О	EST	5 [LENGTH]	7 TYPE	<sup>6</sup> BPGM
10	[NAME OF FIL	E AS ON DISC		
DRIVE]	[BASE ADDRESS IF ABS]	24 [UNI	DEFINED]	
[LAST BYTE ADDRESS]	DEF RUNTIME	DEF ASCIIS		EF RSE

## Memory Contents

# ASCII Representation

124	105	123	123	107	000	002	167	TESTG
124	105	123	124	102	040	040	040	TESTB
040	040	000	000	000	000	000	000	
064	221							

The first four bytes contain the ASCII representation for the name TEST. The next two bytes, with the least significant byte first, contain the length of the binary program in bytes. The type of file that the program is stored under is represented in eight bits (one byte), and the binary program number is stored in the last byte. The next 10 (decimal) bytes show the ASCII representation for the file name under which the program is stored, with ASCII blanks (040) to fill the rest of the bytes. The following six bytes are undefined, and the last two bytes contain the address of the first byte in the binary program.

## System Table

The system uses this table to locate the routines and tables it will need to interpret the binary program. The system table must always be present in a binary program and must always define the subsequent tables in the proper order. During operations the system will need to have the address of a routine to handle parsing, initialization, execution, or error conditions. It will expect the address to be at the proper location as shown below:

Bytes From Base Address	Sample System Table
32	DEF RUNTIME
34	DEF KEYWORDS
36	DEF PARSING
40	DEF ERMSG
42	DEF INIT

When the system looks for a run time routine, it will add 32 (octal) to the base address of the program and access the run time routine table at run time. Likewise, it will add 34 (octal) to the base address to find the parse routine table, and so on. The system will expect the tables and the initialization routine to be in exactly these places in the program.

# Placement of Binary Program Routine Tables

The addresses in the parse and run time routine table will be made absolute by the system when the program is executed. To indicate the end of the tables whose addresses will be absolute, the system looks for two bytes of 377's. Only the parsing routine and run time routine tables are required to have absolute addresses, so all other routine tables must follow the two bytes of 377's.

## ASCII Keyword Table

The system will check a binary program for a BASIC keyword before it will try to process the keyword. In the ASCII table, all of the key words are arranged sequentially. When a BASIC statement is entered into the CRT, the system attempts to match the characters with a keyword in the table. The order of the keywords will affect the parsing and execution of the keyword, as the first keyword in the table will be processed by the first parsing routine in the parsing routine table and executed by the first run time routine in the run time routine table.

The system attempts to find a match by comparing each character in the table with each character in the keyword until it reaches a character with the most significant bit set. This indicates the end of a keyword, and, if no match has been found, the system assumes that the next character begins a new keyword and increments the number of the token. The search stops when a match has been found or a byte containing 377 is found.

Example: The following code creates an ASCII keyword table with one keyword, TEST. The ASP instruction creates an ASCII string with the most significant bit set on the last character, and the BYT 377 instruction signifies the end of the ASCII keyword table.

1110 KEYWORDS ASP "TEST" 1120 BYT 377

# Parsing Routine Table

If the system accesses a BASIC statement keyword that a binary program has listed in the ASCII keyword table, it will use the parsing routine provided in the program. Functions will be parsed by the system. The position of the keyword in the ASCII keyword table determines which parsing routine will need to be executed. If the keyword does not need to be parsed, then the corresponding position in the parsing routine table must be filled with two bytes of 0's to reserve the space corresponding to the ASCII keyword table.

The system will always skip the first two bytes after the location of the parsing routine table. The next two bytes are used as the address of the first parsing routine.

Example: The following parsing routine table has only one routine, TESTPARS, to parse the keyword TEST.

1080 PARSING BYT 0,0 Two dummy bytes 1090 DEF TESTPARS First parsing routine

### Run Time Routine Table

Each keyword also has a run time routine associated with it. More than one run time routine may be listed in the table, so the system distinguishes between them in the same way as in the parsing routine table. When the system encounters a keyword that is listed in a binary program, it passes control to the proper routine corresponding to the position of the ASCII keyword in the keyword table.

Example: The run time routine table for the program TEST contains one routine address "TEST." which corresponds to the keyword "TEST" and the parsing routine TESTPARS.

1060 RUNTIME BYT 0,0 Two dummy bytes 1070 DEF TEST. First run time routine

Error Message Table

When an error is flagged in XCOM, the executive loop calls an error reporting routine that displays the error message. If the error number is less than 128 (200 octal), then it is a ROM error message and the bank-addressed ROM is selected whose number is in RAM location ERRROM. If the error message number is greater than 128 (200 octal), then the message will be from the binary program whose number is in ERRBP#.

The error message table is similar to the ASCII keyword table. It is constructed of entries which are strings of ASCII characters, the last character of each string having the most significant bit set. The table is terminated by a BYT 377.

Error messages in ROMs are numbered 0 through 177 (octal). Binary programs are numbered from 377 down to 200 (octal). The first nine error messages for ROMs and binary programs are for default errors. They will give only warning messages if defaults are on (refer to the owner's manual). The other error messages will always display the appropriate error message. Example error message table:

```
10 ERMSG BYT 200,200,200,200 | ININE DUMMY BYTES (377,367),
20 BYT 200,200,200,200,200 | I WITH THE MSB SET
30 ASP "SYNTAX-CHECK KEYWORD." !!ERROR 366 OCTAL
40 ASP "ROW OUT OF RANGE, >16." !!ERROR 365
50 ASP "COL OUT OF RANGE, >32." !!ERROR 364
60 BYT 377
70 INIT BSZ 0
80 RTN
```

Initialization Routine

The program TEST has no need to initialize pointers, hooks, or flags. Therefore, the INIT routine returns control to the system. For times when further initialization is needed, refer to paragraph 3.5.

1130 INIT RTN

External Address Table

If any of the system locations have been used in the program, a table must be included to define the labels as absolute addresses.

Example: In the program TEST the addresses SCAN, PTR2-, and STBEEP are used and must be defined for the system.

1220 SCAN DAD 21110 1230 PTR2- DAD 177715 1240 STBEEP DAD 10441

## 6.2 Attributes

Attributes define the type of a token. The system uses the attribute type to determine how parsing is to occur, how allocation and deallocation are to be performed, and how decompiling is to be done. The system is told how the keyword is to be handled at these times. The attributes must be defined immediately before the run time code in the program memory as shown in line 90:

1190	BYT 241
1200 TEST	JSB=STBEEP
1210	RTN

There are two types of attributes: primary and secondary. All keywords have primary attributes, but only functions have secondary attributes. The secondary attributes tell how many and the type of parameters the function will need and may occupy one or more bytes.

## Attribute Location

The attributes must be placed directly before the run time routine code. The primary attributes must be the first byte before the run time routine. The secondary attributes would precede from the first byte to the last byte. The system checks the attributes from the bottom up, starting with the primary attribute and ending with the last parameter.

The following program listing:

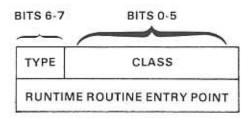
040 055 BYT 040,055

is the octal representation of these attributes:

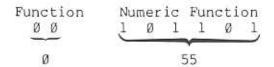
055 Primary attribute - numeric function
040 Secondary attributes - two numeric parameters

# Primary Attributes

The primary attribute consists of one byte of information containing the type of the keyword in the two most significant bits and the class of the keyword in the next six bits as shown:



The user must define the primary bits in order to tell the system exactly how he wants the system to recognize the keyword. For instance, if the keyword is to be a numeric function its attributes would be:



### Type

Bits 7 and 6 define how the keyword may be used. A keyword may be a BASIC statement or another command for calculator mode. A BASIC statement may be defined as legal after a THEN or illegal after a THEN. System commands are BASIC statements used only in calculator mode. Functions may be used in BASIC programs or in calculator mode.

The codes for each of the types are the following:

Bits 7,6	Octal	Туре
1 1	3	BASIC statement, illegal after THEN
1 Ø	2	BASIC statement, legal after THEN
0 1	1	System commands (nonprogrammable)
ØØ	Ø	Functions and others

## Class

The class will give the system further information on how to process the keyword. The class should follow directly after the type of keyword. For example, a function that returns a number will be in the numeric function class. Keywords that are to be invisible when the program is decompiled have their own class. All BASIC statements that are not functions and all system commands are reserved words.

Example: The keyword "INPUT" uses two tokens to compile but only one shows when it decompiles. The first token puts the system into a pseudo-calculator mode to allow characters to be entered from the keyboard and outputs a "?." The other, which is hidden during decompiling, takes the system out of the pseudo-calculator mode and does the actual storing of the input values. The class of the second keyword keeps it from being printed in the program listing. The keyword "LOAD" is in the class of reserved words. Refer to the sample program LINPUT in section 7.

	40000	Service Service	
USA	1111	C. 1.	asses

	Bi	ts	5-	-Ø		Octal	Class
1	Ø	Ø	Ø	Ø	1	41	Reserved words
1	Ø	Ø	1	Ø	Ø	44	Invisible at decompile time
1	Ø	1	1	Ø	1	55	Numeric function (such as, SIN, IP)
1	Ø	1	1	1	Ø	56	String function (such as, CHR\$, VAL\$

## Secondary Attributes

At parse time, if the system parser finds a match for a keyword in the binary program ASCII table, it will then check the attribute type. If the keyword is a statement, control passes to the binary program parse routine. If the keyword is a function, the secondary attributes determine the type and number of parameters to use.

One byte is needed if the function uses one or two parameters, and a second byte is needed if there are more than two but less than seven parameters. More parameters require more bytes. The first four bits of the first byte indicate the number of parameters that the function will accept. The next two bits define the type of the first parameter, and the second parameter is defined by the last two bits. Thereafter the consecutive pairs of bits define extra parameters.

# Parameter Types

Туре	Description
ØØ	Numeric
Øl	Numeric array
1 Ø	String
1 1	Strange

<u>0010</u> <u>00</u> <u>00</u> 2 Parameters First Second

#### 6.3 Assembler Instructions

The instruction set is used to communicate between the assembly language programmer and the CPU. Assembly language instructions can move data, perform arithmetic operations, and execute other functions. There are two types of instructions: those which operate directly on the CPU and are translated into machine language; and pseudo-instructions which act as messages to the Assembler ROM.

The typical instruction is broken up into five fields. The first field is the line number, for the convenience of the programmer. When assembled, the program will not have line numbers, instead it will show the value of the instruction counter. The instruction counter is the offset in bytes from the start of the program. The next field is an optional label field. System labels may be defined in the global file. Other labels must be defined in the routine. The opcode comes after the label field and is the heart of the instruction because it tells the CPU or Assembler ROM what is to be done. Following the opcode is the operand(s) for the instruction, and at the end of the instruction or in the label field the programmer may place a "!" followed by a comment.

In assembler mode the system will automatically space the elements typed in the proper fields. The programmer has only to distinguish the fields by at least one space. The registers may be referred to by their octal numbers, and the system will add the "R" in its proper place.

Example: Line number 120 may be typed in as follows:

120 LDMD 46,22 IR MULTI-BYTE LOAD

After pressing [END LINE], it will appear in the program listing as:

120

LDMD R46,R22

'IA MULTI-BYTE LOAD

## Line Numbering

Each line of a program source code must begin with a line number (which will not appear in the assembled code). A line number may be up to 99999 and may be entered individually or automatically, using [AUTO] for automatic line numbering. When a program is assembled the line numbers will appear as relative addresses of the instructions, that is, the instruction location counter.

#### Labels

A label may be from one to eight characters long. The label field starts in the second space after the line number. A digit may not be used as a first character, and no spaces may be used in a label because a space denotes the end of a label. When variable storage is needed in the program, a label may be used after the run time routine. To simulate control loops and branch execution, a label may be used to designate the location of the jump.

# Opcodes

The opcodes for assembly language instructions may be entered after typing at least two spaces after the line number or at least a single space after a label. Entries in the opcode field are restricted to valid instructions. Blanks are not allowed within the opcode field.

Opcodes may be single-byte, multi-byte, or pseudo-operations. The pseudo-operations may act upon bytes but are only messages to the Assembler ROM and do not generate executable code.

## Operands and Addressing

Depending upon the kind of instruction to be performed, the operand may be a register, a label or address, a pointer to a value, or a relative location which must be offset by an absolute address. The DRP will point to the register that will be operated upon according to the opcode. If the opcode calls for direct addressing, where the value is at a location outside of the CPU register bank, the operand will contain the address of the value in memory. If the opcode calls for direct addressing, where the value is pointed to by a label that is located outside of the CPU register bank, the operand will contain the address of the pointer in memory.

Indexed addressing can be used to access an area of memory by adding a base address to an offset such as in table searching. The absolute address of a label can be obtained by adding BINTAB to the relative address of the label.

#### Comments

A comment must begin with an exclamation point "!." A comment may be typed beginning in the first or second space after the line number or one or more spaces after the other elements of the instruction. Comments may be as long as needed, though the limit is 160 characters per line.

#### Constants

Constants may be entered in octal, BCD, or decimal notation. A BCD value is entered by immediately following the value with a "C," while a decimal value is followed by a "D"; otherwise the system assumes octal values. Constants will be stored as one or more bytes, depending on whether it indicates a single- or multi-byte operation. After the program is assembled constant values are placed immediately after the machine code.

# Syntax and Explanation

Each of the opcodes are discussed in detail in the next three subsections. The opcode is shown above its explanation, then following the explanation is an example of how the instruction may be used.

The first two letters of the opcode signify its operation, but the designation for a single-byte, "B," or multi-byte, "M," operation must be added at the end. In addition, if a type of addressing other than register immediate is needed, then the letter for that addressing mode must be added, "D" for direct or "I" for indirect. Instructions using direct or indirect addressing will have opcodes of four characters. A register being used for indexing must be entered in the operand field with an "X" instead of an "R." Pseudo-instructions always have opcodes of three characters.

The examples are designed to give the programmer a few hints for using the instructions in binary programs and clarify some points about the syntax of the instruction set.

## Syntax Guidelines

LDB Instructions shown in capital letters must be entered exactly as shown (in either upper- or lower-case).

- [] Items shown between brackets are optional. If several items are stacked between brackets, any one or none of the items may be specified.
- ... Three dots (ellipsis) following a set of brackets indicate that the items between the brackets may be repeated.
- ( ) Contents of.
- Complement.
- B/M Single- or multi-byte instruction.
- AR Address register location. Location of first byte addressed by the ARP. Can be a register, R\*, or R#.
- DR Data register location. Location of first byte addressed by the DRP. Can be a register, R\*, or R#.
- A Address mode for load/store. Can be blank (for immediate), D (for direct), or I (for indirect).
- ARP Address Register Pointer. A 6-bit register used to point to one of 64 CPU registers. The byte to which ARP points is often used as the first of two consecutive bytes forming a memory address.
- DRP Data Register Pointer. A 6-bit register used to point to one of 64 CPU registers. The location to which DRP points is often used as the destination for data loaded into the CPU.
- R(x) CPU register addressed by (x).
- M(x) Memory location addressed by (x) which must be 16-bit address.
- PC Program Counter. CPU registers R4 and R5. Used to address the instruction being executed.

- SP Subroutine Stack Pointer. CPU registers R6 and R7. Used to point to the next available location on the subroutine return address stack.
- EA Effective Address. The location from which data is read for load-type instructions or the location where data is placed for store-type instructions.
- ADR Address. The two-byte quantity directly following an instruction that uses the literal direct, literal indirect, index direct, or index indirect addressing mode. This quantity is always an address.

## LOAD/STORE Instructions

The instructions for loading and storing data have access to all eight addressing modes, and they can be single- or multi-byte.

## LD CPU Instruction

Data register is loaded with the contents of the effective address determined by the operand and the addressing mode.

Format: LDBA DR, operand single-byte LDMA DR, operand multi-byte

ST CPU Instruction

Contents of data register are stored in effective address determined by the operand and the addressing mode.

Format: STBA DR, operand single-byte STMA DR, operand multi-byte

# Addressing Modes

The CPU allows several addressing modes. These include literal, register, indexed, and stack modes of memory access.

Not all addressing modes are available to all instructions. The load (LD) and store (ST) instruction have access to all addressing modes except stack addressing, and they are used here for illustration. For a list of the addressing modes used by a particular instruction, refer to appendix B.

Most addresses are referred to as two-byte quantities. Because addresses are two consecutive bytes, only the least significant byte is referenced. For instance, the address register (AR) is actually a single byte within the CPU register bank that is pointed to by the address register pointer (ARP). When the address register contains an address, the CPU register pointed to contains the least significant byte of the address. The next register (ARP + 1) contains the most significant byte of the address.

The multi-byte feature of the CPU allows data to be manipulated in quantities of one to eight bytes. Therefore, in the following descriptions, only the address of the first byte is specified.

In the following descriptions, the effective address (EA) points to the first byte of data to be affected by the instruction.

Register Mode

This mode allows the CPU registers to contain addresses as well as data. There are three types of register addressing: register immediate, register direct, and register indirect.

Register Immediate

Examples:

LDB R36, R32 Loads contents of R32 into R36.

STM R40,R50 Stores the contents of R40-R47 into

R5Ø-R57.

Register Direct

Examples:

LDBD R36,R32 Loads CPU register R36 with the contents

of the system memory location addressed

by R32-R33.

STMD R40,R50 Stores contents of R40-R47 into system

memory beginning with the location

addressed by R50-R51.

# Section 6: Writing Binary Programs

# Register Indirect

## Examples:

LDBI R36,R32 If R32 contains 105731, and location

105731 contains 110437, the contents of 110437 is loaded at location R36.

STBI R36, R32 If R32 contains 105371, and 105731

contains 110437, then the contents of R36 is stored at location 110437.

#### Literal Mode

The operand is a literal quantity stored in memory immediately following the opcode. A literal string, ten octal bytes or less, is a BCD, octal, or decimal constant or a label. The programmer is responsible for ensuring that the number of bytes of the literal string matches the DRP setting. The assembler does not check for a mismatch. Literal mode includes literal immediate, literal direct, and literal indirect forms of addressing.

## Literal Immediate

### Examples:

LDB R36,=10D Loads 10 decimal (12 octal) into CPU

register R36.

LDM R40, =0,0,0,0,0,0,0,120 Loads 120 octal (a floating point 5

in BCD format) into register R40-R47.

LDM R32,=LABEL Loads R32-R33 with the relative address

of LABEL.

## Literal Direct

## Examples:

LDBD R34,=ROMFL Loads the contents of the memory

location addressed by the label ROMFL

into CPU register R34.

STMD R74, =CHIDLE Stores the contents of CPU register R74

through R77 into four memory locations beginning with the location addressed

by the label CHIDLE.

Literal Indirect

Example:

STBI R30,=ADDR

Stores the contents of CPU register R30 into the memory location addressed by another memory location which is itself addressed by the two-byte literal quantity specified by the label ADDR.

Index Mode

Indexing is useful for accessing data when the data is stored in a table. In indexed addressing, a fixed base address is added to an offset to create the desired address. The CPU performs this addition using CPU registers R2 and R3. After an index instruction, these registers contain the effective address (the sum of the base and the offset). Neither the original base nor the offset is altered in memory. There are two types of indexed addressing: index direct and index indirect.

Index Direct

Example:

LDBD R36, X30, TABLE

Loads into CPU register R36 the contents of the memory location addressed by registers R2 and R3. R2 and R3 contain the sum of the contents of R30-R31 and the address TABLE.

Index Indirect

Example:

STMI R36, X30, OFFST

Stores the contents of CPU register R36 and R37 in memory, beginning with the location addressed by another memory location which is addressed by CPU registers R2 and R3. Registers R2 and R3 contain the sum of the address in R30-R31 plus the offset specified by the

label OFFST.

STBI R36, X34, 66

Stores the contents of R36 in the location addressed by R2 and R3 (sum of

the address in R30-R31 plus 66).

### Stack Instructions

In stack addressing, a register pair serves as a pointer to the stack in memory. A load or store is performed at the top of the stack, and the register pair is decremented or incremented to the new top of the stack. Instructions push and pop are available to push data onto and pop data from stacks in the main memory. These stacks can be addressed using direct or indirect addressing.

PU

Pushes single byte or multi-byte using direct or indirect addressing. The stack pointer is incremented (increasing stack) or decremented (decreasing stack).

## Examples:

PUBD R32,+R12

Pushes single byte from R32 onto the R12

stack. The stack pointer is

incremented.

PUBI R32,-R46

The stack pointer is first decremented and then the single byte contained in R32 is pushed onto the R46 stack.

PO

Pops single byte or multi-byte off stack using direct or indirect addressing. The stack pointer is incremented (increasing stack) or decremented (decreasing stack).

POBD R32,+R20

Pops single byte contained in R12 onto the R20 stack. The stack pointer is incremented after the operation.

POBD R32,-R20

The stack pointer is first decremented and then R32 is loaded with the byte

pointed to by R20.

## Stack Addressing

You can address a stack from nearly any CPU register pair. Registers R6 and R7 are hardware-dedicated and always point to the subroutine return stack, a fixed stack of 512 bytes. A subroutine jump will automatically push an address onto this stack and a return will load the program counter with the address on the top of the stack, causing execution to begin at that address on the next cycle. The R6 stack is also affected by SAD and PAD instructions (save and restore status), which push three bytes onto the R6 stack and remove them respectively.

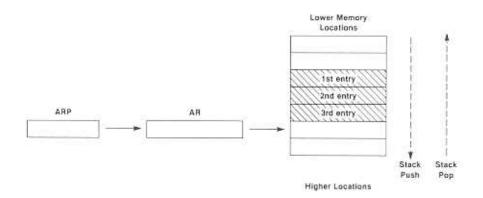
Another stack used by many of the system routines at run time is the R12-R13 operation stack. This stack is used to pass parameters between system routines. The documentation for each routine using this stack describes what the routine expects on the R12 stack and what it leaves after it has finished.

Stacks may be increasing or decreasing. An increasing stack is one which is filled in the direction of higher memory locations and from which data is removed in the direction of lower memory locations. In a decreasing stack, data is pushed in the direction of lower memory locations, and taken off in the direction of higher memory locations. To avoid confusion, it is best to address a particular stack using only instructions for an increasing stack or only instructions for a decreasing stack, but not both.

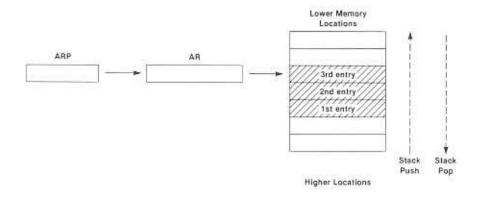
For stack addressing, the stack pointer is contained in the AR. Multiple stacks are handled by having multiple stack pointers within the CPU register space. A stack is activated by setting the ARP equal to the location of that stack pointer.

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For an increasing stack, the AR must point to the next location available on the stack. For a decreasing stack, the AR points to the occupied location on top of that stack.



Increasing Stack



Decreasing Stack

#### Stack Direct

In this addressing mode, the stack is presumed to contain data. Stores to the stack (pushes) fill the stack. Loads from the stack (pops) empty the stack.

For a push onto an increasing stack, the AR points to the location where data is to be stored. Following the store, the AR is incremented by the number of bytes stored. For a pop operation from an increasing stack, the AR is first decremented by the number of bytes to be popped off. The AR then points to the location of the data to be removed from the stack.

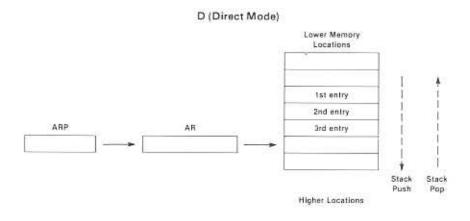
For a pop from a decreasing stack, the AR points to the location of the data to be removed. Following the removal, the AR is incremented by the number of bytes moved. For a push operation onto a decreasing stack, the AR is first decremented by the number of bytes to be stored on the stack. Then the data is pushed onto the stack.

#### Stack Indirect

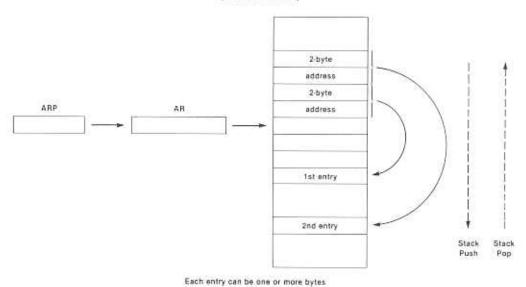
In this mode, the stack is presumed to contain an ordered list of addresses. These addresses point to the location from which data is read by pops or to the location into which data is stored by pushes.

For a push onto an increasing stack, the AR points to the effective address. After storing data in M(EA), the AR is incremented by two. For a pop instruction from an increasing stack, the AR is first decremented by two in order to point to the effective address. The effective address is then loaded into the CPU register designated by the DRP.

PUBD DR,+AR Push byte direct with increment.



## I (Indirect Mode)

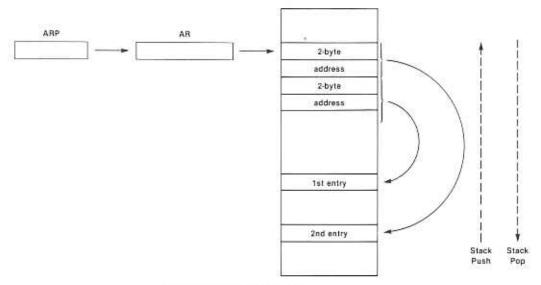


INCREASING STACK

The instructions available for use with an increasing stack are:

PUBD DR, +AF	Push byte direct with increment.
PUMD DR, +AF	
PUBI DR, +AF	
PUMI DR, +AF	Push multi-byte indirect with increment.
POBD DR, -AF	
POMD DR, -AF	Pop multi-byte direct with increment.
POBI DR, -AF	
POMI DR, -AF	

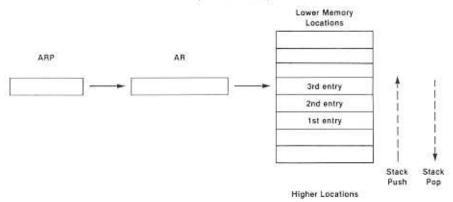
## I (Indirect Mode)



Each entry can be one or more bytes

## DECREASING STACK

## D (Direct Mode)



The instructions available for use with a decreasing stack are:

PUBD	DR,-AR	Push byte direct with decrement.
PUMD	DR,-AR	Push multi-byte direct with decrement.
PUBI	DR,-AR	Push byte indirect with decrement.
PUMI	DR,-AR	Push multi-byte indirect with decrement.
POBD	DR,-AR	Pop byte direct with increment.
POMD	DR,-AR	Pop multi-byte direct with decrement.
POBI	DR,-AR	Pop byte indirect with decrement.
POMI	DR,-AR	Pop multi-byte indirect with decrement.

Arithmetic and Logical Instructions

AD

Add may be used to combine the value of the data register and the contents of the operand. This operation may be performed on single bytes or multiple bytes, and direct addressing or constants may be used. In BCD mode addition will take place using four-bit digits. The result is always stored in the data register.

## Example:

ADB R20, R30

Adds the contents of R30 to R20 and

places the result in R20.

ADMD R20,=BINTAB

Takes the location of the beginning of the binary program and adds it to the

value in R20, R21. The result is

stored in R20, R21.

AN

Each bit in the data register is compared to the corresponding bit in the operand. If the bits being compared are both 1, then the result is a 1. If either bit is 0, then the result is 0. The operand may be a value in memory that is addressed directly. Although this instruction is available only for multi-byte operations, single-byte operations are possible with the DRP set to a boundary register.

# Example:

ANM R20, R30

Converts all of the 1's in R20-R21 to 0's if the same bits in R30-R31 are 0's.

If R20-R21 contain:

1 1 0 1 1 0 1 1

00101100

and R30-R31 contain:

01 000 111

00001011

then the result is:

01 000 011

00001000

CM

The compare is used to simulate the logical operations of a high level language. It is done by subtracting the operand from the data register and setting the appropriate status indicators; the result of the operation is not stored. In binary mode the subtraction is two's complement, and in BCD mode the subtraction is ten's complement. Compares may be either single- or multi-byte operations, and direct addressing may be used. When used previous to a logical jump, an IF-THEN BASIC statement may be simulated.

In order to simulate the relation:

DR <ar< th=""><th>CMM DR, AR</th><th>CY flag should be <math>\emptyset</math></th></ar<>	CMM DR, AR	CY flag should be $\emptyset$
	JNC LABEL	Jump if DR <ar< td=""></ar<>
DR>=AR	CMM DR, AR	CY flag should be 1
	JCY LABEL	Jump if DR>=AR
DR=AR	CMM DR, AR	ZR flag should be 1
	JZR LABEL	Jump if equal
DR#AR	CMM DR, AR	ZR flag should be Ø
	JNZ LABEL	Jump if not equal

The jump instructions JNZ, JZR, JCY, and JNC are explained later in this section.

OR

Each bit in the data register is compared to the corresponding bit in the operand. If either bit is a 1, then that bit in the data register is set to 1. Otherwise the bit in the data register is set to  $\emptyset$ . This logical operation may be performed on single bytes or multiple bytes, but must use register immediate addressing only.

## Example:

ORB R20, R30

Leaves a 1 in R20 if the corresponding bit in R30 is set.

If R20 contains: 0 0 1 0 1 1 0 0 and R30 contains: 0 0 0 0 1 0 1 1 1 then the result is: 0 0 1 0 1 1 1 1

SB

Subtraction is simulated by adding the complement of the operand to the data register. Ten's complement is used in BCD mode, and in binary mode two's complement is used. The result of the subtraction is stored in the DR. The operand may be addressed immediately or directly, and can be a single- or multi-byte instruction. The CY flag is set to 1 if the result is positive and cleared if the result is negative.

# Example:

SBM R20, R30

In binary mode, takes the two's complement of R30-R31 and adds that to R20-R21. The result is put in R20-R21.

If R20-R21 contain:

1 1 0 1 1 0 1 1

00101100

and R30-R31 contain:

01 000111

00001011

then the complement of R30-R31

10 111 000

1 1 1 1 0 1 0 1

is added to R20-R21.

The result is:

11 011 011

00101100

The operation is done in binary mode. Since registers are shown in octal, the previous example would look like this:

Before:

R2Ø

R21

R3Ø

R31

Ø54

333

Ø13

107

Two's complement

Result:

R20

R21

R30

R31

223

041

Ø13

107

# Section 6: Writing Binary Programs

# Example:

SBB R20,R30

In BCD more, takes the ten's complement of the two digits in R30 and adds that to the two digits in R20.

If R20 contains:	00101100
which in BCD are the decimal digits:	28
and R30 contains:	0 1 0 0 0 1 1 1
which in BCD is:	47
Then the ten's complement of R30:	53
is added to R20:	28
and the result in R20 is:	80

#### XR

In the "exclusive or" logical operation the bit that corresponds in the data register is set to 1 when the bits being compared are not the same. When both bits are 1 or both bits are 0, the bit in the data register is set to 0. The CY and OVF flags are cleared.

# Example:

XRM R20, R30

Compares the individual bits in R20-R21 and R30-R31. If they are not the same, sets that bit to 1 in the DR; otherwise it is set to 0.

If R20-R21 contain:

11 011 011

00100100

and R30-R31 contain:

01 000 111

00001011

The result in R20-R21 is:

10011100

00100111

#### Shift Instructions

All shift instructions can be done in BCD or binary mode. In BCD mode the shift will affect a BCD digit, or four bits, and in binary mode it will affect a binary digit, a bit. Shifts may also be single- or multi-byte operations, and the result of a shift will be determined by the nearest boundary in the direction of the shift. In single-byte shifts the boundary is actually the register being shifted, whereas in multi-byte operations the boundaries are those in the CPU register bank. In arithmetic and logical operations the boundaries are normally toward the higher-numbered registers. With shifts, the boundary may be to the left, higher-numbered registers, or the right, lower-numbered registers depending on whether you are shifting right or left.

Shifts are made into one of the shift registers: the E register or the CY flag. In BCD mode shifts are made into and out of the E register, and in binary mode shifts are made into and out of the CY flag.

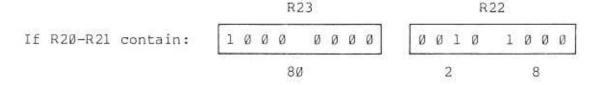
EL

The extended left shift will take the most signficant digit, put it into the shift register, move the rest of the contents one digit to the left and put the previous contents of the shift register into the least significant digit.

# Example:

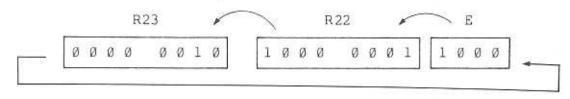
ELM R20

In BCD mode, shifts the most significant digit of R20-R21 (1000) into the E register. The other 12 bits will move left four bits, and the least significant digit will be filled with the previous contents of the E register (0001).



and "E" contained previously: 0 0 0 1

then the shift would take place as follows:



ER

The extended right shift moves the least significant digit to the shift register and the contents of the shift register into the most significant digit. It works in the same way as the extended left shift except that the movement is toward the right boundary.

# Example:

ERB R21

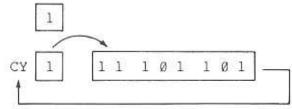
In binary mode, shifts the least significant bit (LSB) to the CY flag, then moves the previous contents of the CY flag to the MSB position.

Ø 1 1

If R2l contains: 1 1 0 1 1 and the CY flag is: 1

then the result would be:

LR

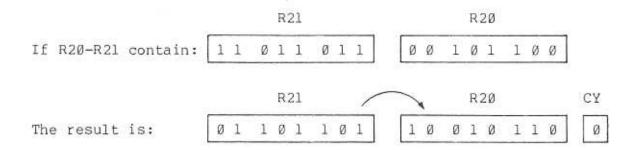


When a logical right shift is performed, the LSB is moved into the shift register and the MSB is cleared. The digit is maintained in the shift register and may be shifted back using the extended shift instructions.

## Example:

LRM R21

In binary mode, shifts the LSB into the CY flag and clears the MSB.



### LL

The logical left shift moves the most significant digit of the data register into the shift register and clears the least significant bit. If the shift causes a sign change then the OVF is set to 1.

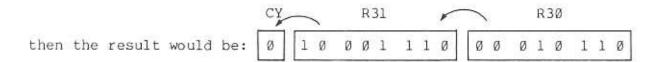
## Example:

LLM R30

In binary mode, shifts the MSB of R30 into the CY flag and clears the LSB of R31.

If R30-R31 contain:





DC

The decrement is simulated by adding the complement of 1 binary in binary mode, to the quantity in the data register. The quantity may be single or multiple bytes.

#### Example:

DCB R31

Subtracts one from the quantity in R31.

R31

If R31 contains: 1 0 0 0 0 0 0

R31 then the result is: Ø 1 1 1 1 1 1

The OVF flag is set to 1, because the sign changed.

## IC

When an increment is performed, 1 is added to the quantity in the data register. In BCD mode, the quantity is incremented by decimal 1, and in binary mode, it is incremented by a binary 1. In BCD mode the OVF flag is cleared (single- or multi-byte).

# Example:

ICM R20		, the decimal quantity i increased by 1.
	R21	R20
If R20-R21 contain:	1001 1001	0010 0101
which in BCD is:	9 9	2 5
	R21	R2Ø
then the result is:	1001 1001	0010 0110
which in BCD is:	9 9	2 6

#### NC

This complement instruction will give the nine's complement in BCD mode and the one's complement in binary mode. The nine's and one's complement are performed by taking the number of digits to be complemented and subtracting each digit individually from 9 in BCD mode and 1 in binary mode. The result is placed in the data register.

## Example:

NCB R20

In binary mode, flips all bits (one's complement operation).

R2Ø

If R20 contains:

Ø 1 Ø Ø Ø 1 1 1

R20

then the result is:

1 0 1 1 1 0 0 0

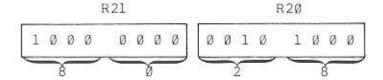
# Example:

NCM R20

In BCD mode, takes the nine's complement of the contents in R20-R21 by subtracting each digit from a BCD 9 (1001).

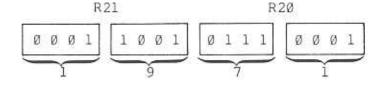
If R20-R21 contains:

which in BCD is:



then the result would be:

which in BCD is:



TC

The contents of the data register is replaced by the two's complement in binary mode or the ten's complement in BCD mode. Two's and ten's complement is found by incrementing the one's or nine's complement. In BCD mode, the OVF flag is cleared.

## Example:

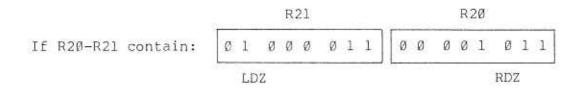
TS

The status of the contents of the data register are tested, and the appropriate status indicators are set. The OVF and CY flags are cleared in all cases, and the E register is not affected. This instruction is a single- or multi-byte instruction. The status indicators are discussed in section 2.

## Example:

TSM R20

Will set the status indicators and clear the OVF and CY flags.



the resulting flags will be set:

DCM	May be 1 (BCD) or 0 (binary).
E	- Not affected.
CY	Cleared.
OVF	Cleared.
OD:	Set to 1.
NG	Set to 0.
Z	Set to 0 (since quantity is nonzero)
LDZ	Set to 0.
RDZ	Set to Ø.

CL

The clear instruction permits the clearing of any byte or of any multi-byte portion of the CPU register bank. The DR is set equal to  $\emptyset$  and the flags CY and OVF are cleared.

## Example:

CLB R47

Clears R47

JSB

When a subroutine jump is made, the control of the program is given to a set of instructions with the intention of returning to the program at the next instruction after the jump was made. In order to return, the program counter for the next instruction must be stored. This return location is pushed onto the R6-R7 stack, and when the RTN instruction is executed, it is loaded back into the program counter. A subroutine jump that is made to a relocatable address in a binary program must be indexed from the absolute start of the program (BINTAB).

# Examples:

JSB =NUMVA+

Increments the program counter (PC) to the address of the next instruction after the JSB. That address is pushed onto the R6-R7 stack, and the PC is loaded with the address the jump is to be made to NUMVA+ (located at 22403). When the system executes a RTN, it pops the address of the next instruction off of the R6-R7 stack and loads that value into the PC.

JSB X14, ROUTINE

Makes a jump to ROUTINE by adding the value of ROUTINE as a label to the location of the start of the program (BINTAB) which is stored in R14-R15. In all other aspects it is the same as JSB=.

#### Jump Instructions

This group of instructions gives the capability for branching control to addresses that are defined by the label that the jump is being made to. If a condition is true, then the jump is made; otherwise, the jump is ignored and the next instruction is executed. These branching instructions use relative addressing. Labels that are used must be contained inside the program. The program counter (PC) is loaded with the value of the address, and program control moves to that location in the program memory. The maximum number of bytes that may be jumped is 177 octal (forward) higher-addressed bytes or 200 octal (backward) lower-addressed bytes.

Each conditional jump has a complement, except the jump on no overflow, which jumps on the opposite of the relation. For instance, the jump on negative is simply the opposite of the jump on positive and may be used in the same circumstances depending on the personal preference of the programmer. All of the jumps will be discussed.

JMP

The unconditional jump always occurs when executed. It is the only jump that does not check the status of any system flags.

Example:

JMP ALWAYS

Will always jump to ALWAYS, a location in the program.

JNO

Since the system has no jump on overflow, a jump on no overflow must be used for both cases. If the OVF flag is set to 1, then the jump is ignored and the next instruction will be executed. In the case of an overflow, the code after the jump instruction will perform the necessary steps, and then if necessary, continue the program.

Example: If a flag (E) is to be incremented when an overflow occurs:

ADM R20, R30

Executes the operation that may set

an overflow (OVF).

JNO RESUME

If there has been no overflow, the

program will begin at RESUME.

ICE

If JNO is ignored, then an overflow

has occurred, and the program

increments the E flag.

RESUME BIN

Resumes the program.

JPS, JNG

Jump on positive and jump on negative are made by checking the status of the most significant bit (NG) flag and taking the "exclusive or" of NG and the OVF. In the case of two positive numbers added together resulting in a negative number (NG=1), the jump on positive takes that into consideration and would jump because NG=1 and OVF=0 and the "exclusive or" would be 1 and the jump would be made.

Example: If R20 contains 073 and R30 contains 125 then the addition:

ADDITION ADB R20, R30

Adds 073 to 125 (octal) and sets NG=1 and

OVF=Ø.

JNG ADDITION

Since the exclusive OR of NG=1 and OVF=0 is 1 and JNG expects it to be 0, then the jump will not be made even though the NG flag says it is negative.

JOD, JEV

The least significant bit flag (OD) shows whether a number is even or odd. If the number is even, OD is set to  $\emptyset$  and JEV, jump on even, will take place. If the number is odd, OD=1, then JOD, jump on odd, will take place. This conditional jump works in binary and BCD modes.

Example: To find out if the 16-bit binary number stored in R36-R37 is a prime number, all even numbers may be ignored by the following code:

TSM R36 Checks to see if the number is even.

JEV NOTPRIME Since the number is odd, it might be prime.

JZR, JNZ

When making comparisons and when decrementing a counter, the jump on  $\emptyset$  and jump on not  $\emptyset$  are useful. If two quantities are equal, comparing them will cause the ZR flag to be set to 1. To simulate a conditional IF-THEN statement, a comparison is made prior to the jump. To simulate a controlled FOR-NEXT loop, the loop counter is decremented and the conditional jump made.

Example: To simulate IF X=80 THEN RESUME (R20 contains 120 octal which is 80 decimal):

CMB R20,=120 Compares R20 to 120. Since they are equal, the ZR flag is set to 1.

JZR RESUME Since ZR=1, the jump is made to the location RESUME.

To simulate the FOR-NEXT loop, the number of times that the loop will be executed is decremented and a check is made to see if that number of loops has been done.

Example: If R20 contains the number of times the loop is to be executed, then FOR X=1 to 20 would be:

DCB R20 After the statements have been executed, R20 is decremented. If R20 is equal to

zero, the ZR flag is set to 1.

JNZ LOOP If the loop has not been done the specified number of times, it must be done again starting at the beginning of

the statements (LOOP).

JCY, JNC

When the carry flag (CY) is set to 1, it indicates an addition has become too large for the register to handle. This happens often in subtraction and in comparisons. To simulate the statement IF-THEN with a "greater than or equal to" or "less than" relation, a compare is made between the values, and the CY flag is checked.

Example: If R20 and R30 contain the first and second numbers to be included in the compare, then the statement IF QUANTITY1 > QUANTITY2 THEN RESUME could be:

CMB R20, R30

Compares R20 to R30 by adding the negative of R30 to R20 and sets the status flags. If R20 is greater than or equal to R30 then CY=1. If R20 is less than R30, CY=0.

JCY RESUME

Jumps to the location RESUME if R20 is greater than R30 (CY=1).

JEZ, JEN

The jump on E equal to zero and the jump on E not equal to zero check the status of the E register for parsing routines and user defined flags. In parse routines the E flag will be set to 1 if the token searched for is found, and Ø if not found. After returning from a parse routine it is convenient to set an error message or to do another procedure if the token is not found. Also, if the E register is used as a programming flag, it may be set on a special condition to jump to a procedure.

Example: To demand a numeric parameter at parse time:

JSB=NUMVA+ JEZ ERR Try to parse a numeric value.

Jump if not found to error reporting.

JLZ, JLN

JLZ: Jump on left digit 0 (left BCD digit).

JLN: Jump on left digit not 0.

Example: If R20 contains 011, the following code would take the jump:

TSB R2Ø JLZ TRUE Section 6: Writing Binary Programs

JRZ, JRN

JRZ: Jump on right digit Ø (right BCD digit).

JRN: Jump on right digit nonzero.

Example: If 011 is in R20, the following code would not take the jump:

TSB R20 JRZ TRUE

ARP and DRP Load Instructions

These two instructions are available for loading the address register pointer or the data register pointer. They are not normally needed because the assembler automatically generates the ARPs and the DRPs where required.

ARP

Sets the address register pointer to the address register.

DRP

Sets data register pointer to the data register.

Use of R\*

When entering source code, you may substitute R\* for the AR or the DR in any CPU instruction. This causes the DRP or the DRP to be loaded with the least significant six bits of CPU register RØ. The effect is that the DR and the AR are specified by the contents of RØ. The CPU uses the DRP1 and ARP1 opcodes to implement this feature.

## Example:

LDB RØ, = 26

Loads RØ with 26.

LDB R\*,R30

Loads CPU register specified by RØ.

(which is now R26) with the contents

of R30.

STB R40,R\*

Stores contents of R40 into register

(R26 now) specified by RØ.

To avoid confusion, Rl is not allowed in either the DR or the AR fields. This means that CPU register Rl is inaccessible except by a multi-byte RØ operation or when RØ=l and the ARP or the DRP is specified by  $R^*$ .

Other Instructions

There are a few other CPU instructions which you can use.

BCD

Sets decimal mode (DCM=1). Arithmetic operations will be in BCD format.

BIN

Sets binary mode (DCM=0). Arithmetic operations will be in binary format.

CLE

The four bits of the extend register are set to 0.

DCE

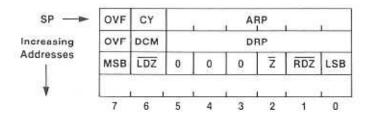
The extend register is decremented by 1. This instruction is always a binary operation, regardless of the DCM flag status.

PAD

Restores ARP, DRP, and status (usually after a SAD instruction) by popping them off the stack. The stack pointer is decremented by three, and all status flags except E are altered by the contents of the three stack locations that are read.

The first byte processed is read as the least LSB in bit  $\emptyset$ , the  $\overline{RDZ}$  bit 1,  $\overline{Z}$  in bit 2,  $\overline{LDZ}$  in bit 6, and MSB bit 7. The second byte is read as the DRP in bits  $\emptyset-5$ , DCM status in bit 6, and overflow flag in bit 7. The third byte is read as the ARP in bits  $\emptyset-5$ , carry flag in bit 6, and overflow flag in bit 7.

Following a PAD instruction, the stack has been read as shown here:



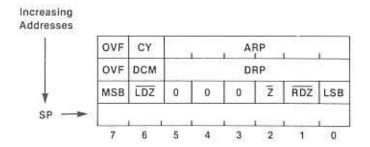
### RTN

The subroutine return stack pointer is decremented by two. Then the return address is read from the stack and written into the program counter.

#### SAD

Three bytes are pushed onto the stack to save the ARP, the DRP, and the status flags (except E). The first byte contains the ARP in bits  $\emptyset$ -5, CY in bit 6, and the overflow flag in bit 7. The second byte contains the DRP in bits  $\emptyset$ -5, DCM status in bit 6, and the overflow flag in bit  $\overline{7}$ . The third byte contains the LSB in bit  $\emptyset$ , RDZ in bit 1,  $\overline{Z}$  in bit 2, LDZ in bit 6, and the MSB in bit 7. The stack pointer is incremented by three. Status is not affected by this operation.

Following a SAD instruction, the stack contents are as shown here:



## 6.4 Assembly of CPU Instructions

When the address field of an instruction consists of a DR and an AR, each source statement is usually assembled into three bytes of machine code. These bytes are assembled in order as:

- 1. DRP: DRP set to point to DR.
- 2. ARP: ARP set to point to AR.
- Opcode: Perform operation.

A stack push instruction such as PUBD would be assembled and appear as shown here:

Byte Number	Machine Code	Source Code
000227	110 006 342	PUBD R10,-R6

When the address field of an instruction consists of a DR and a label, as in the case of literal direct and literal indirect addressing (such as, LDMI R32, =ADDRS), each source statement is usually assembled into four bytes of machine code:

- 1. DRP: DRP set to point to DR.
- 2. Opcode: Perform operation.
- 3. Low-order byte of literal quantity.
- High-order byte of literal quantity.

When the address field of an instruction consists of a DR, an AR, and a label, as in the case of indexed direct and indexed indirect addressing (such as, LDBI R36, X32, TABLE), five bytes of machine code may be generated for each source statement:

- 1. DRP: DRP set to point to DR.
- 2. ARP: ARP set to point to AR.
- 3. Opcode: Perform operation.
- 4. Low-order byte of address.
- 5. High-order byte of address.

The ARP and the DRP During Assembly

An optimizing feature of the Assembler ROM is the deletion of "unnecessary" ARP and DRP instructions during assembly.

If an instruction if not labeled (that is, there is not an entry in the label field) and the ARP (and/or DRP) is already set to the correct value, the previously set ARP (and/or DRP) is not generated during assembly.

## Example:

Byte Number	Machine Code	Source Code
000227	110 006 342	LABEL POBD R10,-R6
000232	342	POBD R10,-R6

In this example, both the ARP and the DRP are specified beginning with byte 227. Since they are now correctly set for the next instruction, they are automatically deleted when the second POBD R10,-R6 instruction is assembled. This results in the machine code shown in byte 232.

Not all previously set ARPs and DRPs are deleted during assembly. Times when they are not deleted include:

- Labeled instructions: Since a jump from anyplace in code may cause execution to resume at the label, the first ARP and DRP are not deleted after an instruction that contains an entry in the label field.
- Returns: After executing a subroutine jump, then returning, the first ARP and DRP encountered are not deleted.
- PAD: Following a PAD instruction, the first ARP and DRP are not deleted.

#### Pseudo-Instructions

Instructions to the assembler are pseudo-instructions. Each may be entered by typing a pseudo-opcode in the same field as the opcode for an instruction, followed by any additional operand.

Pseudo-instructions perform these three functions when encountered during assembly:

- Assembly control.
- · Data definition.
- Conditional assembly.

#### ABS base address

If the base address is less than 100000 (octal) then a ROM file will be generated at assembly time. Otherwise a binary program file will be generated and all labels are given absolute addresses, not relative addresses. The ABS statement must precede a NAM statement, if used.

### FIN

Signifies the end of the source code. This pseudo-instruction is required for assembly.

## GLO file name

If no file name is specified , GLO declares this source code to be a global file to be used in the assembling of the current source code. If there is a file name, it is the name of the global file to be used in the assembling of this source code. Comments are not allowed on the same line as the GLO instruction, and the instruction must precede ABS and NAM.

LNK file name

Will load another file containing more source code and continue assembling. Allows assembly of larger programs than would otherwise be possible.

LST

Causes the code to be listed on the current PRINTER IS device at assembly. The printed lines will be truncated at whatever the current line length is.

An address that is undefined when its label is encountered will be printed in object code as 326, 336, or 377, depending upon whether it is a DEF, a relative jump, or a GTO statement.

NAM binary program #, unquoted string

Sets up the program control block for a binary program. Should be preceded only by GLO, ABS, LST, UNL, DAD, EQU, or comments.

ORG address

Specifies a base address which is added to all following defined addresses (DADs). This pseudo-instruction is most useful in global files.

UNL

Turns off the list feature which was turned on by the LST pseudoinstruction. After an UNL, code is not listed during assembly.

Pseudo-Instructions for Data Definition

ASC numeric value, unquoted string ASC quoted string

Inserts into the object code the ASCII code for the number of characters specified of the unquoted string. Inserts the entire quoted string.

ASP numeric value, unquoted string ASP quoted string

Same as ASC except that the parity bit of the string's final character is set. (During operation, the system determines the end of an ASCII string in some system tables by checking to see if the character's parity bit is set. When the bit has been set, the system assumes the next character begins a new string or entry in the table.)

BSZ numeric value

Inserts into the object code the octal  $\$ number of bytes of  $\$ 0's specified by the numeric value.

BYT numeric value [,numeric value...]

Inserts literal values into the object code.

DAD Label DAD address

Assigns either an absolute address or a constant to a label. DAD and EQU are similar; DAD is usually used for addresses, while EQU is used for values other than addresses. ORG affects only DADs.

DEF label

Inserts the two-byte address associated with the label.

EQU Label EQU numeric value

Assigns either an absolute address or a constant to a label. This instruction is similar to the DAD pseudo-instruction.

GTO label

Generates four bytes of object code which load the program counter with the address, minus one, of the label. The label must be an absolute address.

The CPU relative jump instructions can cause jumps of from 177 (octal) to -200 (octal) memory locations. The GTO pseudo-instruction is useful for jumping beyond this range.

The GTO instruction is primarily for use in ROMs. It should not be used in a binary program unless that program has been declared an absolute program.

VAL label

Inserts the one-byte literal octal value associated with the label.

Pseudo-Instructions for Conditional Assembly

These instructions permit you to control assembly with conditional assembly flags. A conditional assembly flag has the same constraints as a label--it can be no more than eight characters in length, and the first character cannot be a digit.

A conditional assembly flag is treated the same as a label by the system. (For example, an assembly flag can be located by a label search.) For this reason, a conditional assembly flag should be unique, and should not duplicate a label.

AIF assembly flag name

Tests the specified conditional assembly flag and, if true, continues to assemble the following code. If the flag test false, the source code after the flag is treated as if it were a series of comments until an EIF (end of conditional assembly) instruction is found.

CLR flag name

Changes the specified conditional assembly flag to false.

EIF

Terminates any conditional assembly in process. Only one conditional assembly can be handled at a time. If a second one is encountered while the first is still active, the second will override the first.

SET flag name

Sets the specified conditional assembly flag to true.

# 6.5 Multiple Binary Programs

There can be up to five binary programs in memory at one time. There is a table of two-byte addresses called BINBAS that contains the base addresses in the order in which the binary programs were loaded. Bytes that are not used are zero. Anytime the system calls a binary program, it first fetches from BINBAS the base address for that program and stores it in BINTAB.

The ASCII keyword tables and the binary programs are searched in the order they are loaded. This is also how initialization routines are called.

#### SAMPLE BINARY PROGRAMS

#### 7.1 Introduction

This section includes five binary programs. In addition to being listed here, these programs are on the disc you received with your Assembler ROM. Source code file names end in "S", while object code file names end in "B."

Each of these programs is designed to illustrate assembly language programming, and each provides a function or keyword that is useful to the HP-87 operating system.

At the end of each program listing is a table of system routine addresses used by the program. Inserting the disc and placing a GLO GLOBAL pseudo-opcode near the beginning of the program eliminates the need for these addresses in some of the sample programs. Certain programs call system routines whose addresses are not available on the global file disc.

The string highlight program includes instruction on how to use a binary program following the listing.

## 7.2 String Highlight

Source Code: HGL\$S

Object Code: HGL\$B

```
1010 I* This binary program implements a string function called HGL$
1020 !* which accepts one string parameter and returns that string with
1030 1* the most significant bit of each character set.
1040 !* This binary program is a translation of the UDL$ binary program * 1050 !* from the HP-85 Assembler Rom manual.
1060 !*
1070 1*
                  (c) Hewlett-Packard Co. 1982
1080 1*
1090 !* Bn example of how this function might be used is:
1100 1*
1110 1* 100 INPUT A$
1120 1* 110 DISP HGL$(A$)
1130 !*
NAM 53, HGLBIN ! SET UP THE PROGRAM CONTROL BLOCK
DEF RUNTIM ! PTR TO RUNTIME ADDRESS TABLE
DEF ASCIIS ! PTR TO KEYWORD TABLE
DEF PARSE ! PTR TO PARSE ADDRESS TABLE
DEF ERMSG ! PTR TO ERROR MESSAGE TABLE
DEF TAIT ! PTR TO TAIT BOUTINE
1150
         DEF RUNTIM
DEF ASCIIS
DEF PARSE
DEF ERMSG
DEF INIT
1160
1170
1180
1190
1200
                               I PTR TO INIT ROUTINE
1240 RUNTIM BYT 0,0 ! DUMMY TOK# O RUNTIME ADDRESS
1250 DEF REV. ! TOK# 1 RUNTIME ADDRESS
1260 DEF HGL$. ! TOK# 2 RUNTIME ADDRESS
1270 BYT 377,377 ! TERMINATE RELOCATION
1290 ASCIIS ASP "HGL$B" ! KEYWORD #1
1300 ASP "HGL$" ! KEYWORD #2
1310 ERMSG BYT 377 ! TERMINATE ASCII TABLE & ERMSG TABLE
1330 INIT RTN
                               I NO INITIALIZATION TO BE DONE
    1350
1350 HGL$.
1370
1380
1390
1400
1410
1420
1430
1440
1450
1460
1470
1480
1490 MORE
1500
1510
1520
1530
1540
1550 DONE
```

# Section 7: Sample Binary Programs

1560		STMD R75, -PTR1-	AND RESTORE IT BEFORE RETURNING
1570		RTN	I DONE
1580	******	*************	**************
1590		BYT 0,56	I NO PARAMETERS, STRING FUNCTION
1600	REV.	BIN	I FOR ADDRESS MATH
1610		LDM R43,=40D,0	I LOAD THE LENGTH OF THE STRING
1620		DEF BATE-	! AND THE ADDRESS OF THE STRING
1630		BYT 0	! (MAKE IT A THREE BYTE ADDRESS)
1540		ADMD R45, - BINTAB	MAKE THE ADDRESS ABSOLUTE
1850		PUMD R43, +R12	I PUSH IT TO THE OPERATING STACK
1660		RTN	DONE
1670	DATE	ASC "40.102:veR 2891	.oC drakcaP-ttelweH )C("
1680	DATE -	BSZ O	I PLACE HOLDER FOR ADDRESS LOAD
1590	. * * * * * * * =	******	***********************
1700	BINTAB	DAD 104070	
1710	RSMEM-	DAD 31741	A secretarion of the second
1720	PTR1-	DAD 177711	DEFINE ADDRESSES
1730	PTR2-	DAD 177715	1
1740		FIN	I TERMINATE ASSEMBLY

1. In assembler mode load the source code:

ALOAD "HGL\$S" [END LINE]

2. To assemble the source code:

#### ASSEMBLE "HGL\$B" [END LINE]

If you want a printed copy of the object code as it assembles, you must designate a PRINTER IS device (that is, PRINTER IS 701). There must also be an LST instruction at the beginning of the code. The object code is now assembled and stored on your disc.

3. To use this function, return to BASIC mode. Type:

## BASIC [END LINE]

4. Load the object code. Type:

#### LOADBIN "HGL\$B" [END LINE]

Before running this program you may wish to set a breakpoint.With the system monitor inserted, type:

#### BKP REL (100)

The REL instruction sets the breakpoint at an absolute address in memory. The breakpoint information will appear on the CRT. It will also be printed if you specify a PRINTER IS device. For example:

#### BKP REL (100),701

The program will now halt at the address specified in the breakpoint. Your breakpoint will look similar to this when HGL\$ ("string") is typed:

DR AR OV CY NG LZ ZP RZ OD DC E BKP1 BKP2 PIR1 PTR2 114334 57 55 0 0 0 1 0 0 1 0 00 114333 000000 0370014 0370013 000 MEM 114233:0 ROO 005 001 242 053 334 230 100 204 110 107 114 102 254 000 002 053 HGLB, 110 107 114 044 102 040 040 040 307 200 350 212 075 210 001 001 HGL \$ B 233 230 053 016 010 013 310 200 041 000 324 230 267 230 233 230 040 040 000 000 000 000 000 000 R20 233 230 277 230 307 230 277 230 P O 015 000 000 000 000 060 360 001 320 230 321 230 233 230 017 231 T MMHGL\$ R50 110 053 230 002 000 041 000 001 324 230 377 114 044 302 110 107 114 244 377 236 030 BHGL\$ ■ 000 000 000 000 000 041 350 001 RED R70 016 004 000 000 000 017 350 001 056 145 012 343 130 343 055 243

After execution is halted at the breakpoint, you may single step a specified number of instructions using the TRACE instruction. For example, to trace the next 10 program steps, type:

#### TRACE 10

The TRACE instruction will give you status information for each of those 10 steps, as well as the contents of memory. TRACE 10 will output the following information:

PC 114335	DR 57	AR 55	0	CY Ø	NG O	LZ 1	ZR 1	RZ 1	00	DC O	E 00	BKP1 114334	BKP2 000000	PTR1 0370014	PTR2 0370013	R0M 000
PC 031741	DR 57	AR 55	0 V 0	CY O	NG O	LZ 1	ZR:	RZ 1	0 D 0	DC O	E 00	BKP1 114335	BKP2 000000	PTR1 0370014	PTR2 0370013	R0M 000
PC 031742	DR 57	AR 55	0	CY	NG O	LZ 1	ZR 1	RZ 1	0D 0		E 00	BKP1 031741	BKP2 000000	PTR1 0370014	PTR2 0370013	R0M 000
PC 031743	DR 57	AR 55	0 V	CY O	NG O	LZ 1	ZR 1	RZ 1	7.7	DC Q	E 00	BKP1 031742	BKP2 000000	PTR1 0370014	PTR2 0370013	ROM 000
PC 031744	DR 65	AR 55	0 V O	CY 0	NG O	LZ 1	ZP 1	RZ 1	00 0	DC 0	E 00	BKP1 031743	BKP2 000000	PTR1 0370014	PTR2 0370013	R0M 000
PC 031747	DR 65	AR 55	0 V	CY O	NG O	LZ 1	ZR Ó	RZ 0	0 D 1	DC O	E 00	BKP1 031744	BKP2 000000	PTR1 0370014	PTR2 0370013	R0M 000
PC 031750	DR 65	AR 55	0٧	CY O	NG O	LZ 1	ZR O	RZ 0	0 D 1	DC O	E OO	BKP1 031747	BKP2 000000	PTR1 0370014	PTR2 0370013	R0M 000
PC 031751	DR 65		0 V 0	CY O		LZ 1	ZR O	2 0	0.0	DС О	E 00	BKP1 114333	BKP2 000000	PTR1 0370014	PTR2 0370013	RDM 000
R00 00 R10 30 R20 2 R30 00 R40 0 R50 1 R60 00	07 33 41 15 10	1 200 230 000 000 053 000	324 000 230 000	2 05 0 2 3 0 4 2 0 0 0 0	12 1	4 351 075 010 267 000 000	5 06: 21: 01: 23: 06: 04: 25:	3 1 0 0 0 3 3 0 2 0 3 1 0 5 2	02 01 10 33 50 00	7 204 001 200 230 001 000 000	110 040 233 320 324 302	0 107 1 0 040 0 1 230 2 0 230 3 4 230 3	14 102 14 044 000 000 77 230 21 230 77 377 07 114	102 040 000 000 307 230 233 230 110 107 244 377	002 053 040 040 000 000 277 230 017 231 114 044 236 030 055 243	HGLB, + HGL\$B ? G ? P O T ##HGL\$ BHGL\$# .e cXc~#

To continue execution after a breakpoint or after a TRACE instruction, press [RUN].

To run this program without halting, type:

HGL\$ ("string")

after the LOADBIN instruction.

#### 7.3 CRT Control

Source Code: ALPHAS

Object Code: ALPHAB

```
1010 1* This binary program implements three CRT control statements:
         AWRITE [(row), (column)][, (string)]
1020 1*
          AREAD (string variable)
1030 1*
1040 1*
          START CRT AT (absolute line #)
1050 |* AWRITE allows you to do one of three things:
1050 1*

    force ALPHA mode without moving the cursor position

1070 1*
             2) force ALPHA mode and move the cursor to a position which*
                is relative to the top left of the current screen
1080 1*
1090 1*
            3 force ALPHA mode and move the cursor to new position
               and output a string at that location, leaving the cursor*
1100 1*
                positioned at the beginning of the string.
1110 1*
          In all cases the cursor is not actually displayed, until some *
1130 |*
           other normal cursor movement occurs.
1140 1* AREAD allows you to read a string of characters from the CRT into *
        a string variable. Usually the cursor will have been moved to *
1150 1*
1160 I*
          the correct position with the AWRITE statement.
1170 I* START CRT AT allows you to scroll the display up or down or jump *
        to an entirely different page, all under program control.
1180 1*
          NOTE: this routine does not change the cursor's location in
1190 1*
          CRT memory, so the cursor may get last off of the screen when *
1200 1*
         this command is used. It can be brought back by use of the AWRITE statement, or by using the Home Cursor key.
1210 !*
1220 1*
1230 I* ALPHAB returns the revision date of the binary program.
1250 1*
1260 !* An example of how this binary might be used in BASIC is:
1270 I*
        110 FOR I+1 TO 1000
          120 START CRT AT IP(RND*50)
1280 1*
1290 1*
         130 AWRITE RND*16, RND*80 @ AREAD A$
1300 1*
          140 AWRITE RND*16, RND*80, A$
         150 NEXT I
1310 1*
1320 !* This is guaranteed to turn any intelligent display into nonsense. *
1330 1*
! BINARY PROGRAM NUMBER
1350 MYBPGM# EQU 52
1380
            NAM 52, ALFA
                                  I NAME BLOCK FOR BINARY
                                 ADDRESS OF RUNTIME ADDRESSES
            DEF RUNTIM
1370
                                 I ADDRESS OF ASCII TABLE
            DEF ASCIIS
1380
                                 ADDRESS OF PARSE ADDRESSES
1390
            DEF PARSE
                                 I ADDRESS OF ERROR MESSAGES
            DEF ERMSG
1400
1410 DEF INIT
1420 RUNTIM BSZ 2
                                 HDDRESS OF INITIALIZATION ROUTINE
PLACE HOLDER
RUNTIME LABEL FOR 'AWRITE'
                              RUNTIME LABLE (S.)
RUNTIME FOR 'AREAD'
CRT TOP LINE
RUNTIME FOR REVISION
     DEF AREAD.
1470
1440
            DEF STARTAT.
1450
1460
            DEF REV.
```

```
1470 PARSE BSZ 2
                                   I PLACE HOLDER
1480 DEF ALPHAP I PARSE LABEL FOR 'AWE
1490 DEF AREADP I PARSE LABEL FOR 'ARE
1500 DEF STARTATP I PARSE FOR TOP LINE
1510 BYT 377,377 I END OF RELOCATABLES
                                   PARSE LABEL FOR 'AWRITE'
PARSE LABEL FOR 'AREAD'
1530 ASCIIS BSZ 0
1540 ASP "AWRITE" | TOKEN 1
1550 ASP "AREAD" | TOKEN 2
1560 ASP "START CRT AT" | TOKEN 3
1570 ASP "ALPHAB" | TOKEN 4
1580 ERMSG BYT 377 | END OF F
                                    I END OF ASCII TABLE
1730 ALPHAP PUBD R43,+R6 | SAVE TOKEN NUMBER
1740 | JSB *NUMVA+ | TRY TO GET A NUMBER
1750 | JEZ DKRY | MUST BE JUST 'AWRITE'
1880 AREADP PUBD R43,+R6 ! SAVE THE TOKEN
1890 JSB *SCAN ! LET'S DO A SCAN
1900 JSB *STRREF ! MUST BE A STRING REFERENCE
1910 JMP OKAY | FINISH THE PARSE
1930 BYT 0,56 | NO PARAMETERS, STRING FUNCTION
1940 REV. BIN | FOR ADMD R45, *BINTAB
1950 LDM R43, *40D,0 | LOAD THE LENGTH OF THE STRING
1960 DEF DATE | AND THE ADDRESS OF THE STRING
1970 BYT 0 | (MUST BE THREE BYTE ADDRES
1980 ADMD R45, *BINTAB | MAKE THE ADDRESS ABSOLUTE
1990 PUMD R43, +R12 | PUSH IT ALL ON THE OPERATING STACK
2000 RTN | DONE
                                               (MUST BE THREE BYTE ADDRESS)
2010 ASC "30.102:veR 2891 .oC drakcaP-ttelweH )c("
2020 DRTE BSZ 0
                                    ! PLACE HOLDER FOR THE LABEL (ADDRESS)
```

			BASIC STATEMENT
2050	ALFA.	BIN !	FOR MATH
2060		LDBD R37, *CRTSTS	CHECK CRT STATUS
2070		JPS INALPHA!	JIF ALREADY IN ALPHA MODE
2080		JPS INALPHA!	IF NOT, MAKE IT SO
2090	INALPHAI	CMMD R12, = TOS	ANYTHING ON THE R12 STACK
2100		JZR NO-ADR I	JIF JUST 'AWRITE'
2110		JSB *DECUR2 !	JIF JUST 'AWRITE' KILL BOTH POSSIBLE CURSORS
2120		JSB =HMCURS !	MOVE THE CURSOR TO THE HOME POSITION
2130		LDMD R14. *BINTAB	BECAUSE I'M RELATIVE
2140		CLM R43	FRKE O STRING LENGTH
2150		LDM R20.R12	COPY OF R12
2160		SBM R20. = 25.0	SUBTRACT 25
2170		CMMD R20. = TOS	WHAT'S ON R12
2180		JNZ A-ONLY	JIF ONLY X.Y
2190		POMD R43R12	GET LENGTH AND ADDRESS OF STRING
2200	B-DNLY	STMD R43.X14.58V-\$	SAVE LENGTH AND ADDRESS
2210		ISB •TWOB	GET TWO BINARY NUMBERS OFF OF R12
2220	CAL CADE	DCM R56	DECREMENT 'Y'
2230	arrest to the state of the stat	ING GOT-IT	ITE ADDRESS FIGURED DUT
2240		8DM R46 =120 0	ADD TO GET TO NEXT LINE
2250		TMP CRICARR	TRY FOR BNOTHER ONE
2260	GOT-IT	STM R46 R24	CORY ADDRESS DISPLACEMENT TO 26
2270	30 1 1	TER -MOUCES	MOVE THE CURSOR
2280		I DMD R43 X14 SBV-\$	ANYTHING ON THE R12 STACK JIF JUST 'AWRITE' KILL BOTH POSSIBLE CURSORS MOVE THE CURSOR TO THE HOME POSITION BECAUSE I'M RELATIVE FAKE O STRING LENGTH COPY OF R12 SUBTRACT 25 WHAT'S ON R12 JIF ONLY X,Y GET LENGTH AND ADDRESS OF STRING SAVE LENGTH AND ADDRESS GET TWO BINARY NUMBERS OFF OF R12 DECREMENT 'Y' JIF ADDRESS FIGURED OUT ADD TO GET TO NEXT LINE TRY FOR ANOTHER ONE COPY ADDRESS DISPLACEMENT TO 26 MOVE THE CURSOR GET LENGTH AND ADDRESS OF STRING BACK GET LENGTH JIF NO LENGTH SET MEMORY POINTER TO STRING ADDRESS GET LENGTH GET A CHARACTER WAIT FOR CRT NOT BUSY
2290		I DM P56 P43	GET LENGTH
2300		17P NO-8DP	ITE NO LENGTH
2310		STMD D45 *PTD7	SET MEMORY POINTER TO STRING ADDRESS
2320		I DM P36 P43	GET LENGTH
2330	BLOB	I DRI P32 =PTP2-	GET A CHARACTER
2340	TIEUT	19B #CHKSTS*	WRIT FOR CRT NOT BUSY
2350		CTDD D22 *CDTDDT	STORE IT
2350		DCM PRE	BNY CHARACTERS LEET
2370		1N7 DLOD	THE THERE GOE
2370	NO-DDD	DTKI ILUF	OLI DONE
2300	NU-HDK	TS 114 ;	STORE IT ANY CHARACTERS LEFT JIF THERE ARE ALL DONE ************************************
2400		BYT 241	BASIC STATMENT
2410	eeren	BYT 241	FOR MATH
2420	THE PIECE	I DBD D37 *CDTCTC I	CET CPT STATUS
2430		BIN LDBD R37, *CRTSTS ! JPS INALPHA# ! JSB *ALPHA. ! JSB *DECUR2 ! POMD R73, -R12 ! STM R73, R55 ! PUMD R73, +R12 ! CLB R57 !	JIF ALREADY IN ALPHA MODE
2440		ISB #BI PHB	IF NOT MAKE IT SO
2450	THE PHE	ISB *DECUP2	KILL THE CURSORS
2450		POMD P73 -P12	CET STRING STUFF
2470		STM D73 D55	COPY TO 55
2490		DIMD 073 4012	DUCK THE STHEE BACK
2400		CIB 057	CLEON MCD
2500		JSB *RESMEM	LET'S GO RESEARVE SOME MEMORY
2500		JOD -KEUHEN	COPY 55 TO 75
2510 2520		STM R55,R73 !	COPY SINK ADDRESS
2530		STM R65,R75	SET MEMORY POINTER
2540		STMD R65,*PTR2 ! PUMD R73,+R12 !	PUSH STRING STUFF ONTO R12
2550		TSM R55	HOW BIG CAN I GO
2550		JZR DO-STO	: 13:50 B
2570		LDMD R34, = CRTBYT	JIF 0 GET CURRENT POSITION
2580		PUMD R34,+R6	SAVE IT
2590		JSB =BYTCR!	SET CURRENT POSITION
- No. of the last		The second secon	mean in membranism to it. to be seen in a men's

```
2600 ALOOP JSB *INCHR ! GO GET A CHARACTER
2610 STBI R32, *PTR2- ! STORE IT
             JSB =RTCUR.
2620
                                   ! MOVE 1 BYTE
            DCM R55
                                   I ANY MORE
2630
            JNZ BLOOP
           JNZ RLOOP | JIT THERE THE
POMD R34, -R6 | GET OLD CRTBYT BACK
JSB *BYTCR! | SET CURRENT POSITION
JSB *STOST | SAVE IT AWAY
RTN | ALL DONE
2640
2650
2660
2670 DO-STO JSB •STOST
2680 RTN
2700 | * START CRT AT THE SPECIFIED NUMBER
BYT 241
2730 STARTAT, JSB *ONEB
                                   I GET A NUMBER OFF OF R12
2740 BCD
2750 LLM R#
                                   ! FOR MATH
                                   1 *16
2760
            BIN
                                   I FOR THE REST
            STM R#,R#
                                   ! COPY IT
2770
           LLM R#
LLM R#
ADM R#,R#
STM R#,R#
                                   1 *32
2780
2790
                                   1 *64
2800
                                    1 *80
          STM R#,R# | COPY TO 46
LDMD R#,*ASIZE | GET ALPHA SIZE INTO 76
DRP R46 | GET READY FOR 'MOD'
JSB *MOD | MOD IT
STM R#,R34 | COPY RESULT TO 34 FOR 'SADI'
JSB *SADI | SET CRT START ADDRESS
RTN | ALL DONE
2810
2820
2830
2840
2850
2860
2870
2910 NUMVA+ DAD 22403
2920 GETCMA DAD 23477
2930 NUMVAL DAD 22406
2940 STREXP DAD 23724
            DAD 10220
2950 ERROR+
           DAD 177715
2960 PTR2-
2970 SCAN
             DAD 21110
2980 STRREF DAD 24056
2990 STREX+ DAD 23721
3000 BINTAB DAD 104070
3010 CRTSTS DAD 177702
3020 ONEB DAD 12153
3030 PTR2 DAD 177714
             DAD 177714
3040 CHKSTS DRD 13204
3050 CRTBAD DAD 177701
3060 CRTDAT DAD 177703
3070 ALPHA. DAD 12413
3080 TOS
             DAD 101744
                                   I DEFINE ADDRESSES
3090 DECUR2 DAD 13467
3100 HMCURS DAD 13661
3110 TWOB DAD 56760
3120 MOVCRS DAD 13771
3130 RESMEM DAD 31741
3140 CRTBYT DAD 100206
3150 BYTCR! DAD 14003
3160 INCHR DAD 14262
```

Section 7: Sample Binary Programs

3170 RTCUR.	DAD 13651	(1
3180 STOST	DRD 46472	4
3190 ASIZE	DAD 104744	1 *
3200 SAD1	DAD 13723	1
3210 MOD	DAD 14216	1
3220	FIN	I TERMINATE ASSEMBLY

## 7.4 Line Input

Source Code: LINPUTS

Object Code: LINPUTB

```
10 !*************************
20 1*
           A KEYWORD THAT IS PARSED INTO MORE THAN ONE TOKEN
30 1*
40 1*
          A TOKEN WITH A CLASS OF 44 (MISC IGNORE AT DECOMPILE)
50 1*
            that makes it invisible when the program is listed
50 I*
70 1*
                 (c) 1982 Hewlett-Packard Co.
BO 1*
90 1*
       This binary program implements the BASIC statement 'LINPUT'
100 1*
110 I* which acts exactly the same as the BASIC statement 'INPUT' except *
120 I* that it will only allow you to input a string value and that
130 I* string value may contain commas and/or quotes. The keyword stands
140 I* for Line INPUT, as it allows the inputing of a line regardless
150 1% of what characters are in that line.
180 1*
190 !* An example of how a BASIC program might use LINPUT is:
200 1*
210 1*
        100 DISP "Address of destination";
220 !*
        110 LINPUT Dest_addr$
230 !*
        120 PRINT# 1; Dest_addr$
240 !*
NRM 51, LINPUT | SET UP PCB, BPGM # IS 51
260
           DEF RUNTIM
                                 I POINTER TO RUNTIME ADDRESS TABLE
270
                                 | POINTER TO TABLE OF ASCII KEYWORDS
280
           DEF ASCIIS
                                  I POINTER TO TABLE OF PARSE ADDRESSES
           DEF PARSE
290
                                  I POINTER TO TABLE OF ERROR MESSAGES
300
            DEF ERMSG
                                  I POINTER TO INITIALIZATION ROUTINE
           DEF INIT
310
330 I* The way an INPUT statement works in the series 80 computers is *
340 I* this: the keyword is actually parsed into two tokens, so the Job *
350 !st of doing an INPUT is split into three parts; two are performed by st
360 !* the two INPUT tokens and the third is performed by the system.
370 I* The first of the two tokens outputs a question mark to the CRT and*
380 1* puts the computer into a pseudo-calculator mode, which is known - *
390 !* as Idle-in-Input, by setting CSTAT (R16) to a 4, and then sets the*
400 1* immediate break bits in XCDM (R17) using "or"with 240(octal). Then*
410 1* the first token terminates its execution by returning to the
420 I* interpreter. The interpreter will see the immediate break bits in * 430 I* R17 and will drop out into the exec Loop. The exec will see that *
440 |* the computer is in Idle-in-Input mode and will simply loop on
450 1% itself. At this point, the user starts typing his input (causing *
460 1* keyboard interrupts, which set bits in R17 and SVCWRD, which cause*
470 !* the exec to call the character editor (CHEDIT), which echoes the
480 | * keys to the CRT, clears the SVCWRD flag, and returns to the exec).*
```

```
490 I* This continues until the END LINE key is pressed, which causes
500 I* CHEDIT to set a flag in the E register which will tell the exec *
510 !* that END LINE has been pressed. This will cause the exec to resume*
520 I* execution of the BASIC program by re-entering the interpreter. * 530 I* The third part of the INPUT is carried out by the second token of *
540 1* the INPUT statement. It takes the input line, parses and executes *
550 l* it, then stores the values in the appropriate variables.
560 I* LINPUT statement works in very much the same way. As a matter
570 !* of fact, the first two LINPUT tokens do nothing but call *
580 !* the runtime code for the first of the INPUT tokens. The difference*
590 I* comes in the second token. For LINPUT, all we want to do is input *
600 1* a literal string with no expressions allowed, so we have no need *
610 !st to parse and execute the input line. All we have to do is reverse st
620 1* the string so that the first character is at the highest address *
630 1* and then store it in the string variable.
BYT 0,0
DEF REV.
650 RUNTIM BYT 0,0
                                      | DUMMY ADDRESS FOR TOKEN #0 RUNTIME
             DEF REV. | ADDRESS FOR TOKEN #1 RUNTIME ROUTINE
DEF LINPT. | ADDRESS FOR TOKEN #2 RUNTIME ROUTINE
DEF LIN$. | ADDRESS FOR TOKEN #2 RUNTIME ROUTINE
670
680
                                       I ADDRESS FOR TOKEN #3 RUTNIME ROUTINE
700 PRRSE BYT 0,0 | DUMMY ADDRESS FOR KEYWORD #0 PARSE 710 BYT 0,0 | DUMMY FOR KEYWORD #1 PARSE (A FUNCTION)
           DEF LINPRS | ADDRESS FOR KEYWORD #2 PARSE ROUTINE
BYT 377,377 | TERMINATE RELOCATION OF ADDRESSES
720
730
750 |* The runtime table has three entries even though the RSCII and #
760 I* parse tables have only two. The third entry in the runtime table
770 1* will only be used in conjunction with the second entry. If you
780 I* look at the parse routine for the second keyword (LINPUT) you will*
790 1* see that it pushes out both tokens 2 and 3. Normally, you want to *
800 1* keep a one for one relationship between entries in the ASCII,
810 1* PARSE, and RUNTIME tables, but there are times when you can play *
B20 1* tricks like this (if you're careful).
840 ASCIIS ASP "LINPUTG" ! KEYWORD #1 (REVISION DATE FUNCTION)
850 ASP "LINPUT" ! KEYWORD #2
850 FRMSG BYT 327 ! TERMINETE ASCIT AND FRAGE TO
             BYT 377
860 ERMSG
                                       I TERMINATE ASCII AND ERROR MESSAGE TABLES
880 ERR89 JSB *ERROR+ ! SET ERROR FLAGS IN R17 AND 'ERRORS'
890 BYT 89D ! SYSTEM ERROR MESSAGE #89 'INVALID PARAM'
910 LINPRS LDM R55, 2,51,371 | LOAD TOKEN#, BPGM#, AND SYSTEM TOKEN
910 LINPRS LDM R55, *2,51,371 | LUMD TUKEN#, BPGN#, HND STSTEN TUKEN
920 STMI R55, *PTR2- | STORE THEM ALL OUT TO PARSE STREAM
930 JSB *SCAN | SCAN THE INPUT STREAM FOR NEXT TOKEN
940 JSB *STRREF | TRY TO GET A STRING VARIABLE REFERENCE
950 JEZ ERRB9 | JIF NOT THERE, ERROR CONDITION
960 LDM R55, *3,51,371 | ELSE LOAD SECOND TOKEN#, BPGM#, AND SYS
970 STMI R55, *PTR2- | STORE THEM OUT TO PARSE OUTPUT STREAM
980 INIT RTN | DONE FOR PARSING AND INITIALIZING
1000 I* LINPT, is the runtime code for the first of the two LINUT tokens.*
1010 1* It is responsible for the output of the question mark to the CRT
1020 |* and putting the computer into Idle-in-Input mode.
1040 BYT 241 | ATTRIB., BASIC STATEMENT LEGAL AFTER THEN
1050 LINPT. JSB *INPUT. | DO QUESTION MARK AND SET R16*4
1060 RTN | DONE, WALT FOR INPUT
```

```
1080 I* LIN$, is the runtime code for the second of the the two LINUT
1090 1* tokens. It is responsible for reversing the string in memory so it*
1100 !* will be ready for storing into the string variable, and then doing*
1110 1* the actual store (by calling STOST). The R12 stack will already
1120 I* have been set up for the variable store by the tokens parsed by
1130 I* STRREF.
BYT 44
1150
                                       ! ATTRIBUTE, MISCELLANEOUS IGNORE
1160 LINS.
               BIN
                                       I BIN MODE FOR COUNTING
              LDMD R32, *INPTR
1170
                                       ! FETCH ADDRESS OF STRING THAT WAS INPUT
              STM R32, R14
CLM R36
                                      I SAVE A COPY
1180
1190
                                      I PRE-SET LENGTH TO ZERO
1200 CHRCNT POBD R35,+R32 | GET THE NEXT BYTE FROM INPUT STRING
1210 CMB R35,*15 | IS IT A CARRIAGE RETURN CHARACTER?
1220 JZR ENDOF$ | JIF YES, WE'VE FOUND THE END AND LENGTH
1230 ICM R36 | ELSE INCREMENT THE LENGTH
1230
1370 JCY DONE
1380 LDBD R30,R14
1390 POBD R31,-R32
1400 STBD R30,R32
          LDBD R30,R14 | ELSE GET BYTE FROM FRONT
POBD R31,-R32 | AND A BYTE FROM THE BACK
STBD R30,R32 | STORE THE FRONT BYTE IN BACK
PUBD R31,+R14 | AND THE BACK BYTE IN FRONT
JMP DOLOOP | LOOP TIL STRING IS REVERSED IN PLACE
PUMD R36,+R12 | PUSH THE LENGTH OF STRING TO STACK
PUMD R65,+R12 | PUSH THE ADDRESS OF STRING TO STACK
JSB *STOST | STORE THE STRING TO THE VARIABLE AREA
RTN | DONE
                                      I JIF YES
1410
1420
1430 DONE
1440
1450
1460
               RTN
                                       DONE
1480 I* This is the runtime code for the revision date function, which is*
1490 I* a string function with no parameters which always returns the same*
1500 I* string value, the copyright notice and the revision code.
```

1520		BYT 0,56	ATTRIBUTES, STRING FUNCTION, NO PARAMETERS
1530	REV.	LDM R43, =44,0	LOAD LENGTH OF THE STRING
1540			AND THE ADDRESS OF THE STRING
1550		BYT 0	(IT NEEDS TO BE A THREE BYTE ADDRESS)
1560		BIN	I BIN MODE FOR ADDRESS MATH
1570		ADMD R45, -BINTAB	MAKE THE ADDRESS ABSOLUTE
1580			! PUSH THE LENGTH AND ADDRESS TO THE STACK
1590		RTN	I DONE
1600		ASC "82.111 .veR 2891	drakcaP-ttelweH )c("   THE REVISION STRING
1610	DATE	BSZ 0	NEED LABEL HERE TO GET RIGHT ADDRESS
1620	] ******	********	**********
1630	BINTAB	DAD 104070	1
1640	ERROR+	DAD 10220	1
1650	INPTR	DAD 101143	1
1660	INPUT.	DAD 16314	LABEL DEFINITIONS
1670	PTR2-	DAD 177715	
1680	SCAN	DAD 21110	f .
1690	STOST	DAD 46472	T.
1700	STRREF	DRD 24056	F
1710		FIN	I TERMINATE ASSEMBLY

## 7.5 Taking the KYIDLE Hook and Buffering the Keyboard

Source Code: KEYS

Object Code: KEYB

```
1010 1*
1020 |* TAKING THE 'KYIDLE' HOOK AND BUFFERING THE KEYBOARD
1030 1*
1040 1*
                  (c) 1981 Hewlett-Packard Co.
1050 1*
1050 1*
        THIS BINARY PROGRAM TAKES OVER THE 'KYIDLE' HOOK AND PUTS ALL
1070 I* KEYS PRESSED INTO A BUFFER EXCEPT FOR THOSE KEYCODES LISTED IN
1080 I* THE TABLE STARTING AT 'KEYTAB' (RIGHT NOW, THOSE KEYS TO BE LEFT
1090 1* FOR THE SYSTEM TO HANDLE ARE THE SOFT KEYS AND THE RESET KEY. THIS*
1100 I* COULD EASILY BE CHANGED BY MODIFYING THE 'KEYTAB' TABLE), THE *
1110 I* BINARY ALSO WATCHES FOR 'SHIFT END LINE' AND 'SHIFT UP ARROW' *
1120 I* (WHICH IS THE 'HOME' KEY, ('UP ARROW' AND 'HOME' ACTUALLY GENERATE*
1130 !* THE SAME KEYCODE AND CAN ONLY BE DIFFERENTIATED BY CHECKING TO SEE*
1140 I* IF THE SHIFT KEY IS UP OR DOWN.)) WHEN 'END LINE' OR 'UP ARROW' IS*
1150 1* PRESSED WITH THE SHIFT KEY DOWN, THE BINARY PROGRAM CHANGES THE
1160 I* KEYCODE TO A DIFFERENT UNIQUE KÉYCODE SO THE BASIC PROGRAM CAN
1170 I* TELL THE DIFFERENCE. THIS, AND SIMILAR TECHNIQUES, COULD BE
1180 I* APPLIED TO MOST OF THE KEYBOARD.
1210 1*
1220 I* The following is a sample BASIC program showing how this binary
1230 I* program can be used:
1240 1*
1250 !*
         100 TAKE KEYBOARD
1260 I* 110 A$=KEY$
       120 IF A$ = " THEN 110
1270 1*
         130 IF A$="E" THEN 200
1280 1*
         140 DISP "THAT WAS THE " & A$ & " KEY."
1290 1*
1300 1*
        150 GOTO 110
1310 1*
        200 RELEASE KEYBOARD
1320 1*
        210 DISP "DONE"
        220 END
1330 1*
1340 1*
1360 MYBPGM# EQU 50
                                  | BINARY PROGRAM NUMBER
1370
            NAM 50, KEYS
                                 I NAME BLOCK FOR BINARY
     1380
1390
1400
1410
1420
1430 RUNTIM BSZ 2
1440 DEF TAKE.
1450
```

```
I RUNTIME FOR 'KEYS'
1460
             DEF KEYS.
1470
             DEF REVDATE.
                                    RUNTIME FOR REVISION
1480 PARSE BSZ 2
                                     I PLACE HOLDER
            BSZ Z ! PLACE HOLDER

DEF COMPARS ! PARSE ROUTINE FOR 'TAKE KEYBOARD'

DEF COMPARS ! PARSE ROUTINE FOR 'RELEASE KEYBOARD'

BYT 377,377 ! END OF RELOCATABLES
1490
1500
1510
1530 ASCIIS BSZ 0
            ASP "TAKE KEYBOARD"
                                      TOKEN 1
1540
             ASP "RELEASE KEYBOARD" | TOKEN 2
1550
             ASP "KEY$" ! TOKEN 3
ASP "REV DATE" ! TOKEN 4
1560
1570
1580 ERMSG BYT 377
                                     I END OF ASCII TABLE
1600 !* BECAUSE THIS PROGRAM TAKES OVER 'KYIDLE', SOME SPECIAL TRICKS *
1610 1* ARE NEEDED. 'KYIDLE' IS AN INTERRUPT HOOK WHICH MEANS THAT THE
1620 1* BASE ADDRESS OF THIS BINARY PROGRAM MAY NOT BE IN 'BINTAB'. A * 1630 1* METHOD IS NEEDED FOR THE HOOK ROUTINE ('USEKEY' IN THIS CASE) TO *
1840 I* KNOW WHAT THE BASE ADDRESS IS. SINCE THE 'KYIDLE' HOOK IS ? BYTES *
1650 I* LONG AND IT ONLY TAKES 4 BYTES TO DO 'JSB *USEKEY' & 'RTN', 3
1660 I* BYTES ARE LEFT UNUSED (AND THAT WE CAN BE SURE NO DNE ELSE IS
1670 I* GOING TO USE, AS LONG AS THIS BINARY HAS THE HOOK, WHICH IS AS
1680 1* LONG AS IT MATTERS), TWO OF THESE BYTES ARE USED TO STORE THE 1690 1: BASE ADDRESS OF THIS BINARY PROGRAM. WE'VE NAMED THE LOCATION
1700 |* 'MYBTAB' AND DEFINED ITS ADDRESS AS 4 HIGHER THAN THAT OF 'KYIDLE'*
1710 |* (103703 AND 103877 RESPECTIVELY.)
1720 1* THE 'INIT' ROUTINE DOESN'T HAVE TO DO ANYTHING IN THIS PROGRAM *
1730 I* SINCE 'LOAD' AND 'SCRATCH' CAN'T BE PERFORMED WHILE THE BINARY
1740 I* HAS THE HOOK, AND DURING A "RESET" THE SYSTEM WILL HAVE ALREADY
1750 I* PUT 'RTN's BACK INTO 'KYIDLE'. WE ONLY TAKE THE HOOK WHEN A
1760 1* 'TAKE KEYBOARD' COMMAND IS EXECUTED, SO THERE'S NOTHING FOR INIT *
1770 I* TO DO.
1780 I* THE BASIC PROGRAM WRITER NEEDS TO BE VERY CAREFUL, HOWEVER,
1790 I* USING THIS BINARY, BECAUSE IF HE WERE TO EXECUTE A 'STOP' OR 'END'*
1800 I* COMMAND WHILE THE HOOK IS TAKEN, THE KEYBOARD WILL EFFECTIVELY BE * 1810 I* LOCKED UP EXCEPT FOR THE 'RESET' KEY AND, THUS, 'RESET' WOULD THEN*
1820 I* BE THE USERS ONLY RECOURSE.
! ALL DONE
1840 INIT RTN
1860 I* NEITHER 'TAKE KEYBOARD' OR 'RELEASE KEYBOARD' HAVE ANY PARAMETERS *
1870 1* SO THEY BOTH USE THE SAME PARSE ROUTINE, WHICH SIMPLY PUSHES OUT * 1880 1* THE THREE BYTE SEQUENCE FOR THE KEYWORD AND THEN DOES A 'SCAN' FOR*
1890 I* THE SYSTEM, SO THAT R14 WILL HAVE THE NEXT TOKEN WHEN WE RETURN. *
1910 COMPARS LDM R56, -50,371 | BPGM # AND SYSTEM TOKEN
1920 LDB R55,R43 | GET THE BINARY PROGRAM TOKEN #
1930 STMI R55, =PTR2- | STORE IT ALL DUT TO PARSE STACK
1940 JSB -SCAN | DO A SCAN FOR THE SYSTEM
             RTN
1950
```

```
1970 I* 'REV DATE' IS A STRING FUNCTION WITH NO PARAMETERS WHICH RETURNS *
1980 I* AS ITS STRING VALUE THE COPYRIGHT STATEMENT AND REVISION CODE OF *
1990 I* THE BINARY PROGRAM.
BYT 0,56 NO PARAMETERS, STRING FUNCTION
2010
2020 REVDATE. BIN ! FOR ADMD R45, =BINTAB
2030 LDM R43, •40D,0 ! LOAD THE LENGTH OF THE STRING
2040 DEF DATE ! AND THE ADDRESS OF THE STRING
2050
            BYT O
                                              (MUST BE THREE BYTE ADDRESS)
            RDMD R45, = BINTAB
PUMD R43, +R12
RTN
                                  MAKE THE ADDRESS ABSOLUTE
2080
                                  | PUSH IT ALL ON THE OPERATING STACK
2070
2080
             RTN
            ASC "31.102:veR .oC drakcaP-ttelweH 2891 )c("
2090
2100 DATE BSZ 0 | PLACE HOLDER FOR THE LABEL (ADDRESS)
2120 1* THIS IS THE TABLE OF KEYS THAT THE BINARY PROGRAM SHOULD LET THE *
2130 I* SYSTEM HANDLE, AND IT SHOULD NOT PUT THEM IN THE BUFFER. THE TABLE* 2140 I* IS TERMINATED BY A 377, WHICH IS A KEYCODE THE KEYBOARD CONTROLLER*
2150 I* IC IS INCAPABLE OF GENERATING.
2170 KEYTAB BYT 200
                                   FKI
2180
            BYT 201
                                   1 K2
2190
             BYT 202
                                   1 K3
            BYT 203
BYT 241
2200
                                   1 K4
2210
                                   1 K5
2220
            BYT 242
                                   1 K6
2230
            BYT 234
                                   1 K7
2240
            BYT 204
                                   1 K8
2250
             BYT 205
                                   1 K9
2260
             BYT 206
                                   1 K10
            BYT 207
2270
                                   1 K11
            BYT 245
2280
                                   1 K12
2290
            BYT 254
                                   L K13
2300
            BYT 223
                                    1 K14
            BYT 213
BYT 377
2310
                                  1 RESET
2320
                                   ! END OF INVALID KEY TABLE
2340 !* THIS IS THE RUNTIME ROUTINE FOR THE 'TAKE KEYBOARD' KEYWORD, IT *
2350 I* INITIALIZES POINTERS TO THE BEGINNING AND END OF THE KEYBOARD
2360 I* BUFFER, WHICH EXISTS FARTHER DOWN IN THE BINARY PROGRAM, TAKES
2370 I* OVER THE 'KYIDLE' HOOK, AND INVALIDATES THE KEY REPEAT FLAG. IF * 2380 I* THE KEY REPEAT FLAG IS VALID, THE LAST KEY IS TAKEN FROM THE * 2390 I* BUFFER (USING THE 'KEY$' FUNCTION), AND A KEY IS STILL DEPRESSED *
2400 I* THE LAST KEY WILL BE PUT BACK IN THE BUFFER SO THAT IT WILL REPEAT*
2410 I # AS LONG AS THE KEY IS HELD DOWN.
BYT 241
2430
            LDMD R46, BINTAB | FOR RELATIVE ADDRESSING
LDM R30, KEYBUF | GET ADDRESS OF KEYBOARD BUFFER
ADM R30, R46 | MAKE IT ABSOLUTE
2440 TAKE.
2450
2460
            STMD R30, X46, KEYPTR | INITIALIZE KEY POINTER
ADM R30, =80D, G | POINT TO END OF BUFFER
2470
2480
```

```
2490 STMD R30, X46, KEYEND | INITIALIZE KEYEND
2500 LDM R30, = USEKEY | ADDRESS OF KEYBOARD SERVICE ROUTINE
2510 ADM R30, R46 | MAKE IT ABSOLUTE
2520 STM R30, R43 | COPY TO 43&44
2530 LDB R45, = 236 | 45*'RTN'
2540 LDB R42, = 316 | 42*'JSB'
2550 TAKEIT STMD R#, = KYIDLE | STORE DUT RTN'S OR JSB=USEKEY, RTN, BINTAB
2560 LDB R#, = 377 | INVALID REPEAT FLAG
              STBD R#, X46, LASTKEY ! SET IT
2570
2580
              RTN
2600 I* THIS IS THE RUNTIME ROUTINE FOR THE 'RELEASE KEYBOARD' KEYWORD.
2610 I* ALL IT DOES IS PLACE RETURNS BACK INTO THE 'KYIDLE' HOOK, THUS,
2620 I* GIVING UP CONTROL OF THE KEYBOARD.
BYT 241
2640
2650 RELEAS. LDMD R46, *BINTAB | GET BPGM'S BASE ADDRESS
2690 !* 'USEKEY' IS AN INTERRUPT SERVICE ROUTINE SO IT MUST BE CAREFUL TO*
2700 I* SAVE ALL CPU STATUS AND CONTENTS AND THEN RESTORE THEM WHEN DONE. *
2710 I* THE SYSTEM HAS ALREADY DONE R 'SAD' BEFORE IT DID THE 'JSB' TO
2720 I* 'KYIDLE'. THE ROUTINE CHECKS TO SEE IF THE BUFFER IS FULL AND IF \,\star\, 2730 I* SO THROWS THE CURRENT KEYHIT AWAY. IT THEN CHECKS FOR THE SHIFTED \,\star\,
2740 I* 'UP ARROW' OR 'END LINE' KEYS AND IF SO MODIFIES THE KEYCOD TO
2750 1* MATCH, IT THEN CHECKS THE 'KEYTAB' TABLE TO SEE IF THIS KEY SHOULD*
2760 I* BE IGNORED. IF IT IS IN THE TABLE, THE ROUTINE JUST CLEANS UP A
2770 I* LITTLE AND RETURNS BACK INTO THE SYSTEM KEY HANDLING ROUTINE.
2780 I* OTHERWISE, IT PUTS THE NEW KEYCODE IN THE BUFFER AND UPDATES THE 2790 I* BUFFER POINTER. IT THEN FIGURES OUT WHAT THE DRP SHOULD BE WHEN
2800 1* IT RETURNS FROM THE INTERRUPT SERVICE, AND PLACES A DRP COMMAND
2810 I* WHERE IT WILL BE EXECUTED JUST BEFORE RETURNING (THIS IS SO THE
2820 I* EXTENDED MEMORY CONTROLLER CAN KEEP TRACK OF THE DRP FOR MULTI-
2830 (* BYTE OPERATIONS.) IT THEN RESTORES REGISTERS, THROWS AWAY TWO
2840 1* RETURN ADDRESSES, AND RETURNS TO WHATEVER WAS HAPPENING BEFORE
2850 I* THE KEYBOARD INTERRUPTED.
```

```
JZR KEYLOOP1 | JIF IT IS
CMB R23,R22 | IS THIS KEY I
JNZ KEYLOOP | JIF NO MATCH
JSB X26,FIXUP-R6 | FIX UP THE R6
3120
                                   I IS THIS KEY INVALID
3130
3140
                                  I FIX UP THE R6 STRCK
3150
3160 JMP KEYRTN+ ! FÄLL THROUGH, LET THE SYSTEM HAVE IT
3170 KEYLOOP1 PUBD R22,+R20 ! APPEND TO THE BUFFER
3180 STMD R20,X26,KEYPTR ! UPDATE THE POINTER
3190 RE-START CLB R20
            ICB R20
STBD R20, =KEYCOD | /
JSB X26, FIXUP-R6 | FIX UP THE R6 STACK
TRASH TWO RETURNS
TRASH TWO RETURNS
3200 ICB R20
                                   ) > RESTART THE KEYBOARD SCANNER
3210
3220
3230
                                 RE-ENABLE GLOBAL INTERRUPTS
3240 KEYRTH STBD R#, -GINTEN
3250 DRP
             BSZ 1
                                   I FORCE THE DRP
            PAD
3260
                                   I RESTORE THE STATUS
            RTN
3270
                                   I RLL DONE
3280 KEYRTN+ STBD R#, *GINTEN
                                  ! RE-ENABLE INTERRUPTS
            RTN
3290
3330
3340
3350
3360
3370
3380
3390
            POMD R40, R6

POMD R2, R6

PUMD R30, +R6

LDMI R30, =MYBTAB

| GET 30 BACK
3400
3410
                                  PUT THE RETURN BACK
3420
3430
3450 !* THIS IS THE RUNTIME ROUTINE FOR THE 'KEY$' KEYWORD. IT IS A *
3460 !* STRING FUNCTION WITH NO PARAMETERS WHICH RETURNS A STRING WITH A *
3470 I* LENGTH OF DNE WHOSE SOLE CHARACTER IS THE KEYCODE OF THE FIRST *
3480 I* KEY IN THE KEYBOARD BUFFER. IF THE BUFFER WAS EMPTY, IT RETURNS 3490 I* A NULL STRING (LENGTH=0). WHEN IT TAKES A KEY OUT OF THE BUFFER,
3500 I* IT COLLAPSES ALL THE DTHER KEYCODES IN THE BUFFER AND ADJUSTS THE *
3510 !* BUFFER POINTER.
           3530
3540 KEY$.
3550
3560
3570
3580
3590
3600
3610
            LDM R30,R20 | CDPY 20
POBD R32,+R20 | GET A KEY
3620
3630
            STBD R32,X14,LASTKEY | SAVE LAST KEY FOR POSSIBLE REPEAT
```

```
3650 KEY$1 CMM R22,R20 ! BUFFER COLLAPSED 3660 JZR KEY$2 ! JIF IT IS 3670 P0BD R33,+R20 ! GET A KEY 3680 PUBD R33,+R30 ! MOVE IT DOWN 3690 JMP KEY$1 ! LOOP 3700 KEY$2 DCM R22 ! ADJUST KEYPTR 3710 STMD R22,X14,KEYPTR ! RND RESTORE IT 3720 KEY$2+ CLM R22
3710
3720 KEY$2+ CLM R22
3730 ICM R22
                                          1 > LENGTH OF 1
3740 KEY$2++ PUMD R#,+R12
              LDM R55, *LASTKEY | ADDRESS OF KEYHIT
BYT 0 | ----> R57
                                         1 /
3750
3770 ADMD R55, *BINTAB ! MAKE ADDRESS ABSOLUTE
3780 ICM R55 ! POINT TO AFTER THE KEY
3790 PUMD R55, +R12 ! PUSH ADDRESS OUT
3800 STBD R#, *GINTEN ! RE-ENABLE GLOBAL INTERRUPTS
3810 RTN ! ALL DONE
4020 ERROR+ DAD 10220
4030 PTR2- DAD 177715
4040 SCAN DAD 21110
4050 BINTAB DAD 104070
4060 PTR2 DRD 177714
4070 RDMFL DRD 104065
4080 KYIDLE DRD 103677
                                  | DEFINE SYSTEM ADDRESSES
4090 GINTDS DAD 177401
4100 GINTEN DAD 177400
4110 MYBTAB DAD 103703
4120 KEYCOD DAD 177403
4130 KEYSTS DAD 177402
4140
                FIN
                                         ! TERMINATE ASSEMBLY
```

## 7.6 GET and SAVE

Source Code: GETSAVES

Object Code: GETSAVEB

```
1010 I* This binary program implements the SAVE and GET statements for
1020 !* turning programs into normal strings in a DATA file and turning
1030 | * normal strings back into lines of a BASIC program.
1040 I* The syntax for the two statements is:
          SAVE (file name)[, (beginning line)][, (ending line)]
1050 1*
            GET (file name)
1050 1*
1070 I* SAVE calculates the size of the DATA file needed by listing the
1080 [* program and counting the total length of the strings (plus the
1090 1* three bytes of header per string required by the file manager).
1100 !* It does this by taking over IOTRFC and forcing the select code to *
1110 !* a value that will cause the listed strings to go out through the
1120 I* hook, LSSET is an entry point in the LIST routine that lists the
1130 1* entire program. After the size of the data file is known, it is
1140 I* created (any old one of that name already in existence will be
1150 I* purged first) and then the program is listed again, this time with*
1160 I* the lines (as strings) being printed out to the data file.
1170 I* GET opens the data file, reads a string, copies the string to 1180 I* the input buffer INPBUF, then calls the PARSER, which will parse
1190 l* the line and edit it into the program, if no errors occur. If a
1200 I* parse error occurs, an exclamation point is inserted into the line*
1210 !* after the line number and the line is parsed again as a comment.
1220 I* GET has to create a dummy string variable area in the binary
1230 I* program for the strings to be read into, because RDSTR. does a
1240 I* call to STOST before it returns, and STOST expects all the usual *
1250 I* information on the stack and an associated variable area (in
1260 I* other words, we have to trick the system when we call RDSTR.).
NAM 41, SAVG | SET UP THE PROGRAM CONTROL BLOCK
1280
1290
               DEF RUNTIM
                                        ! PTR TO THE RUNTIME ADDRESSES
              DEF TOKS
                                       I PTR TO THE KEYWORDS
1300
                                   PTR TO THE PARSE ADDRESSES

PTR TO THE ERROR MESSAGE TABLE

PTR TO THE INITIALIZATION ROUTINE

DUMMY RUNTIME ADDRESS FOR TOK# 0

RUNTIME ADDRESS FOR TOK# 1

RUNTIME ADDRESS FOR TOK# 2

RUNTIME ADDRESS FOR TOK# 3

DUMMY PARSE ADDRESS FOR TOK# 0

PARSE ADDRESS FOR TOK# 1

DUMMY PARSE ADDRESS FOR TOK# 3

PARSE ADDRESS FOR TOK# 3

PARSE ADDRESS FOR TOK# 3

PRESE ADDRESS FOR TOK# 3

TERMINATE RELOCATION AND ERROR TABLE
              DEF PARSE
                                       ! PTR TO THE PARSE ADDRESSES
1310
1320
              DEF ERMSG
              DEF INIT
1330
1340 RUNTIM BYT 0,0
1350
               DEF SAVE.
              DEF REVISON.
1360
              DEF GET.
1370
1380 PARSE BYT 0.0
              DEF SAVPARS
1390
               BYT 0,0
1400
               DEF GETPARS
1410
               BYT 377,377
1420 ERMSG
1430 INIT
              RTN
                                       1 NO INITIALIZATION
              ASP "SAVE"
1440 TOKS
              ASP "SAVE" ! KEYWORD #1
ASP "GET SAVE" ! KEYWORD #2
1450
               ASP "GET"
1460
                                        | KEYWORD #3
               BYT 377
1470
                                        I TERMINATE KEYWORD TABLE
```

```
1490 SAVPARS PUBD R43,+R6 | SAVE CURRENT TOKEN
1500 | JSB *STREX+ | GET THE FILE NAME
1500 JSB -STREX+
1630 !***********************************
1640 GETPARS PUBD R43,+R6 | SAVE THE INCOMING TOKEN
1650 | JSB *STREX+ | GET THE FILE NAME
1660 | JEZ ERR | JIF NOT THERE
1670 | JMP PARSCOMN | ELSE FINISH UP
1690 BYT 241 ! BASIC STATEMENT, LEGAL AFTER THEN
1700 SAVE. JSB *CLEAR. ! CLEAR THE CRT
1710 LDMD R10, *BINTAB ! GET DUR BASE ADDRESS
1720 LDM R26, *SAVING ! GET THE RELATIVE ADDRESS OF MSG
1730 ADM R26,R10 ! MAKE IT ABSOLUTE
1740 LDM R36,*20,0 ! LOAD THE LENGTH OF THE MSG
1750 JSB *DUTSTR ! DUTPUT THE MSG
1760 LDMD R41,*IOTRFC ! SAVE THE REAL HOOK CONTENTS
1770 STMD R41, *10 SAVIATEC ! STORE IT DURY
                     STMD R41,X10,SAVIOTEC | STORE IT HWAY
LDMD R40,*SCTEMP | SAVE THE REAL SELECT CODE
STMD R40,X10,SAVSCTEM | STORE IT HWAY
LDM R72,*231,231,11,0,0,0 | LOAD DEFAULT LIST PARAMETERS
1770
1780
 1790
1800
                       STMD R43, X10, FILENAME | SAVE IT AWAY
1950
                      STMD R43,X10,FILENAME | SAVE IT AWAY

CLM R50 | SET UP FOR A FLOATING POINT 1

LDB R57,=10C | THAT FINISHES IT

PUMD R50,+R12 | PUSH TO STACK FOR ASSIGN# 1 TO

PUMD R43,+R12 | PUSH FILE NAME BACK

PUMD R10,+R6 | SAVE OUR BASE ADDRESS

JSB =ROMJSB | SELECT THE MSTORAGE ROM

DEF ASSIG. | ASSIGN BUFFER # 1 TO FILE

VAL MSROM# | ROM TO SELECT

POMD R10,-R6 | RECOVER OUR BASE

CMB R17,=300 | ANY ERRORS?

JNC ITSTHERE | JIF NO, IT WAS THERE AND DATA FILE

LDBD R20,=ERRORS | GET REASON
1960
1970
1980
1990
2000
2010
2020
2030
2040
2050
 2060
 2070
```

```
| FILE NAME ERROR? | FILE NAME ERROR? | 2090 | JZR CREATIT | JIF IT WASN'T THERE | 2100 | GTO RESTORE | ELSE BAIL OUT | 2110 ITSTHERE PUMD R10,+R6 | SAVE OUR BASE | MAKE SURE STACK LODKS GOOD | 2130 | LDMD R73,X10,FILENAME | GET THE FILE NAME BACK | 2140 | SUMD R73,X10,FILENAME | GET THE FILE NAME BACK | 2140 | SUMD R73,X10,FILENAME | GET THE FILE NAME BACK | 2140 | SUMD R73,X10,FILENAME | GET THE FILE NAME BACK | 2140 | SUMD R73,X10,FILENAME | GET THE FILE NAME BACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | GET THE STACK | 2140 | SUMD R73,X10,FILENAME | SUMD R73,X10,FIL
    2130 LDMD R73,X10,FILENAME ! GET THE FILE NAME BACK
2140 PUMD R73,+R12 ! PUSH IT TO THE STACK
2150 JSB *ROMJSB ! SELECT THE ROM
2160 DEF MSPUR. ! PURGE THE FILE
2170 VAL MSROM# ! ROM TO SELECT
2180 POMD R10,-R6 ! RECOVER OUR BASE ADDRESS
2190 JMP CREATIT ! CONTINUE
2200 CALCRIN POMD R10,-R6 ! RECOVER BASE
2200 CRLCRTN POMD R10, -R6
2210 GTO RESTORE
2220 CREATIT PUMD R10, +R6
2230 ANM R17, *77 ! CLEAN UP THE ERROR FLAG
2240 CLB R20 ! AND THE DTHER DNE
2250 STBD R20, *ERRORS ! OUT IN RAM
2260 LDM R36, *COUNT ! GET THE REL ADDRESS OF ROUTINE
2270 ADM R36, R10 ! MAKE IT ABSOLUTE
2280 STM R36, R45 ! SET IT
2290 LDB R47, *236 ! LOAD A RTN OPCODE
2300 LDB R44, *318 ! LOAD A JSB OPCODE
2310 STMD R44, *IOTRFC ! TAKE THE HOOK
2320 LDMD R72, *LLDCOM ! SAVE LIST POINTERS
2330 PUMD R72, +R6 ! ON THE RTN STACK
2340 CLM R70 ! ZERO THE SELECT CODE
 LDMD R53,X10,FILENAME | GET THE FILE NAME
PUMD R53,+R12 | PUSH IT TO STACK
PUMD R40,+R12 | PUSH THE NUMBER OF RECORDS DESIRED
LEM R54,*377,56C,2C,0 | MAKE 256 BYTE RECORDS
PUMD R50,+R12 | PUSH IT TO THE STACK
JSB *ROMJSB | SELECT THE ROM
DEF MSCRE
       2570
       2580
       2590
       2600
      2610
      2620
                                                                                           DEF MSCRE.
                                                                                                                                                                                                                                                  ! CREATE THE FILE
       2630
                                                               VHL MSROM# | ROM # | RECOVER OUR BASE | CREATE?
       2540
      2650
2660
```

```
JCY SAVEX
                                                                               | JIF YES
 2670
                           PUMD R10,+R6 | SAVE OUR BASE

JSB X10,ASNPRT | ASSIGN THE BUFFER AND DO THE MSPRNT

POMD R10,-R6 | RECOVER OUR BASE

CMB R17,-300 | ANY ERRORS IN THE ASSIGN?
 2680
 2690
 2700
2710 CHB RTT, -300

2720 JNC PRINT | JIF NO

2730 SAVEX GTO RESTORE | BAIL DUT

2740 ASNPRT STMD R12, -TOS | MAKE SURE STACK LOOKS GOOD

CLM R50 | FIX UP FOR REAL 1
2710
             CLM R50 | FIX UP FOR KEHL |
LDB R57, *10C
PUMD R50, +R12 | PUSH IT TO THE STE
LDMD R53, X10, FILENAME | GET THE FILE NAME
 2760
 2770
                                                                                 I PUSH IT TO THE STACK
 2780
                           PUMD R53,+R12 | PUSH IT TO THE STACK
JSB =ROMJSB | SELECT THE ROM
DEF RSSIG. | RSSIGN THE BUFFER
                           PUMD R53,+R12

JSB =R0MJSB

DEF RSSIG.

VAL MSROM#

CMB R17,*300

JCY ASNRTN

CLM R40

LDB R47,*10C

PUMD R40,+R12

JSB =R0MJSB

DEF MSPRNT

VAL MSROM#

I ROM#

(FLOATING POINT 1)

PUSH IT TO THE STRCK

JSB =R0MJSB

DEF MSPRNT

VAL MSROM#

RTN

I ROM #

I ROM #

I ROM #

I ROM #

I DONE

I SELECT THE REL ADDRESS (
 2790
 2800
 2810
 2820
 2830
 2840
 2850
 2860
 2870
 2880
 2890
2900
3070 DEF PREOL. | DO THE END OF LINE PRINTING
3080 VAL MSROM# | ROM #
3090 POMD R10, -R6 | RECOVER DUR BASE
3100 POBD R16, -R6 | RESTORE CSTAT
3110 JSB X10, CLOSE | CLOSE THE FILE
3120 RESTORE LDMD R71, X10, SAVIDTFC | GET THE OLD HOOK
3130 STMD R71, -IDTRFC | RESTORE IT
3140 LDMD R70, X10, SAVSCTEM | GET THE OLD SELECT CODE
3150 STMD R70, -SCTEMP | RESTORE IT
3160 FINMSG JSB -CLEAR. | CLEAR THE CRT
3170 LDM R26, -MESAGE | LOAD THE ADDRESS OF THE MSG
3180 ADM R26, R10 | MAKE IT ABSOLUTE
3190 LDM R36, -4, O | LOAD THE LEN
3200 JSB -OUTSTR | DONE
  3220 MESAGE ASC "DONE"
                                                                                I DONE
  I NEED ANOTHER 1
  3240 CLOSE CLM R40
```

```
3250
                 LDB R47, *10C
                                              I FINISH THE 1
                PUMD R40,+R12 | PUSH TO STACK
LDM R46,=1,0 | LENGTH OF THE "*" STRING
PUMD R46,+R12 | PUSH IT TO STACK
LDM R45,+STAR | RDDRESS OF THE ASTERISK
3260
 3270
3280
3290
               BYT 0

ADM R45,R10

CLB R47

PUMD R45,+R12

PUMD R10,+R6

JSB *ROMJSB

DEF ASSIG.

VAL MSROM#

POMD R10,-R6

RTN

I RECOVER THE BASE

RTN

I RECOVER THE BASE

I DONE
                BYT 0
3300
                                              I NEED A THREE BYTE ADDRESS
3310
3320
3330
3340
3350
3360
3370
3380
3390
3410 COUNT BIN | FOR THE MATH
3420 CLB R40 | FOR THE MULTI
          CLB R40 | FOR THE MULTI-BYTE HUD

ADM R36,-4,0 | ADD SDME FOR THE HEADER

LDMD R45,-NXTDAT | GET THE PREVIOUS COUNT

ADM R45,R36 | ADD THE CURRENT LINE LEN

STMD R45,-NXTDAT | SAVE THE NEW COUNT

RTN | DONE
3430
3440
3450
3460
3470
3600 PUBD R3U, +R2+
3610 JMP SAVLOOP | LOOP TIL DUNE
3620 PRINT-IT JSB *ROMJSB | SELECT THE ROM
3630 DEF PRSTR, | PRINT THE STRING
3640 VAL MSROM# | ROM#
JSB *FXLEN. | MAKE SURE THE PROGRAM'S DEALLOCATED
JSB *CLEAR. | CLEAR THE SCREEN
LDM R26, *GETTING | GET ADDRESS OF MESSAGE
ADMD R26, *BINTAB | MAKE IT ABSOLUTE
LDM R36, *17,0 | LOAD THE LENGTH OF THE MESSAGE
JSB *DUTSTR | OUTPUT THE MESSAGE
JSB *DECUR2 | GET RID OF THE CUBECO
3750
3760
3770
3780
3790
                                         MOVE DOWN ONE LINE
3800 JSB -DNCURS
3810 BIN1 LDM R10,R4
3820
                 BIN
                                               ! GOOD FOR ADDRESS MATH
```

YOU WEEKEN	3/44/49/E/04/2000 - 04/2000	
4420	LDMD R24,R26	I GET THE LEN OF THE STRING READ
4430	JZR GOTBUF	I JIF NO CHARACTERS
4440 SW	AP POBD R32,-R26	I GET THE NEXT CHARACTER
4450	PUBD R32,+R30	I PUSH IT TO INPUT BUFFER
4460	DCM R24	I DECREMENT LEN COUNT
4470	DCM R24 JNZ SWAP	I JIF MORE TO DO
4480 G0		I COPY END OF BUFFER PTR
4490	SBM R36, = INPBUF	I MINUS THE START OF BUFFER
4500	STMD R36,X10,BUFLEN	
4510	LDB R24,=15 PUBD R24,+R30	I LOAD A CR CHARACTER
4520	PUBD R24,+R30	PUSH IT OUT FOR PARSER
4530	PUBD R25.+R6	I SAVE A O FLAG ON R6 FOR ERROR TRAP
4540	CMB R36, -81D	I DO WE NEED TO MOVE THE CURSOR DOWN?
4540 4550	JNC PARSIT	JIF NO
4560	JSB -DNCURS	I MOVE CURSOR DOWN A ROW
4570 PA		L COD LINEAR
	KOTI CED KID	! FOR LINEDR
4580	LDMD R20, -ASNTBL	1 SAVE ASSIGN BUFFER POINTER
4590	PUMD R20,+R6 LDMD R42,•LAVAIL	I DN THE R6 STACK
4600	LDMD R42, -LAVAIL	! SAVE SOME SYSTEM POINTERS
4610	PUMD R42,+R6	! ON THE RE STACK
4620	LDMD R42, =RTNSTK	! SAVE SOME MORE
4630	PUMD R42,+R6	I SAME PLACE
4640	LDMD R45, -LWAMEM	! SAVE SOME MORE
4650	PUMD R45,+R6	! AGRIN
4660	LDMD R45, -LAVAIL	
	CDMD R45, *CHYHIL	MOVE LWAMEM
4670	STMD R45, =LWAMEM JSB =RSETGO	UP TO LAVAIL
4680	JSB =RSETGO	I RESET EVERYTHING UP
4690	JSB =PARSER POMD R45,-R6	I TRY TO PARSE THE LINE
4700	POMD R45,-R6	START RECOVERING THINGS
4710	STMD R45, = LWAMEM	Committee Committee of the Committee of
4720	POMD R42,-R6	Ť.
4730	STMD R42, *RTNSTK	Ŷ.
4740	POMD R42,-R6	
4750	STMD R42, -LAVAIL	
	DOND RAZ, -LHVHIL	4
4760	POMD R20,-R6	A
4770	STMD R20, *ASNTBL	1
4780	LDB R16, =1	The second secon
4790	CMB R17, =300 JCY FIXIT	I ANY ERRORS?
4800	JCY FIXIT	I JIF YES
4810	BIN	I CONFIRM MATH MODE
4820	DCM R6	I THROW AWAY ERROR TRAP FLAG
4830	GTO OKGET	LOOP
4840 FI	DCM R6 GTO OKGET XIT POBD R36,-R6	RECOVER ERROR TRAP FLAG
4850	JNZ ERREXIT	
		JIF TWO ERRORS
4860	ICB R36	I SET FLAG
4870	PUBD R36,+R6	PUT IT BACK
4880	ANM R17, •77	CLEAR ERROR FLAGS
4890	CLM R40	CLEAR ERROR FLAGS
4900	STMD R40, = ERLIN#	I CLEAR ERROR FLAGS
4910	STBD R40, *ERRTYP	CLEAR ERROR FLAGS
4920	LDM R24, -INPBUF	GET ADDRESS OF BUFFER
4930	STM R24,R22	COPY
4940	ICM R24, R22	
		MOVE AHEAD TO THE FIRST CHARACTER
4950 MOV		GET THE FIRST CHARACTER
4960	PUBD R20,+R22	MOVE IT BACK ONE PLACE
4970	CMB R20,=40	I A BLANK ?
4980	JZR MOVE-1	JIF YES
4990	JSB *DIGIT	IS IT A DIGIT?

```
MAX LEN STRING VAR (0,1)
             ADDRESS OF FIRST BYTE OF $ VAR
```

```
5010 LDB R20, =41

5020 PUBD R20, -R22

5030 JSB =PRINT.

5040 BIN3 LDM R10, R4

5050 RTN
                                          | JIF YES
               JEN MOVE-1
5000
                                         ! ELSE LOAD A !
                                         I PUSH IT TO THE HOLE
                                          ! SET THE SCTEMP SELECT CODE
                                          I GET PC
               BIN | CALCULATE BASE IN CASE PARSER DESTROYED

SBM R10,*BIN3 | BINTAB

LDMD R36,X10,BUFLEN | GET LENGTH OF BUFFER

LDM R26,*INPBUF | GET THE START ADDRESS

JSB *DRV12. | PRINT THE LINE
5060
5070
5080
5090
                                        I GOT PARSE IT AS A COMMENT
5100
                GTD PARSIT
                LDM R10,R4 | GET CURRENT ADDRESS
BIN | FOR ADDRESS MATH
SBM R10,*ERREXIT | GET BPGM'S BASE ADDRESS
GTD FINMSG | GD DISPLAY 'DONE' MESSAGE
5110 ERREXIT LDM R10,R4
5120
5130
5140
5150 !***********************
        BYT 0,56
5160
                                          I FOR ADDRESS MATH
5170 REVISON, BIN
       LDM R43, -40D, 0
DEF DATE
BYT 0
                                       LEN OF STRING
5180
                                           ! ADDRESS AS TWO BYTE REL
5190
                                           I THERE'S THE THIRD BYTE
5200
              ADMD R45, = BINTAB
PUMD R43, +R12
                                          I NOW IT'S ABSOLUTE
5210
5220
                                          I PUSH TO RETURN STACK
5230 RTN | DONE
5240 ASC "81.202 .veR 2891 .oC drakcaP-ttelweH )c("
5250 DATE BSZ 0
5270 SAVING ASC "SAVE IN PROGRESS"
5280 GETTING ASC "GET IN PROGRESS"
5290 DONE ASC "DONE"
                ASC " *"
 5300
5310 STAR BSZ 0
5320 SAVIOTEC BSZ 7
5330 SAVSCTEM BSZ 10
5340 FILENAME BSZ 5
 5350 BUFADR BSZ 3
 5360 BUFLEN BSZ 2
                BSZ 300
 5370
5380 BUFFER BSZ 0
5400 ASNTBL DAD 100125
5410 ASSIG. DAD 65466
5420 BINTAB DAD 104070
5430 CALVRB DAD 100030
5440 CLEAR. DAD 14225
5450 CONBI3 DAD 4516
 5460 DECURZ DAD 13467
5470 DIGIT DAD 21710
5480 DNCURS DAD 13751
5490 DRV12. DAD 6722
5500 ERLIN# DAD 100114
5510 ERROR DAD 10223
5520 ERROR+ DAD 10220
5530 ERRORS DAD 100123
5540 ERRTYP DAD 100124
5550 FLDCOM DAD 100053
5560 FXLEN DAD 31001
5570 G012N DRD 24707
```

5580	INPBUF	DAD	100236
5590	IOTRFC	DAD	103643
	LAVAIL		100025
5610	LLDCOM	DAD	100050
	LSSET		6445
5630	LWAMEM		
	MSCRE.	DAD	65176
5650	MSPRNT		
	MSPUR.		64604
5670	MSROM#		
5680	NXTDAT	DAD	101645
	NXTMEM	DAD	100022
5700	ONEI	DAD	56736
	OUTSTR	DAD	14020
5720	PARSER	DAD	20000
5730	PREOL.	DAD	70464
5740	PRINT.	DAD	71332
5750	PRSTR.	DAD	66662
5760	PTR2 -	DAD	177715
5770	PTR2+	DAD	177716
5780	RDSTR.	DAD	67314
5790	ROMJSB	DAD	6223
5800	RSETGO	DAD	5700
5810	RTNSTK	DAD	100033
5820	SCTEMP	DAD	101721
5830	ST240+	DAD	21067
5840	STREX+	DAD	23721
5850	TOS	DAD	101744
5860	FIN	FIN	

#### REFERENCE MATERIAL

### 8.1 Overview

This section consists of:

- · An alphabetical listing of the global file.
- System operation and routines.
- · Parsing flow diagrams.
- · General hook flowcharts for the following:

CHIDLE

DCIDLE

IOSP

IOTRFC

IRQ20

KYIDLE

PRSIDL

RMIDLE

SPARØ and SPAR1

- System run time table tokens and attributes.
- Error messages.
- The assembler instruction set.
- An assembler instruction coding table.
- A keycode table.
- Some programming hints.

#### 8.2 The Global File

The global file as it appears on the disc is listed here. It gives the permanent addresses in memory of many of the system routines. The global file also contains locations of system pointers, buffers, variables, and constants which may be referenced in a binary program.

Although it is usually more convenient, it is not necessary to use the file GLOBAL as a label table. You may create your own on a disc, or you may specify the addresses of the system routines called in a binary program by adding them to the label table within the program.

	Name	Address	Description	
1000	******	*******	*************	**
1010				*
		GLOBAL FILE FOR USE W	ITH THE ASSEMBLER ROM.	*
1030	1 *			*
1040	(c) 1	982 Hewlett-Packard Co		*
1050	<b>*</b>			*
1060	*			*
1070	!******	********	*********	**
1080	! *NOTE: E	Beware of looking up a	routine in the global file and using	*
1090	I # 1 t with	rout also looking up th	e documentation. This is especially	*
1100	*true if	the routine has an en	try point address between 80000 and	*
1110	1 * 7 7 7 7 7 7	as it may need to be c	alled through ROMJSB.	*
			***********	**
1130		GLO	THE COMMENT OF STREET,	
1140	ABS5	DAD 54525 DAD 177515	I ABS FUNCTION RUNTIME CODE	
1150	ACTB-3	DAD 177515	I I/O MODULE ADDRESSES	
1160	HCIB-P	DMD 177512	1 I/O MODULE ADDRESSES	
	ACTBAS	DAD 177520	1 I/O MODULE ADDRESSES	
	ACTBS+	DAD 177521	1 I/O MODULE ADDRESSES	
		DAD 103560	I ACTIVE MSUS FOR MASS STORAGE ROM	Q
1200	ADD10	DAD 53030	I ADD TWO REAL NUMBERS IN R40 AND R5	
		DAD 52745	! ADD 2 REAL OR INTEGER NUMBERS OFF	SIHCK
		DAD 103416	PLOTTER ROM STOLEN RAM BASE ADDRES	5
		DAD 21656	CHECK TO SEE IF RZO IS ASCII A-Z D	m a-z
		DAD 12466	I FORCE ALPHA ALL MODE I FORCE ALPHA DR ALPHA ALL IF NOT GR	apuai i
		DAD 12542	FORCE ALPHA NORMAL	DEDDFF
		DAD 12413	ADVANCED PROG ROM STOLEN RAM BASE	anno
	APRBAS	DAD 103420	# OF BYTES IN ALPHA (4K OR 16K)	HIDDIN,
1280	ASI ZE	DAD 104744 DAD 103426	ASSEMBLER ROM	
	ASMBAS		24 BYTES ASSIGN FILES	
	ASNTBL	DAD 100125 DAD 65466	ASSIGN A DISC BUFFER TO A FILE	
		DAD 77157	ATNZ FUNCTION	
	ATN2. AUTO#	DAD 100103	! AUTO LINE # LAST VAL	
1330	HOTO#	DID 100103	THOSE CITE II WIND THE	

Name	Address	Description
Name  1840 CSIZE. 1850 CURS 1860 CURSON 1870 CVNUM 1880 DALLED 1890 DALLED 1900 DATE 1910 DATE. 1920 DCIDLE 1930 DCLIN# 1940 DCSLOP 1950 DEFA+. 1970 DEFA 1970 D	Address  DAD 66570 DAD 14030 DAD 105347 DAD 72401 DAD 101104 DAD 47123 DAD 101133 DAD 32073 DAD 104035 DAD 34607 DAD 34607 DAD 35132 DAD 13467 DAD 61576 DAD 61576 DAD 61576 DAD 61576 DAD 100152 DAD 100477 DAD 62257 DAD 54736 DAD 104044 DAD 21710 DAD 100542 DAD 101136 DAD 101136 DAD 52441 DAD 21710 DAD 101136 DAD 52441 DAD 52435 DAD 13607 DAD 52441 DAD 52435 DAD 13607 DAD 13751 DAD 13607 DAD 13751 DAD 13607 DAD 13751 DAD 13607 DAD 13751 DAD 13751 DAD 13607 DAD 13751 DAD 13607 DAD 13751 DAD 13607 DAD 13751 DAD 13751 DAD 13607 DAD 32231 DAD 100160 DAD 6722 DAD 32161 DAD 32231 DAD 14525 DAD 100017 DAD 3564 DAD 66623 DAD 100542	Description  'CSIZE' STATEMENT TURN CURSOR ON CURSOR ON FLAG FORMAT A REAL NUMBER FOR OUTPUT DEALLOCATED FLAG DE-ALLOCATED FLAG DE-ALLOCATE THE BASIC PROGRAM JULIAN DAY YEAR DATE FUNCTION DECOMPILE A BASIC PROGRAM LINE NUMBER REVERSE A STRING FROM EXTENDED MEMORY TURN CURSOR OFF TURN MATH DEFAULTS ON TURN MATH DEFAULTS OFF DEFAULT ERROR FLAG DEFAULT ERROR FLAG DEFAULT ERROR FLAG DISCREDIS TO ENGRESS CONVERSION DIRECTION FLAG FOR DISC READ/WRITE DIGITIZE HOOK FOR CRT DIGITIZING SEE IF R2O CONTAINS A DIGIT (ASCII CODE) DISPLAY BUFFER SET SELECT CODE TO CRT IS DEVICE I BYTE DISPLAY LINE LENGTH DISP BUFFER PTR DIVIDE 2 REAL NUMBERS IN R40 AND R50 DIVIDE 2 REAL OR INTEGER NUMBERS ON STAK DEMAND CARRIAGE RTM, BANG (1), OR @ SIGN MOVE CURSOR DOWN IN ALPHA MEMORY DRAW A LINE ON THE CRT DEG/RAD/GRAD FLAG OUTPUT VECTOR ROUTINE EDITOR MODE (INSERT/REPLACE) EXTENDED MEMORY MOVON
2300 ERRBUF	DAD 100476	ERROR BUFFER (44 BYTES)
2310 ERROM#	DAD 100121	ROM# OF LAST ERROR
2320 ERROR	DAD 10224	! REPORT ERROR ROUTINE
2330 ERROR+	DAD 10220	! REPORT ERROR AND THROW AWAY 1 RTN ADDR.
2340 ERRORS	DAD 100123	! RUN TIME ERRORS
2350 ERROM	DAD 100120	! ROM# OF ERROR

Name	Address	Description
2360 ERRSC 2370 ERRTYP 2380 ERRTEMP 2390 EXEC 2400 EXP5 2410 EXSTAT 2420 EXTFIL 2430 FETSVA 2450 FETSVA 2450 FETSVA 2450 FETSVA 2450 FLIP 2480 FLIP 2500 FNAM+5 2510 FNAM+5 2520 FNAME. 2510 FNAME. 2510 FNAME. 2510 FNAME. 2520 FNAME. 2530 FWEURR 2530 FWEURR 2530 FWEURR 2550 FWEURR 2560 FWEURR 2580 FWEURR 2680 GO12 2670 GO12 2680 GO12 2680 GO12 2680 GO12 2680 GCHR. 2690 GCHR. 2690 GEMINI	DAD 101141 DAD 100124 DAD 100124 DAD 104200 DAD 72 DAD 53174 DAD 177425 DAD 110010 DAD 11565 DAD 50333 DAD 45505 DAD 45505 DAD 100053 DAD 100053 DAD 103510 DAD 103510 DAD 103510 DAD 32355 DAD 27034 DAD 54665 DAD 100006 DAD 100006 DAD 100000 DAD 24543 DAD 100000 DAD 24543 DAD 100000 DAD 24726 DAD 24726 DAD 24772 DAD 24771 DAD 24771 DAD 24772 DAD 24771 DAD 24771 DAD 24772	Description    ERROR SELECT CODE     ERROR TYPE     12 BYTES TEMP     BEGINNING OF MAIN EXEC LOOP     EXP FUNCTION (8^X)     EXTENDED IO STATUS     EXTENDED FILE TYPE TABLE     FAST BACKSPACE (SHIFTED BACKSPACE KEY)     FIND BINARY PROGRAM (BY BPGM #)     FETCH ARRAY VARIABLE ADDRESS     FETCH SIMPLE VARIABLE ADDRESS     TBYTE TAPE, TEMP     FIRST LINE DECOMPILE     TOGGLE THE KEYBOARD 'FLIP' STATUS     FILE NAME 1ST HALF     FILE NAME 2ND HALF     FIND A BASIC PROGRAM LINE IN MEMORY     PARSE AN ARRAY REFERENCE     FRACTIONAL PART FUNCTION     FRAME THE CAT     FAM USER BIN PROG     PTR TO CURRENT PGM     FWA PROGRAM AREA     FWA USER BIN PROG     PTR TO CURRENT PGM     FWA USER AREA     GET STRING & NUMERIC WITH DPTIONALS     TOGGLE BETWEEN GRAPH AND ALPHA     GET O, 1, OR 2 NUMERIC VALUES     GET O OR 1 NUMERIC VALUES     GET O OR 2 NUMERIC VALUES     GET 1, 2, OR 4 NUMERIC VALUES     GET 1, 2, OR 4 NUMERIC VALUES     GET 1 CHARACTER AT PARSE TIME     CLEAR THE GRAPHICS CRT DISPLAY     GEMINI FLAG     COMPARE TWO NUMBERS FOR >=     GET H CLOSE PARENTHESIS     PARSE ONE NUMBERS     PARSE TWO NUMBERS     PARSE FOUR NUMBERS     PARSE ONE NUMBERS     PARSE TWO NUMBERS     PARSE FOUR NUMBERS     PARSE TWO NUMBERS     PARSE TWO NUMBERS     PARSE FOUR NUMBERS     PARSE FOUR NUMBERS     PARSE TOW NUMBERS     PARSE FOUR NUMBERS     PARSE TOW NUMBERS     PARSE TOW NUMBERS     PARSE FOUR NUMBERS     PARSE TOW
2720 GEQ\$, 2730 GEG. 2740 GET) 2750 GET1N 2760 GET2N 2770 GET4N 2780 GETCMA 2790 GETPAP 2800 GETPAR 2810 GINTDS 2820 GINTEN 2830 GLINE 2840 GLOAD 2850 GNAM 2860 GNAM+5 2870 GOTOSU	DAD 3667 DAD 62734 DAD 23450 DAD 24557 DAD 24630 DAD 24635 DAD 23477 DAD 24740 DAD 24562 DAD 177401 DAD 177400 DAD 104740 DAD 72510 DAD 103515 DAD 30317	COMPARE FOR GREATER THAN OR EQUAL TO COMPARE TWO NUMBERS FOR >= GET A CLOSE PARENTHESIS PARSE ONE NUMBER PARSE TWO NUMBERS PARSE FOUR NUMBERIC PARAMETERS DEMAND A COMMA AT PARSE TIME GET SOME OPTIONAL PARAMETERS GET A SET NUMBER OF NUMERIC PARAMETERS GLOBAL INTERRUPT DISABLE NUMBER OF DOTS ON A LINE OF GRAPH SCREEN GLOBAL STATEMENT FOR MASS STORAGE COPY, RENAME, ETC. PARSE A GOTO/GOSUB LINE NUMBER OR LABEL

Name	Address	Description
Name  2880 GR\$, 2890 GR. 2900 GRAD. 2910 GRAPH. 2920 GRAPH. 2920 GRAPH. 2940 GSIZE 2950 GSTOR, 2960 HLFLIN 2970 HMCURS 2980 HORN 2990 ICOS 3000 IDRAW, 3010 IMERR 3020 IMOVE. 3030 INF10 3050 INIT. 3050 INPBUF 3060 INPBUF 3060 INPTOS 3110 INTEGR 3120 INTEGR 3150 INTEGR 3250 INTEGR 3250 INTEGR 3250 INTEGR 3250 INTEGR 3250 INTEGR	Address  DAD 03614 DAD 62705 DAD 62274 DAD 12560 DAD 12560 DAD 12574 DAD 104742 DAD 72711 DAD 14110 DAD 13661 DAD 10400 DAD 77254 DAD 64706 DAD 103724 DAD 64706 DAD 103724 DAD 100233 DAD 100233 DAD 100236 DAD 100167 DAD 100167 DAD 100167 DAD 101143 DAD 100204 DAD 101143 DAD 101143 DAD 101143 DAD 101143 DAD 101143 DAD 101144 DAD 101140 DAD 177422 DAD 177421 DAD 103643 DAD 103643 DAD 103643 DAD 103643 DAD 103643 DAD 103643	Description    COMPARE STRINGS FOR GREATER THAN   COMPARE NUMBERS FOR GREATER THAN   SET THE COMPUTER TO GRAD MODE   FORCE GRAPH ALL MODE   SWITCH TO GRAPH UNLESS IN ALPHA ALL   FORCE GRAPH NORMAL MODE   H OF BYTES IN GRAPH SCREEN (12K OR 16K)   GSTORE' STATEMENT   DISP STRING WITHOUT CR AND LF   HOME CURSOR ON CURRENT CRT PAGE   LOWER LEVEL 'BEEP' ENTRY POINT   ARC COSINE FUNCTION   I'DDRAW' STATEMENT   IMAGE ERROR INTERCEPT RAM HOOK   I'MOVE' STATEMENT   READ ONE CHARACTER IN FROM CRT MEMORY   INFINITY FUNCTION (RETURNS BIGGEST #)   I'NIT' KEY EXECUTION   3 PERMANENT BYTES IN FRONT OF INPBUF PARSER INPUT BUFFER   INPUT COMPLETION ADDRESS   R10 SAVE DURING INPUT   INPUT TOP OF STAK   INPUT LINE POINTER   INPUT RUNTIME ROUTINE   INT FUNCTION   INTEGER DIVISION (\(\)) RUNTIME   GET AN INTEGER AT PARSE TIME   MULTIPLY TWO BINARY NUMBERS   CONVERT A TAGGED INTEGER TO A REAL IND DATA   INPUT LIVE OF THE CONVERT
3270 IPS 3280 IPLOT. 3290 IRG20	DAD 54770 DAD 64660 DAD 103742	P FUNCTION RUNTIME CODE  'IPLOT' STATEMENT  I /O INTERRUPT RAM HOOK
3300 IRQ20+ 3310 IRQPAD 3320 IRQRTN 3330 ISIN 3340 ITAN 3350 KEYCNT	DAD 103751 DAD 103757 DAD 103760 DAD 77244 DAD 77264 DAD 100153	I/O INTERRUPT RAM HOOK   I/O INTERRUPT RAM HOOK   I/O INTERRUPT RAM HOOK   ARC SINE FUNCTION   ARC TANGENT FUNCTION   KEYBOARD REPEAT COUNTER
3360 KEYCOD 3370 KEYHIT 3380 KEYLA. 3390 KEYSTS	DAD 177403 DAD 101142 DAD 13360 DAD 177402	KEYBOARD CODE AND EDJOB I/O ADDRESS   KEYBOARD ASCII   KEY LABEL RUNTIME ROUTINE   KEYBOARD STATUS I/O ROUTINE

Name	Address	Description
Name  3400 KEYTAB 3410 KRPET1 3420 KRPET2 3430 KYIDLE 3440 LASTIN 3450 LASTIN 3450 LEGCA2 3450 LEGEN2 3540 LEGEND 3520 LEGEN 3550 LEGEN 3550 LEGEN 3550 LEGEN 3550 LINEL 3650 LT\$ 365	Address  DAD 102016 DAD 100154 DAD 100155 DAD 103677 DAD 67262 DAD 100475 DAD 100025 DAD 101525 DAD 101525 DAD 101405 DAD 101265 DAD 101145 DAD 3656 DAD 62662 DAD 101714 DAD 66336 DAD 6352 DAD 104231 DAD 104231 DAD 104233 DAD 52346 DAD 104233 DAD 52346 DAD 104750 DAD 52515 DAD 104233 DAD 13757 DAD 104537 DAD 104537 DAD 104537 DAD 104537 DAD 104537 DAD 104537 DAD 1056144 DAD 103424 DAD 56125 DAD 107424 DAD 56125 DAD 177424 DAD 56125 DAD 13771 DAD 57172 DAD 64634 DAD 57232 DAD 53357 DAD 53357 DAD 53357 DAD 53357 DAD 53357 DAD 53357	BASE ADDR KEY TABL     MAJOR KYBD REPEAT     MINOR KYBD REPEAT     KEYBOARD INTERCEPT     'LABEL' STATEMENT     END OF INPUT BUFFER     LAST AYAIL WD IN PGM AREA     'LDIR' STATEMENT     CALC KEYLABELS (BTM ROW)     CALC KEYLABELS (BTM ROW)     RUN KEYLABLES (BTM ROW)     RUN KEYLABLES (BTM ROW)     COMPARE STRINGS FOR LESS THAN OR EQUAL     COMPARE STRINGS FOR LESS THAN OR EQUAL     DEVICE LINE LENGTH     'LINE TYPE' STATEMENT     'LIST' STATEMENT     'LIST' STATEMENT     LAST LINE DECOMPILE     PGSIZE - DNE LINE     PGSIZE - TWO LINES     NATURAL LOGARITHM FUNCTION     LINE TYPE POINTER TABLE     BASE 10 LOGARITHM FUNCTION     LINE TYPE POINTER TABLE     COMPARE STRINGS FOR LESS THAN     COMPARE STRINGS FOR LESS THAN     COMPARE NUMBERS FOR
3830 MSBASE 3840 MSCRE. 3850 MSHIGH 3860 MSLOW 3870 MSPRNT 3880 MSPUR. 3890 MSREN. 3900 MSTIME 3910 NRRRE+	DAD 103773 DAD 56221 DAD 64604 DAD 64724 DAD 104002	MASS STORAGE ROM STOLEN RAM BASE ADDRESS   CREATE RUNTIME CODE   MS HIGH LEVEL HOOK   MS LOW LEVEL HOOK   PART OF PRINT# RUNTIME CODE   PURGE RUNTIME CODE   'RENAME' STATEMENT   MS TIMEOUT HOOK   SCAN AND PARSE A NUMERIC ARRAY REFERENCE

Name	Address	Description
Name  3920 NARREF 3930 NUMCON 3940 NUMVA+ 3950 NUMVAL 3960 NXTDAT 3970 NXTMEM 3980 DNEB 4000 DNEI 4010 DNER 4020 DNEROI 4030 DNELAG 4050 OPTBAS 4060 OUTCHR 4080 OUTSTR 4080 OUTSTR 4090 P.FLAG 4110 P.FLAG 4120 P.FLAG 4130 P.FLAG 4150 P.FLAG 4150 P.FLAG 4170 P.FLAG 4170 P.FLAG 4180 P.FLAG 4210 P.FLAG 4210 P.FLAG 4210 P.FLAG 4210 P.FLAG 4220 P.FLAG 4230 P	Address  DAD 23465 DAD 23551 DAD 22403 DAD 22406 DAD 101645 DAD 100022 DAD 100036 DAD 12153 DAD 56736 DAD 56777 DAD 57035 DAD 56673 DAD 100065 DAD 100065 DAD 10175 DAD 14130 DAD 14143 DAD 14020 DAD 101710 EQU 6 DAD 12756 DAD 101710 EQU 6 DAD 12756 DAD 1303 DAD 1303 DAD 1303 DAD 104535 DAD 107627 DAD 54374 DAD 6344 DAD 64652 DAD 1077423 DAD 70730 DAD 70730 DAD 70730 DAD 70730 DAD 70730 DAD 73023 DAD 103550	PARSE A NUMERIC ARRAY REFERENCE PARSE A NUMERIC CONSTANT SCAN AND PARSE A NUMERIC EXPRESSION PARSE A NUMERIC EXPRESSION NEXT DATA ADDRESS FOR DISK READ/WRITE NEXT BYTE AVAILABLE MEMBRY NEXT AVAILABLE GOSUB/RTN GET 1 NUMBER OFF STACK AS SIGNED BINARY GET 1 NUMBER OFF STACK AS FIGGED INTEGER GET 1 NUMBER OFF STACK AS FRAL OR INTEGER GET 1 NUMBER OFF STACK AS READ OR INTEGER GET 1 NUMBER OFF STACK AS READ OR INTEGER GET 1 NUMBER OFF STACK AS UNSIGNED BIN. ON GOSUB FLAG 2 BYTE PERMANENT OPTION BASE OUTPUT A BYTE TO THE CRT OUTPUT A STRING TO CRT OUTPUT A STRING TO CRT INDIRECT BUFFER POINTERS INDIRECT BUFFER POINTER INDIRECT BUFFER POINTER OFFSET INTO BASIC PCB TO GET TYPE BYTE PAGESIZE RUNTIME CODE PAGESIZE 16 PAGESIZE 16 PAGESIZE 24 SYSTEM PARSER (PEN #) * 3 FOR INDEXING # OF BYTES / PAGE PI FUNCTION RUNTIME CODE PLOTTER HOOK 'PLIST' STATEMENT 'PLOT' STATEMENT 'PLOT' STATEMENT 'PLOT' STATEMENT 'PLOTT STATEMENT 'PLOTT STATEMENT 'PLOTTER ON/OFF FLAG PRINT# STRING ARRAY TO DISC FILE PRINT# STRING TO PRINTER IS DEVICE PRINT LINE RUNTIME CODE  1 BYTE PRINTER LINE LENGTH PRINT# A SUMBER TO A DATA FILE PRINT# A NUMBER TO A DATA FILE PRINT# A STRING TO B DATA FILE PRINT# B STRIN
4290 PRDVF+ 4300 PREDL. 4310 PRINT. 4320 PRLINE	DAD 103550 DAD 70464 DAD 71332 DAD 71641	SPECIAL CHARACTER FLAG FOR LIST TIME PRINT# END OF LINE (DUMP BUFFER) SET SELECT CODE TO PRINTER IS DEVICE PRINT LINE RUNTIME CODE
4340 PRNTR. 4350 PRNUM. 4360 PRSIDL 4370 PRSTR.	DAD 75631 DAD 67220 DAD 103733 DAD 66662	! I BYIL PRINTER LINE LENGTH ! PRINTER IS STATEMENT ! PRINT# A NUMBER TO A DATA FILE ! PARSER RAM HOOK ! PRINT# A STRING TO A DATA FILE
4380 PRTBUF 4390 PRTPTR 4400 PS.C. 4410 PTR1 4420 PTR1+ 4430 PTR1-	DAD 107454 DAD 100062 DAD 100222 DAD 177710 DAD 177712 DAD 177711	PRINT BUFFER PRINT BUFFER PTR PRINTER SELECT CODE I/O ADDRESSES FOR EMC POINTERS I/O ADDRESSES FOR EMC POINTERS I/O ADDRESSES FOR EMC POINTERS

Name	Address	Description
Name  4440 PTR1-+ 4450 PTR2 + 4460 PTR2-+ 4470 PTR2-+ 4480 PTR2-+ 4490 R60+12 4500 R60+14 4520 R60+2 4510 R60+6 4550 R60K 4560 RAD, 4570 RAD10 4580 RAID+1 4590 RAID+1 4590 RAID+2 4600 RDARR* 4610 RDARR* 4610 RDARR. 4620 RDNUM. 4630 RECBUF 4660 REFNUM 4670 RELMEM 4670 RELMEM 4690 RESET. 4700 RESET. 4700 RESET. 4700 RESET. 4700 RESET. 4730 RMEM	Address  DAD 177713 DAD 177714 DAD 177716 DAD 177715 DAD 177717 DAD 60010 DAD 60012 DAD 60004 DAD 60006 DAD 60006 DAD 60000 DAD 60000 DAD 62267 DAD 54472 DAD 103307 DAD 103310 DAD 70312 DAD 70106 DAD 67503	Description  I/O ADDRESSES FOR EMC POINTERS I/O ADDRESSES FOR ROMS I/O ADDRESS FOR ROMS INTO ROUTINES I/O ADDRESS FOR ROMS FOR MADDRESS I/O A BANK SELECTABLE ROM I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O CURSOR RIGHT ON CURRENT SCREEN I/O A CHECKSUM ON BK OF MEMORY I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O CURSOR RIGHT ON CURRENT SCREEN I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O CURSOR RIGHT ON CURRENT SCREEN I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O CURSOR RIGHT ON CURRENT SCREEN I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O CURSOR RIGHT ON CURRENT SCREEN I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O ADDRESS I/O A CHECKSUM ON BK OF MEMORY I/O ADDRE
4690 RESET. 4700 RESMEM 4710 RESULT 4720 RETRHI 4730 RMEM 4740 RMIDLE 4750 RND10 4760 RND1Z. 4770 RDMEND	DAD 5407 DAD 31741 DAD 100070 DAD 13234 DAD 105343 DAD 103706 DAD 53741 DAD 55713 DAD 104145	RESET KEY RUNTIME CODE RESERVE SOME TEMPORARY MEMORY LAST CALCULATOR MODE RESULT WAIT FOR RETRACE HIGH FROM CRT RESERVED MEMORY COUNT EXEC LOOP RAM HOOK RND FUNCTION (GET A RANDOM NUMBER) RANDOMIZE COMMAND END OF ROM TABLE ENTRIES
4780 ROMFL 4790 ROMINI 4800 ROMISH 4810 ROMLST 4820 ROMRTN 4830 ROMTRB 4840 RPLOT. 4850 RSELEC 4860 RSTREG 4870 RSUM8K 4880 RTCUR. 4890 RTCUR. 4890 RTCURS 4900 RTNSTK 4910 RULITE 4920 S10 4930 SRD1	DAD 104085 DAD 6055 DAD 6223 DAD 104143 DAD 6207 DAD 104105 DAD 64666 DAD 177430 DAD 22346 DAD 37670 DAD 13651 DAD 100033 DAD 177704 DAD 103367 DAD 103367	ROM FLAG FOR INITIALIZATION ROUTINES CALL BPGM'S AND ROM'S INIT ROUTINES JSB TO A BANK SELECTABLE ROM- LAST ENTRY IN ROM TABLE RE-SELECT ROM O AND RETURN BASE OF ROM TABLE 'RPLOT' STATEMENT BANK SELECTABLE ROM SELECTION ADDRESS RESTORE REGISTERS DO A CHECKSUM DN BK OF MEMORY MOVE CURSOR RIGHT ON CURRENT SCREEN MOVE CURSOR RIGHT IN ALPHA MEMORY TOP OF GOSUB RETURN STAK RUN LIGHT I/O ADDRESS FOR SAVING R10-11 DURING INTERRUPT SVC
4940 SAVER6 4950 SAVRO	DAD 104066 DAD 103200	DISK BAIL OUT STACK POINTER FOR ERRORS SYSTEM MONITOR REGISTER SAVE AREA

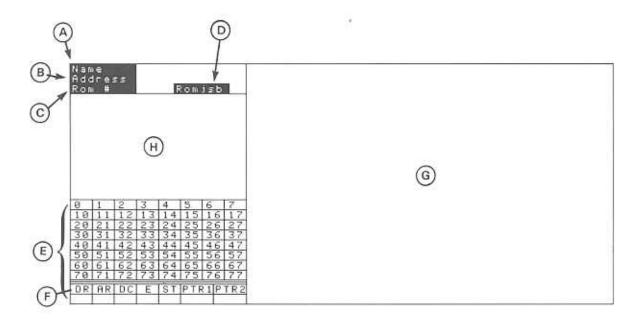
Name	Address	Description
Name  4960 SAVR10 4970 SAVREG 4980 SC10+1 4990 SCAN 5000 SCAN+ 5010 SCRDN 5030 SCRUP 5040 SCTEMP 5050 SEC10 5060 SEMIC\$ 5070 SEMIC. 5080 SEQND+ 5100 SERPOL 5110 SET240 5120 SGN5 5130 SIN10 5140 SKYTXT 5150 SPAR1 5170 STANN 5180 STREXP 5300 STREXP 5310 STREXP	Address  DAD 104063 DAD 22310 DAD 177540 DAD 177540 DAD 211105 DAD 5601 DAD 13671 DAD 13671 DAD 13671 DAD 54260 DAD 72155 DAD 54260 DAD 72155 DAD 54260 DAD 72155 DAD 54260 DAD 72274 DAD 30422 DAD 177425 DAD 10422 DAD 104011 DAD 104022 DAD 103303 DAD 53237 DAD 103303 DAD 103303 DAD 103303 DAD 103303 DAD 103527 DAD 103715 DAD 103715 DAD 24201 DAD 24056 DAD 103715 DAD 24201 DAD 23724 DAD 23724 DAD 23724 DAD 23724 DAD 52734 DAD 52734 DAD 52734 DAD 52724 DAD 52724 DAD 52724 DAD 52724 DAD 52724 DAD 52724 DAD 54363 DAD 101741 DAD 52734 DAD 52724	R10, SAVE FOR PARSE ERRORS   SAVE REGISTERS ON RG   I/O CARD STUFF   GET NEXT TOKEN TO R14 AT PARSE TIME   GCHAR AND SCAN   'SCRATCH' RUNTIME CODE   SCROLL DOWN THE CRT   SCROLL UP THE CRT   S.C. TEMP STORE   SECANT RUNTIME CODE   PRINT STRING;   PRINT NUMBER;   PARSE A LINE NUMBER   MY LISTEN ADDRESS   SET THE IMMEDIATE BREAK BITS IN R17   SGN FUNCTION   SIN FUNCTION   SIN FUNCTION   SAME INTERRUPT RAM HOOK (SYS MONITOR)   SPARE INTERRUPT RAM HOOK (UNUSED)   DISC VOLUME NAME   SYSTEM MONITOR SAVE PTR1 AREA   SYSTEM MONITOR SAVE PTR2 AREA   SYSTEM MONITOR SAVE PTR2 AREA   SOURRE ROOT FUNCTION   CLEAR R16 AND SET240   R6 STACK 500 OCTAL BYTES (320 DECIMAL)   STARDARD BEEP   STORE STRING ROUTINE   STARNGE PARAMETER TYPES INTERCEPT HOOK   PARSE A STRING CONSTANT   SCAN AND PARSE A STRING EXPRESSION   PARSE A STRING EXPRESSION   PARSE A STRING VARIABLE   STRENGE PARAMETER TYPES INTERCEPT HOOK   PARSE A STRING VARIABLE   STATEMENT SIZE PLACE HOLDER POINTER   SUBTRACT TWO REAL NUMBERS IN R40 AND R50   SUBTRACT Z REAL OR INTEGERS ON STACK   SERVICE WORD   SOS CARD ROM DISABLE ADDRESS   TANGENT FUNCTION   TIME OF DAY   TIME OF DAY
5400 TWDB 5410 TWDR 5420 TWDRDI 5430 UNBAS1 5440 UNBAS2 5450 UNEQ\$, 5460 UNEQ, 5470 UNQUOT	DAD 56760 DAD 57020 DAD 57050 DAD 103430 DAD 103434 DAD 3603 DAD 62632 DAD 24366	GET TWO BINARY NUMBERS OFF STACK GET 2 REAL NUMBERS OFF R12 STACK GET 2 REAL OR INTEGERS OFF R12 STACK UNUSED ROM STOLEN RAM BASE ADDRESS UNUSED ROM STOLEN RAM BASE ADDRESS COMPARE STRINGS FOR UNEQUAL COMPARE NUMBERS FOR UNEQUAL PARSE AN UNQUOTED STRING

# 8.3 System Operation and Routines

This section provides documentation for certain areas of system operation. It also shows the input conditions required and the outputs produced by selected system routines. The names and addresses of the system routines detailed here are also on the disc.

The system routines are arranged in alphabetical order. Their area of primary use is noted. Because a routine is listed under a certain application does not limit its use to that area. For example, many utility routines may also be used during run time operations.

The format of the individual system routines is shown here:

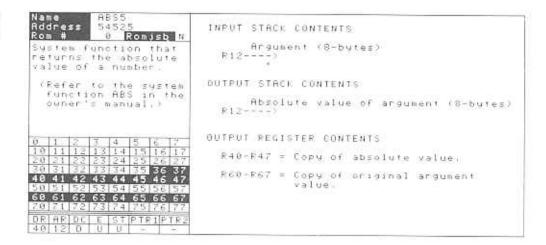


- A. Name: Name of the routine (from the global file).
- B. Address: Permanent octal address of the routine in computer memory.
- C. ROM#: The ROM that must be selected if this routine needs to be called through ROMJSB.
- D. ROM #: The "Y" or "N" entry indicates if this routine needs to to be called through ROMJSB.
- E. Registers: Shaded areas indicate registers used by this routine.
- F. DR, AR, DC, E, ST, PTR1, PTR2: Entries in these boxes indicate exit conditions of this routine. The following symbols are used:

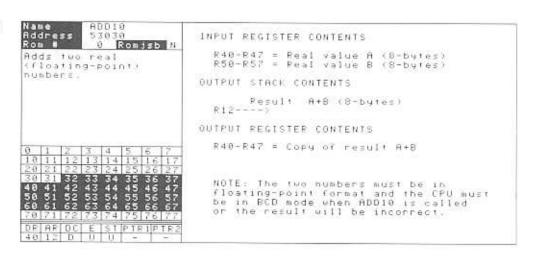
Symbol	Meaning
-	Unchanged.
U	Unknown.
*	Refer to the description (G).

- ${\tt G.}$  Conditions: When applicable, shows input and output stack contents, and output register contents.
- H. Description: Contains description of routine.

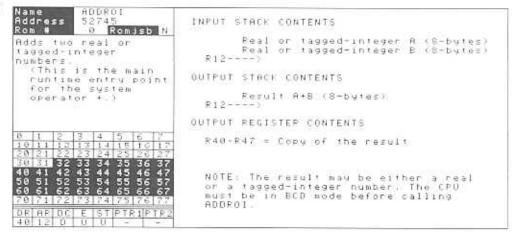
#### ABS5 MATH

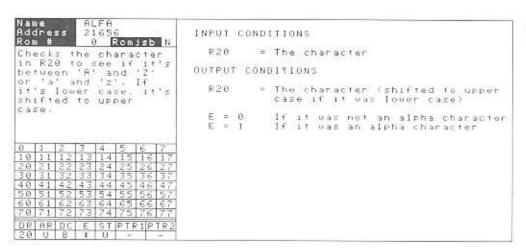


#### ADD10 MATH



# ADDROI MATH

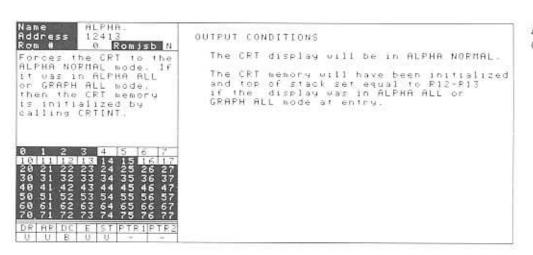




ALFA PARSE

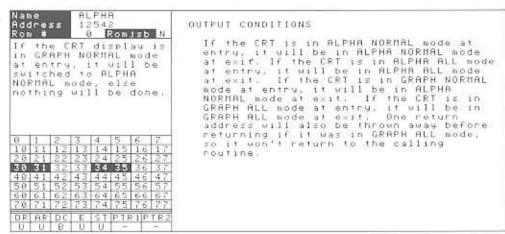
Name Address   ALFAL, Address   12466 Rom # @ Romjsb N Forces ALPHA ALL mode on the CRT.	OUTPUT CONDITIONS  The CRT will be in ALPHA ALL mode. The actual code for ALFAL. is:
	ALFAL. BIN LOBD R30,=CRTSTS   GET CRT STATUS LLB R30   GRAPH/GRAPHALLS ELB R30   ERB R30   JCY ALFAL1   JIF TES JNG ARTN   JIF ALPHA ALL
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTR1PTR2 U U B - U -	ALFALI JSB = CRTHPO

ALFAL. CRT

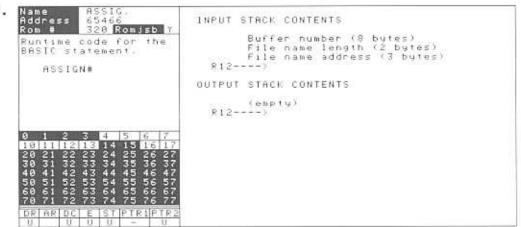


ALPHA CRT

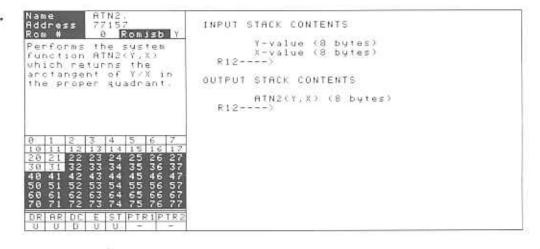
#### ALPHA. CRT

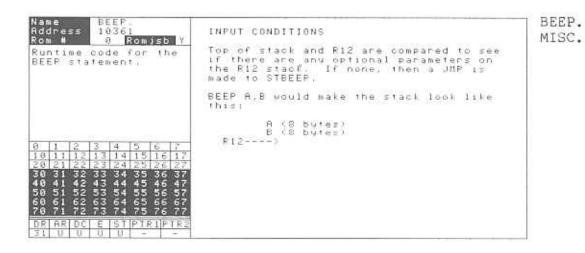


#### ASSIG. DISC

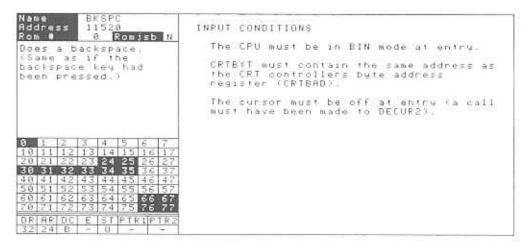


# ATN2. MATH





BKSPC CRT

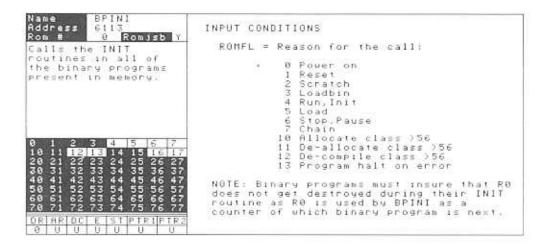


BLKLIN

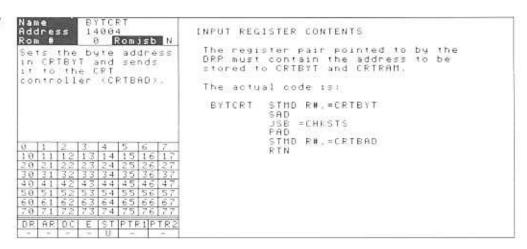
CRT

Name BLKLIN Rddress 14165 Rom ■ 0 Romjsb N	INPUT CONDITIONS
Fills from current CRT bute address to the end of the line	The CRT byte address pointer (CRTBAD) must be pointing to the address where blanking is to start.
with carriage return characters (15 octal),	OUTPUT CONDITIONS
Alters CRIBYT, leaving it pointing to the start of the next line.	The CRT bute address pointer will be pointing to the first character of the next line.
0 1 2 3 4 5 6 7	The actual code for BLKLIN is:
10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37	PERTIN RIN
40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57	LOB R32.=15 LB JSB = OUTCH1 ICB R66
60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77	CHB R66,≔80D JNZ L8
DR AR DC E ST PTR1PTR2	RTN

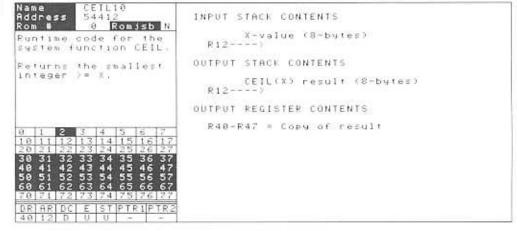
#### BPINI MISC.

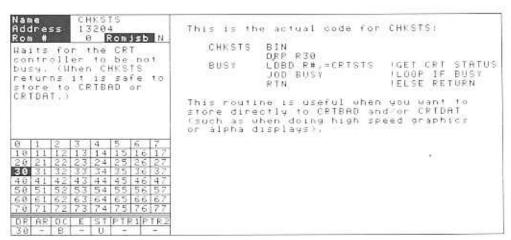


### BYTCRT CRT



# CEIL1Ø MATH

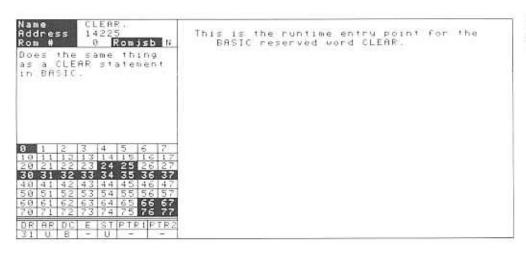




CHKSTS

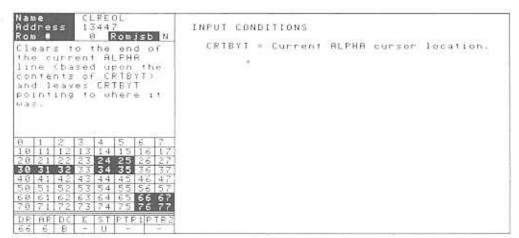
Name CHSROI Address 52672 Rom # G Romjsb N Changes the sign of a real or integer number.	INPUT STACK CONTENTS  Real or tagged-integer (8-bytes)  R12  OUTPUT STACK CONTENTS
	Real or tagged-integer (8-bytes)
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 28 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E STPTRI PTR	The actual code is:     CHSROI

CHSROI MATH

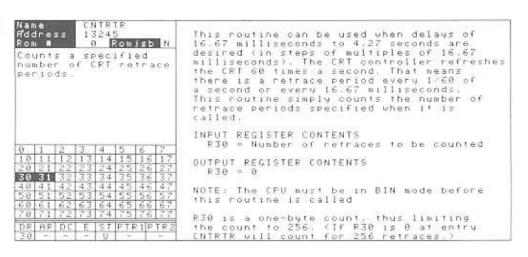


CLEAR. CRT

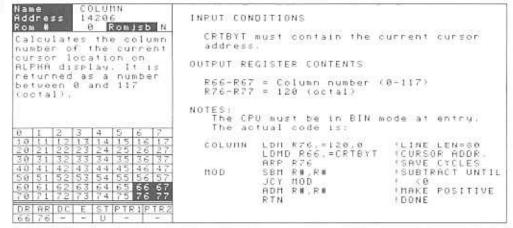
# CLREOL

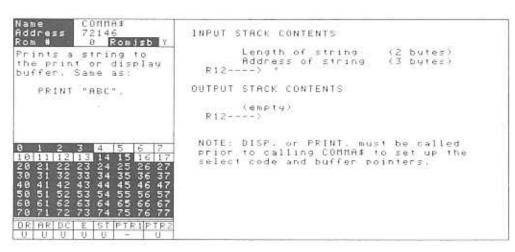


## CNTRTR CRT

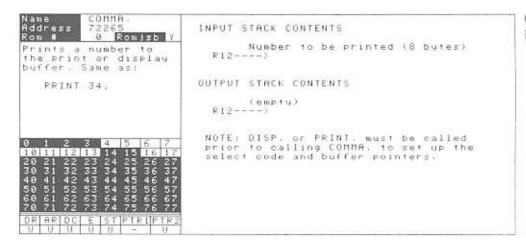


# COLUMN CRT





COMMA\$ PRINT

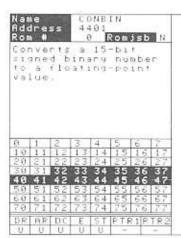


COMMA. PRINT

Name ASIG ASIGNATION A	INPUT REGISTER CONTENTS  PSS-RS7 = 23-bit signed binary number.  The 23 least significant bits are the magnitude (0 TO 2-23-1) and the most significant bit is the sign (0=positive and 1=negative) giving a range of -8388607 TO +8388607
	OUTPUT PEGISTER CONTENTS
0 1 2 3 4 5 6 7 10 11 12 12 14 15 16 17 20 21 22 23 24 25 26 27 39 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTRI PTRZ	R40-R47 = The equivalent floating-point value.

CONBI3 MATH

#### CONBIN MATH



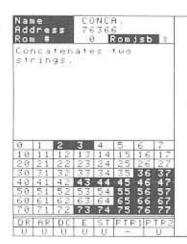
#### IMPUT REGISTER CONTENTS

P36-R37 = 15-bit signed binary number.
The 15 least significant bits are the magnitude (0-32767) and the most significant bit is the sign (0-positive and 1-negative) giving a range of -32767 to \*32767.

#### OUTPUT REGISTER CONTENTS

R40-R47 = The equivalent floating-point value.

# CONCA. MISC.



#### INPUT STACK CONTENTS

A# Length (2 bytes)
A# Address (3 bytes)
B# Length (2 bytes)
B# Address (3 bytes)

#### OUTPUT STACK CONTENTS

A\$ & B\$ Length (2 bytes) A\$ & B\$ Address (3 bytes) R42----)

## CONINT MATH



# INPUT REGISTER CONTENTS

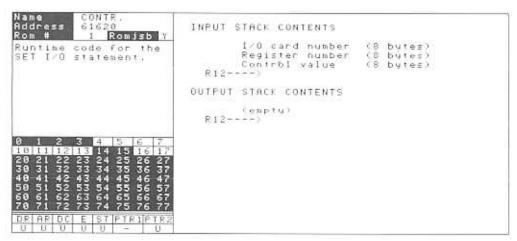
R60-R67 = Floating-point value

#### OUTPUT REGISTER CONTENTS

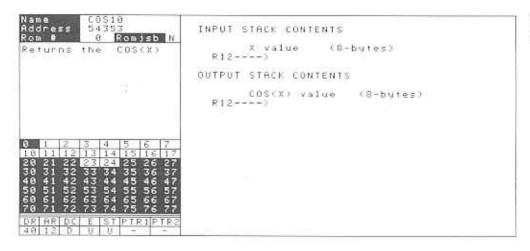
R76-R77 = 16-bit unsigned binary value R32 = Sign of value If R32=0 then value is positive If R32#0 then value is negative

NOTE: This routine doesn't check to insure that R60-R67 contains a floating point number, so if it contains a tagged-integer or some other garbage, you'll get indeterminate results.

CONINT does a SAD at entry and a PAD at exit, so all status is preserved (not including the E register),



CONTR. MISC.



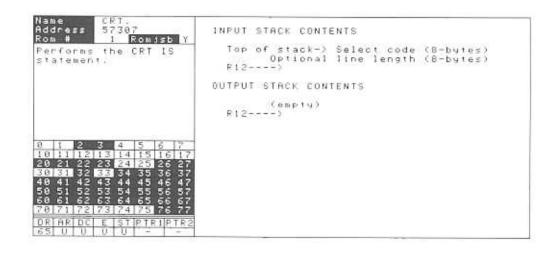
COS10

COTIØ MATH

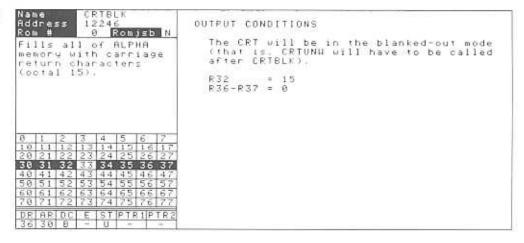
## COUNTK MISC.

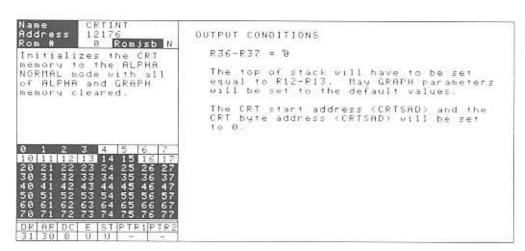
Nam Add Rom Use key	re d	26 <sub>.</sub> ,	14				N.	If a key is not pressed at entry or if it is released before the key repeat wait is done then a call is made to EOJ2, else the key repeat speed is forced to KRPET2 and the service request flags are set in RI7 and SVCHRD.
В	1	2	73	4	5	6	7	
1.0	11	12	13	14	15	16	1.7	
2.0	CA.	22	6.3	24	25	26	27	
10	41	Sel.	200	34	35	3.5	47	
50	51	52	43	54	55	46 56	57	
50	21	62	23	-	65	56	67	
20	24	-	63	2.4			991	
1.0	alos le	72	7.3	64	7.5	76	1.6	
DR	AR	D.C.	E	ST	PIR	IP.	TRE	
U	-U	8	11-1	U	7 -	100	-2.	

# CRT.



#### CRTBLK CRT





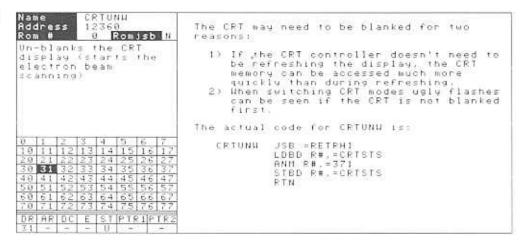
CRTINT CRT

Name Address 12334 Rom # 0 Romisb N	The code for CRTPOF is:
Turns off the CRI's high voltage supply. Nould be used if a device were hooked up to the computer that needed excessive power and didn't need to be run at the same time as the CRI display.	CRTPOF LDB R30.=6 BIN JSB =RETRHI LDBD R#.=CRTSTS ORB R#.R30 STBD R#.=CRTSTS RIN
0 1 2 3 4 5 6 7	
10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27	
20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37	
40 41 42 43 44 45 46 47	
50 51 52 53 54 55 56 57	
60 61 62 63 64 65 66 67	
70 71 72 73 74 75 76 77	
DRIAR DC E STIPTRIPTES	

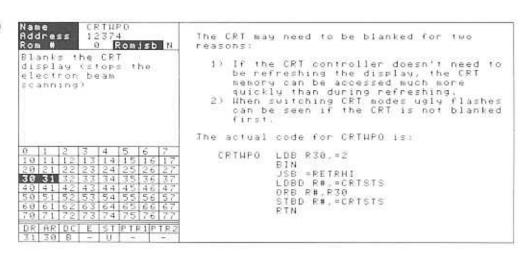
CRTPOF CRT

CRTPUP CRT

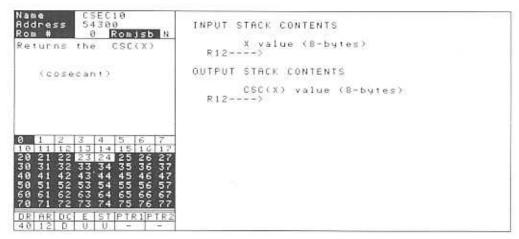
#### CRTUNW CRT

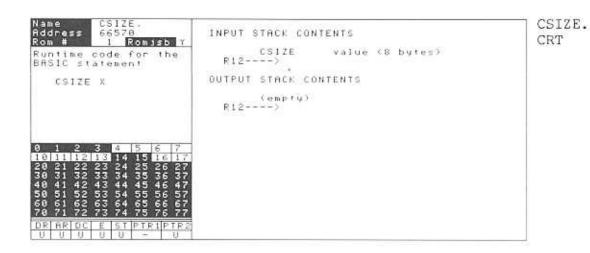


#### CRTWPO CRT

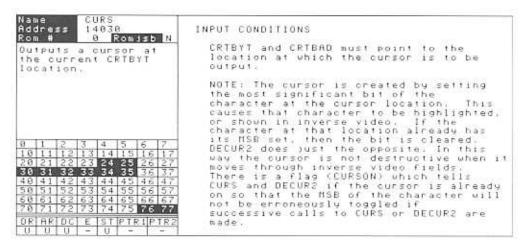


# CSEC10



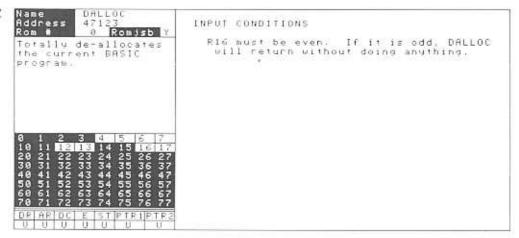


CURS CRT

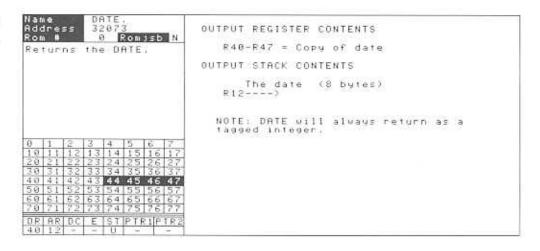


CVNUM PRINT

## DALLOC MISC.

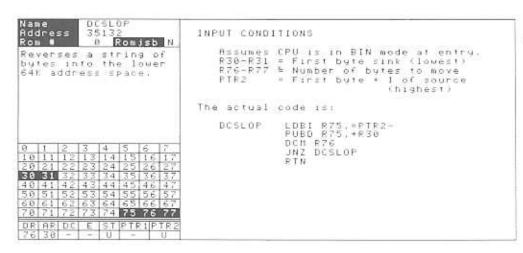


# DATE. MISC.



# DCLIN# PRINT

Name DCLIN# Address 34607 Rom # 0 Romish N	INPUT CONDITIONS
Decompiles a BASIC program line number.	R30-R31 = Pointer to output buffer R65-R67 = Line number as 5 BCD digits
	OUTPUT CONDITIONS
	R30-R31 = Pointer to output buffer (after the line number was pushed out as ASC11 characters).
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17	
20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37	
40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57	
60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77	
DR AR DC E STPIRIPIR2	



DCSLOP MISC.

ด	1	2	3	4	15	6	7
10	1.1	12	13	1.4	15	16	17
20	21	22	23	24	25	26	27
60		SE 42 S	333	34	35	36	37
Ē	EFF		1				
48	41	42	43	44	143	4.5	7
38 48 58	31 41 51	42 52	43 53	54	55	56	52

INPUT CONDITIONS

CRTBYT and CRTBAD must point to the location of the cursor.

NOTE: The cursor is created by setting the most significant bit of the character at the cursor location. This causes that character to be highlighted, or shown in inverse video. If the character at that location already has its MSB set, then the bit is cleared. DECUR2 does just the opposite. In this way the cursor is not destructive when it moves through inverse video fields. There is a flag (CURSON) which tells CURS and DECUR2 if the cursor is already on so that the MSB of the character will not be erroneously toggled if successive calls to CURS and DECUR2 are made.

DECUR2 CRT

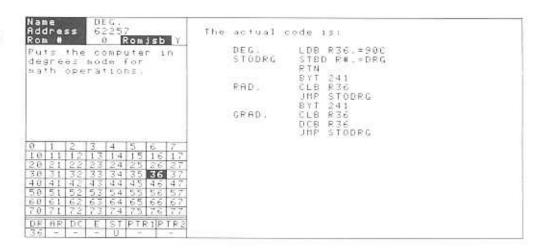
DEFA+. MATH

# Section 8: Reference Material

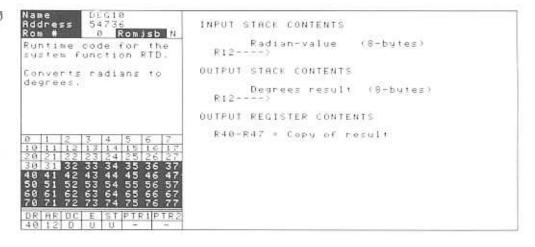
#### DEFA-. MATH

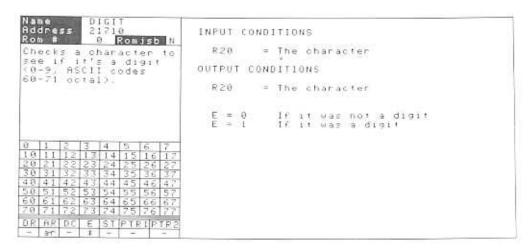
Name Address 61604 Rom # 0 Romisb Y Turns the system math defaults off.	OUTPUT REGISTER CONTENTS R36 = 0
	The actual code is:
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27	DEFA+. CLB R36 1CB R36 JMP STORDF
30 31 32 33 34 35 <b>36</b> 37 40 41 42 43 44 45 46 47 50 51 52 53 54 65 66 67 68 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77	DEFA-, CLB R36 STORDF STBD R#.=DEFAUL
DR AR DC E ST PTR1PTR2	

#### DEG. MATH



# DEG1Ø MATH





DIGIT PARSE

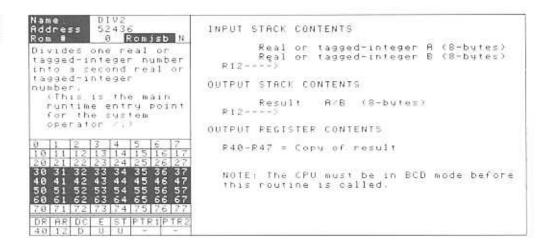
Sets up SCTEMP so that it contains the current CRT 19 select code. It's usually used prior to calling DRV12.  8 1 2 3 4 5 6 7 18 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 38 31 32 30 34 35 36 37 48 41 42 43 44 45 46 47 58 51 52 53 54 55 56 57	Name Address Rom #	DISP. 71311 0 Rom	isb i		
Code. It's usually used prior to calling DRV12.	that it	contains	the		
T0 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 48 41 42 43 4 45 46 47 50 51 52 53 54 55 56 57	code. It	t's usual	14		
10 11 12 13 14 15 15 17 20 21 22 23 24 25 25 27 80 31 32 33 34 35 36 36 48 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57					
20 21 22 25 24 25 26 27 30 31 32 33 34 35 35 35 48 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57		And the same of th	6 7		
50 51 52 53 54 55 56 57	20 21 22 30 31 32	23 24 25	26 27		
	III PARAGOSIA MONTE PARAGOS		46 47 56 57		
70 71 72 73 74 75 76 77  DR AR DC E STPTRIPTR2 40 U U	DRIARIDO	ESTIPIE	76 77		

DISP. CRT

Name Rddress 52441 Rom # 0 Romisb N Divides one real number into a second real number	INPUT STACK CONTENTS  (empty)  R12)  INPUT REGISTER CONTENTS  R48-R47 = Real number
	OUTPUT STACK CONTENTS Result B/A (8-butes) RI2)
8 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47	OUTPUT REGISTER CONTENTS R40 R47 - Copy of result D/A NOTE: The CPU must be in BCD mode before DIV10 is called.
50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST RTRIRTR2 40 12 0 U U	The two arguments must be real numbers or the result will be unknown.

DIV10 MATH

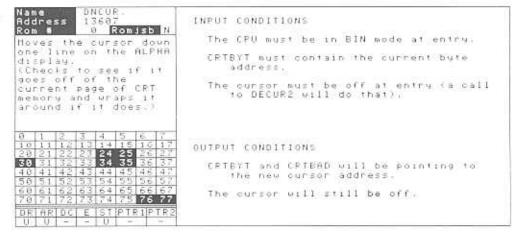
# DIV2 MATH

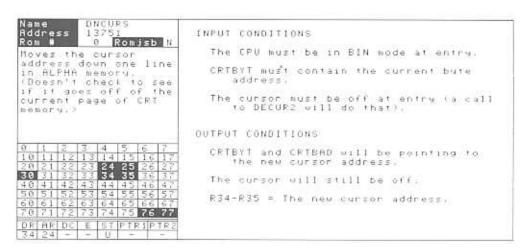


#### DMNDCR PARSE

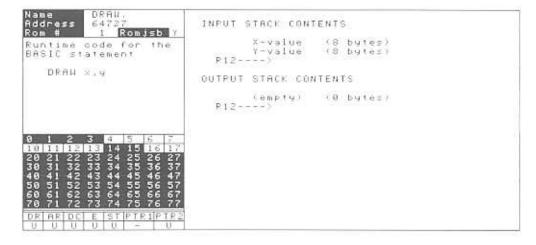
Address 25175 8 Romjsb Y Rom # Romjsb Y Flags an error if R14 is not a carriage, return or an @ token.	INPUT CONDITIONS  R14 = Incoming token  OUTPUT CONDITIONS  If R14 didn't contain a carriage return token or an @ token then ERROR will have been called.
8 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 37 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 08 88 00 0 0 0 0 0 0 0 0 0 0	

# DNCUR. CRT

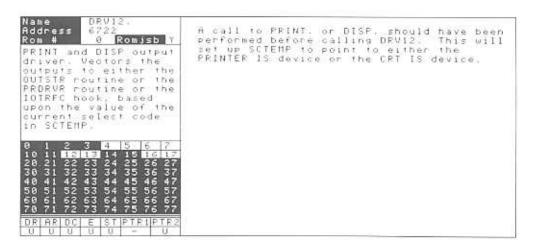




DNCURS CRT



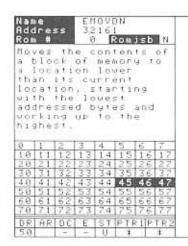
DRAW. CRT



Á

DRV12. PRINT

## EMOVDN MISC.



#### INPUT CONDITIONS

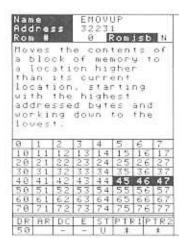
The CPU must be in BIN mode at entry-PTR1 = Last byte of source (low) PTR2 = Last byte of sink (low) P45-P47 = Number of bytes to move

#### OUTPUT CONDITIONS

PTR1 = First byte of source +1 (high) PTR2 = First byte of sink +1 (high) R45-R47 = 0

NOTE: EMOVUP and EMOVON are backwards from MOVUP and MOVON. In other words, EMOVUP does for extended memory what MOVON does for the lower 64K of memory and EMOVON corresponds to MOVUP.

# EMOVUP MISC.



#### INPUT CONDITIONS

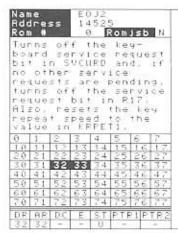
The CPU must be in BIN mode at entry.
PTR1 = First byte of source +1 (high)
PTR2 = First byte of sink +1 (high)
R45-R47 = Number of bytes to move

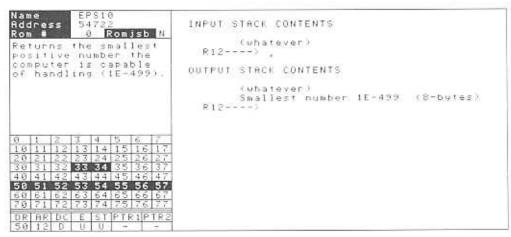
#### OUTPUT CONDITIONS

PTR1 = Last byte of source (low) PTR2 = Last byte of sink (low) R45-R47 = 0

NOTE: EMOVUP and EMOVDN are backwards from MOVUP and MOVDN. In other words, EMOVUP does for extended memory what MOVDN does for the lower 64K of memory and EMOVDN corresponds to MOVUP.

# EOJ2 MISC.

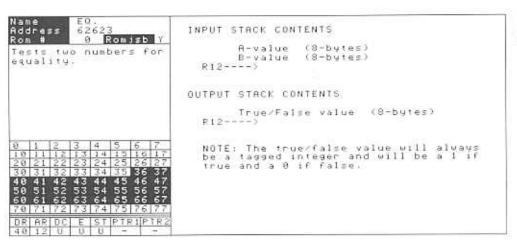




EPS10 MATH

Name EQS. Rddress 3564 Rom # 0 Romjsb N Checks two strings for equality.	INPUT STACK CONTENTS  Length of string 'A' (2-bytes) Address of string 'A' (3-bytes) Length of string 'B' (2-bytes) Address of string 'B' (3-bytes) R12>
	OUTPUT STACK CONTENTS  True/False value (8-bytes) R12)
8 1 2 3 4 5 6 7 18 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 38 31 32 33 34 35 36 37 10 11 12 43 44 45 46 47 59 51 52 53 54 55 56 57 60 51 52 63 64 65 66 67 78 71 72 73 74 75 76 77 DR RR DC E ST PTR PRZ	OUTPUT REGISTER CONTENTS  R78-R77 = Copy of true/false value.  NOTE: The true/false value is =0 if false, =1 if true and is in floating-point format.

EQ\$. MATH

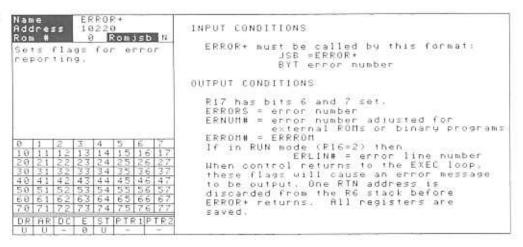


EQ. MATH

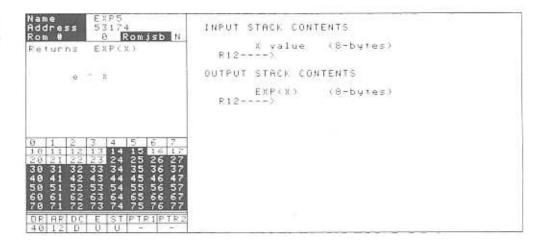
## ERROR MISC.

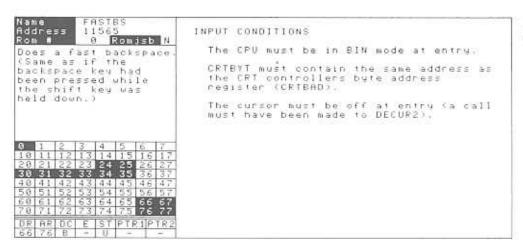
Name Addr Rom Sets repo	f 1	1 ags		4 Roma	sb N	INPUT CONDITIONS  ERROR must be called by this format:  JSB =ERROR BYT error number
						OUTPUT CONDITIONS  R17 has bits 6 and 7 set.  ERRORS = error number  ERNUM# = error number adjusted for  external ROMs or binary program
0 11	12	13	14	5	6 7	ERRON# = EPRRON
101	1 12	1 1 3		1.5	16 17	If in RUN mode (R16°2) then ERLIN# = error line number
20 2	1 22	23	154	25	26 27	When control returns to the EXEC loop
30 3	1 32	33	34	35	36 37	these flags will cause an error message
40 4	1 42				46 47	to be output All pagistage and saved
50 5	1 52				56 57	
<u> 독명 등</u>	1 52	63	164	65	55 57	
1011	1 1 6 2	173	1/4	1.3	76 77	
DRIA	R D	E	ST	PTR	1 PTR	
UIL	1 -	.0	U	-	-	

#### ERROR+ MISC.

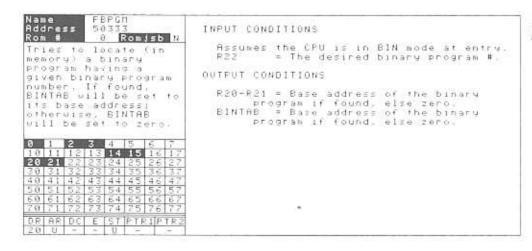


# EXP5 MATH





FASTBS CRT



FBPGM MISC.

Name	INPUT STACK CONTENTS PTC to variable area (3 butes	y
etches the address n array element FEIch Array Variable ddress).	Row index (2 bytes (2 bytes (2 bytes ) (2 bytes ) (1 byte) (1 byte) OUTPUT STACK CONTENTS	1
1 2 3 4 5 6 7 8 11 12 13 14 15 16 17 9 21 22 23 24 25 26 27	(empty) R12) OUTPUT REGISTER CONTENTS R20-R21 = Max len if string array R46 = Header byte of array R76-R72 = Abs. address of variable nam R75-R77 - Abs. address of element value PTR2 = Address of element value	0.0
0 31 32 33 84 85 85 85 87 86 87 88 88 88 88 88 88 88 88 88 88 87 88 88	NOTE: At entry, 'Ptr to variable area' an address which is relative to FHCURR. 'Col index' is present only if the arrais two dimensional. 'Dim flag' is even the array is two dimensional, odd if on	y i f

FETAVA MISC.

#### Section 8: Reference Material

#### FETSVA MISC.

Name Address Rom # 8 Romish N FETch Simple Variable Address, Takes an address which is relative to FUCURR and changes it to an absolute address. 8 1 2 **3** 4 5 6 7 10 11 12 13 14 15 16 17 20 21 32 38 34 35 36 3. 30 31 32 38 34 45 46 47 40 41 42 43 44 45 46 47 40 41 52 51 54 55 55 54 65 66 67 78 71 72 DR AR DC E ST PTR 1 PTR. 46 47 B - U - #

INPUT REGISTER CONTENTS

R65-R67 = Address of variable are (relative to FHCURR)

OUTPUT REGISTER CONTENTS

R46 " Header byte of variable R70-R72 = Abs. address of variable name PTR2 = Abs. address of least significant byte of address of variable name in variable

storage area.

#### FLIP. MISC.

The actual code for FLIP is:

LDB R36, #200 STBD R36, #KEYSTS RTN FLIP.

# FNDLIN MISC.

Name Address Rom # FNDLIN 32355 8 Romist N Finds a specified line (by number) in a BASIC program. 50 51 52 53 54 55 50 51 52 53 54 55 60 61 62 63 64 65 70 71 72 73 74 75 DR AR DC E ST PTR1P 65 6 8 \* U \*

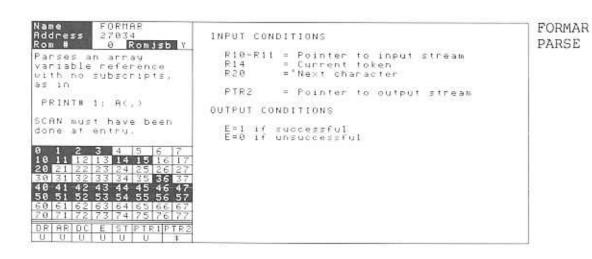
INPUT REGISTER CONTENTS

R75-R77 = Line number to be found

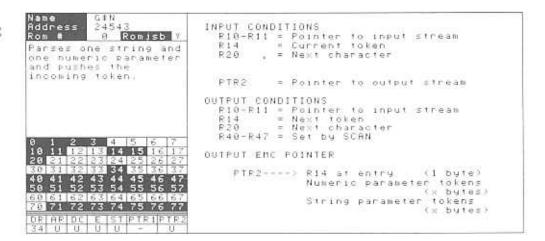
**DUTPUT CONDITIONS** 

E = 0 if the line was found. E = 17 if the line was not found. PTR1 points to the line length byte of the desired line, if found, else it points to the same byte of the next highest line.

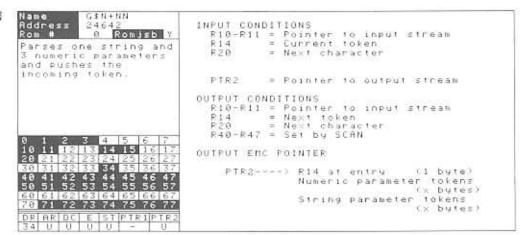
NOTE: Upon return a LDBI R20.=PTR1-would load into R20 the length of the found line, that is PTR1 is really pointing to the least significant bute of the line number. fill registers are saved and restored.



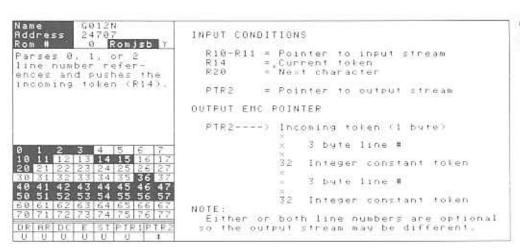
#### G\$N PARSE



#### G\$N+NN PARSE



G/A CRT



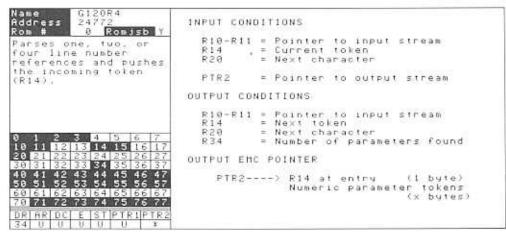
GØ12N PARSE

GØ1N PARSE

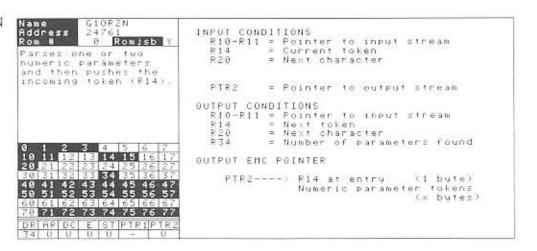
Name GBIN Address 24726 Rom # B Romisb Y	INPUT CONDITIONS
Parses zero or one line number refer- ence and pushes the	R10-R11 = Pointer to input stream R14 = Current token R20 = Next character
incoming taken (R14).	PTR2 = Pointer to output stream
	OUTPUT EMC POINTER
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47	PTR2> Incoming token (1 byte)  × 3 byte line #  × 32 Integer constant token
50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTF1PTR2 1 0 0 0 0 0 0 1 1	NOTE: The line number is optional so the output stream may be different.

GØOR2N PARSE

# G120R4 PARSE

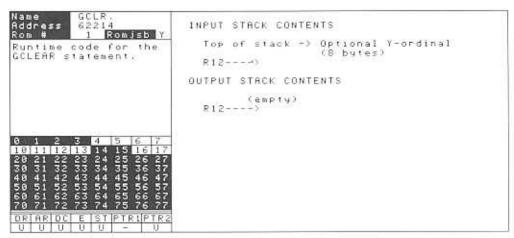


#### G10R2N PARSE

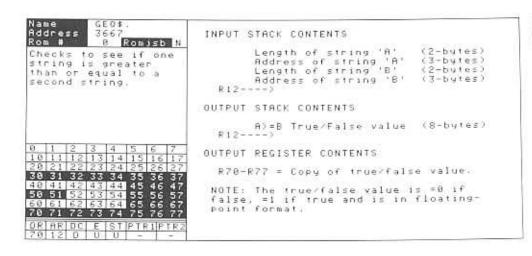


# GCHAR PARSE

Nam Add Rom Get	Γ <b>α</b> :		21		Rom		и	INPUT CONDIT	IONS Pointer to input stream
bla R20 i#	n k	ch If	ara the	01	er har re	to act	n	OUTPUT CONDI	Tions Next non-blank character
1 5								The actual c	
								GCHAR	S 8 0 B 1 N
0 20 30 40 50 50 70	71	72	3 13 23 33 43 53 7 E	74		76	7 17 27 37 47 57 57	GCH1	LOBO R20,R10 CMB R20,=15 JZR GCHRTN POBO R20,+R10 CMB R20,=40 JZR GCH1 PAD RTN
- 1	1110	-	-	-	-	-	-		



GCLR. CRT

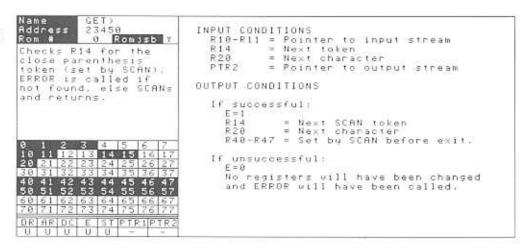


GEQ\$.

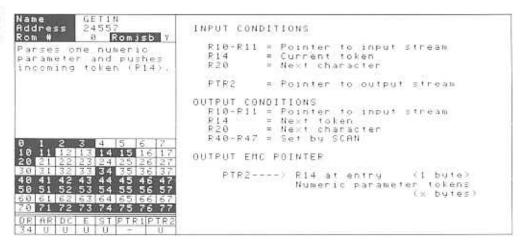
INPUT STACK CONTENTS  A-value (8-bytes) B-value (8-bytes) R12)	Wame GEQ. Address 62734 Rom # 8 Romjsb Y Tests to see if one number is greater than or equal to a second number.
OUTPUT STACK CONTENTS  A)=B True/False value (8-bytes) R12>	
NOTE: The true/false value will always be a tagged integer and will be =1 if true and =0 if false.	

GEQ. MATH

#### GET) PARSE

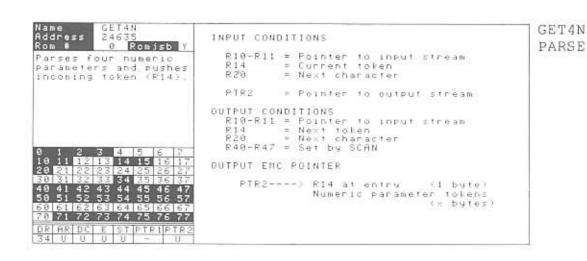


# GET1N PARSE



# GET2N PARSE

```
Name
Address
Rom #
             24630
0 Romisb Y
                                        INPUT CONDITIONS
                                           R10-R11 = Pointer to input stream
R14 = Current token
R20 = Next character
Parses two numeric
parameters and pushes incoming tolen (R14),
                                          PTR2
                                                       = Pointer to output stream
                                        OUTPUT CONDITIONS
                                          R10-R11 = Pointer to input stream
R14 = Next token
R20 = Next character
R40-R47 = Set by SCAN
     11 12 13 14 15 16
                                        OUTPUT EHC POINTER
                 34
                                              PTR2----> R[4 at entry (1 byte)
Numeric parameter tokens
(x bytes)
 60 61 62 63 64 65 56 67
70 71 72 73 74 75 76 77
DR AR DC E ST PTR1F
```

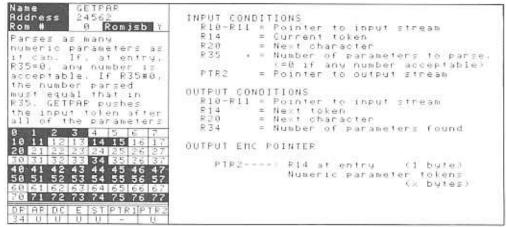


Name GETPH?
Address 24748
Rom # 8 Romjsb y
Parses as many numeric parameters as it can and then pushes the incoming token (Ri4).

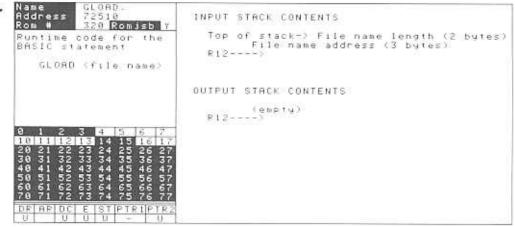
0 1 2 3 4 5 6 7 1 1 2 12 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 48 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 68 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 0 71 72 73 74 75 76 77 0 71 72 73 74 75 76 77 0 71 72 73 74 75 76 77 0 74 75 76 77 0 74 75 76 7

GETPA? PARSE

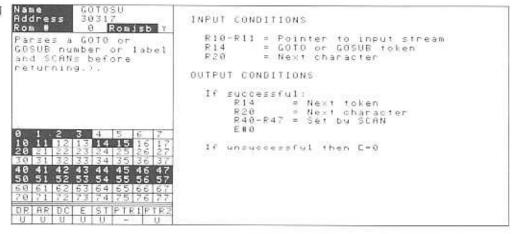
#### GETPAR PARSE

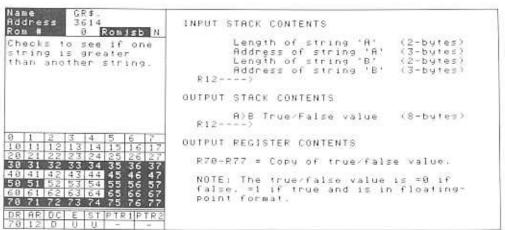


#### GLOAD. CRT

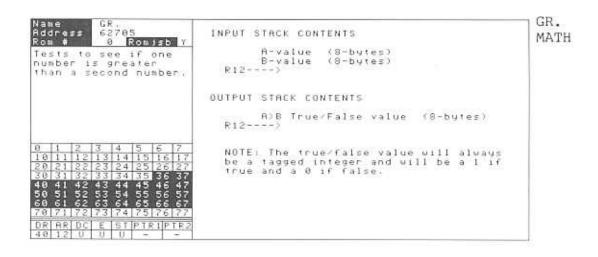


#### GOTOSU PARSE





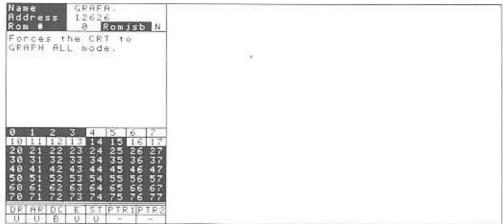
GRS. MATH



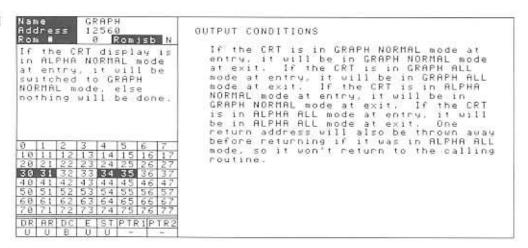
Name Address Rom # GRAD 62274 0 Romisb V The actual code is: HTAM DEG. LDB R36,=90C STBD R#.=DRG RTN Puts the computer in grads mode for STOORG math operations. BYT 241 CLB R36 JHP STODRG RAD. BYT 241 CLB R36 DCB R36 JMP STOBRG GRAD 9 1 2 3 4 5 6 16 11 12 13 14 15 16 20 21 22 23 24 25 26 30 31 32 33 34 35 36 30 31 32 33 34 35 36 48 41 42 43 44 45 46 50 51 52 53 64 55 56 50 51 52 53 54 55 56 57 68 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DP RR DC E ST PIRIPIRS 36 - - - U -

GRAD.

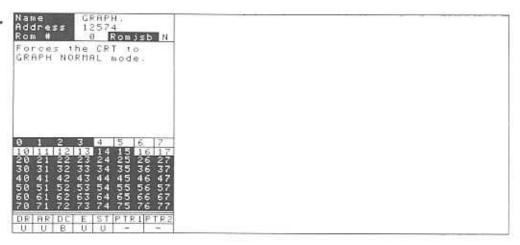
GRAFA. CRT

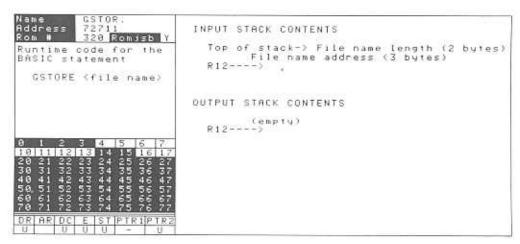


GRAPH CRT



GRAPH. CRT





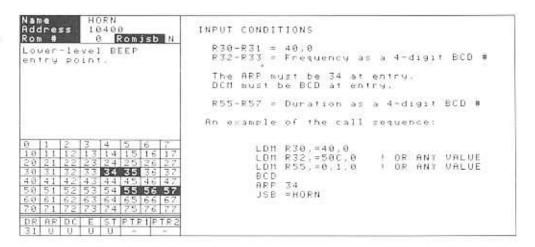
GSTOR. CRT

t ha	t pu e C = 50	ts RT. r p	051	tr av	ing on	to th	e	R26-R27 = Address of the string R36-R37 = Length of the string	
0	1	2	3	4	5	6	7		
28	21	22	23	30	25	26	27		
10	41	32	43	4.4	45	36	47		
50	51	52	53	54	55	56	57		
50		6.2	63	64	63	66	67		

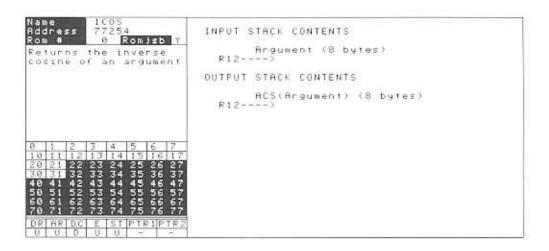
HLFLIN CRT

Name HMCURS Address 13661 Rom # Romjsb N Hoves the cursor to the top left of the ALPHA display.	INPUT CONDITIONS  The cursor must be off (a call to DECUR2 must have been made) prior to calling HMCURS.
	DUTPUT CONDITIONS
	CRTBYT will contain the new cursor address (which will be the same as the contents of CRTRAM.
0 11 12 13 14 15 16 17	R34-R35 = New cursor address.
0 11 12 13 14 15 16 17	The actual code for HMCURS is:
9 21 22 23 24 25 26 27 8 31 32 33 34 35 36 37 10 41 42 43 44 45 46 47 10 51 52 53 54 55 56 57 10 71 72 73 74 75 76 77	HRCURS LDHD R34, = CRTRAN JSB = BYTCRT RTN
OR HR DC E ST PTR 1PTR2	

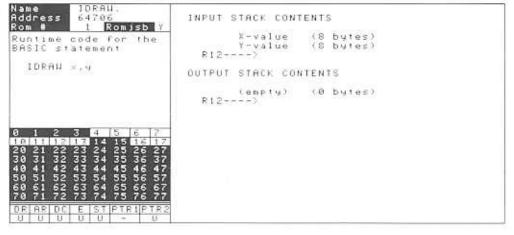
HMCURS CRT HORN MISC.

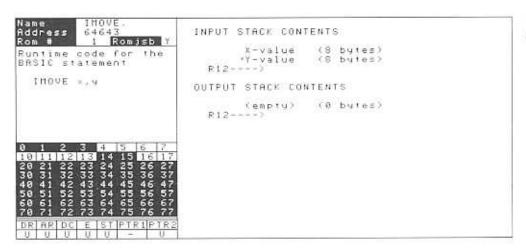


ICOS MATH



IDRAW.





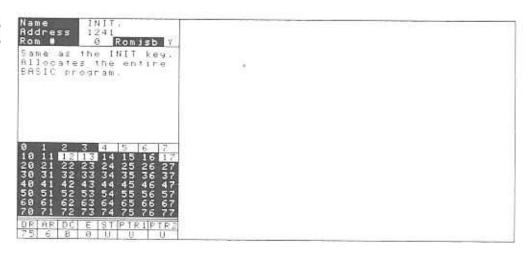
IMOVE. CRT

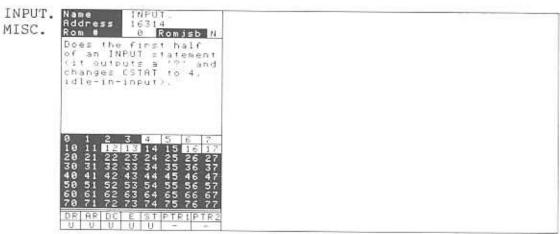
Name INCHR Address 14262 Rom # 0 Romjsb N Reads one character from the location in CRT memory pointed to by the CRT's byte address register (CRTBAD).	INPUT CONDITIONS  The CRT's internal byte address pointer (which is set by storing to CRTBAD) must be pointing to the address of the byte to be read.  OUTPUT REGISTER CONTENTS  R32 = Character from CRT memory.
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17	NOTE: The CPU must be in BIN mode at entry. The actual code for INCHR is:
20 21 22 23 24 25 26 27 30 31 <b>32</b> 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57	INCHR DRP 32  JSB =BUSY   Refer to CHKSTS  ICB R#  STBD R#,=CRTSTS
60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR RR DC E ST PTR1PTR2 32 - 8 - U	LOOP2 LOBD RW, -CRISTS JOD LOOP2 LOBD R32, -CRIDAT RIN

INCHR CRT

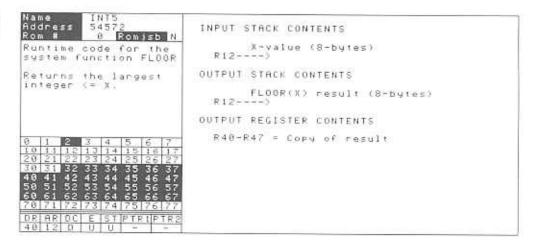
Name INF10 Address 54321 Rom # 0 Romisb N Returns the largest number that can be handled by the computer:	INPUT STACK CONTENTS  (whatever) R12> OUTPUT STACK CONTENTS
9.99999999999E499  9.10 11 12 13 14 15 16 17  20 21 22 23 24 25 26 27  38 31 32 33 34 35 36 37  40 41 42 43 44 45 46 47  50 51 52 53 54 55 56 57	(whatever) 9,999999999998499 (8-bytes) R12>
60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTR1PTR2 48 12 D U U	

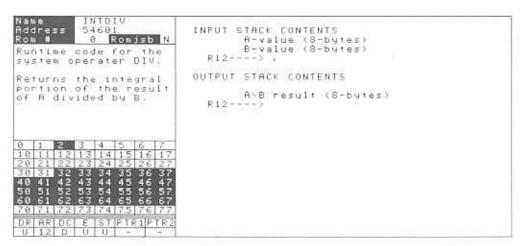
INF1Ø MATH INIT. MISC.



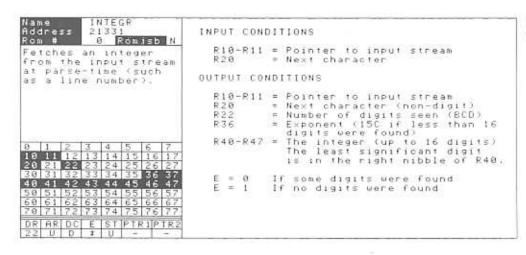


INT5 MATH





INTDIV MATH

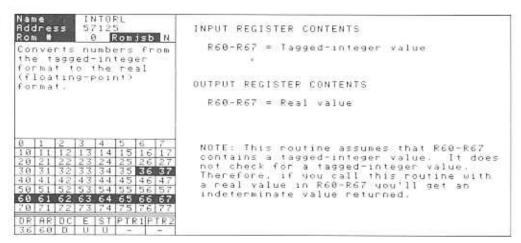


INTEGR PARSE

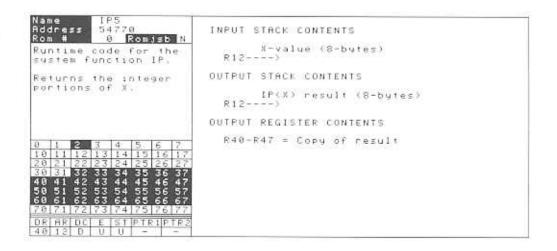
ame 1NTHUL ddress 53673 on 8 Romish N	INPUT REGISTER CONTENTS
ultiplies two 16-bit inary numbers giving	R66-R67 = 16-bit binary number A R76-R77 = 16-bit binary number B
32-bit binary . esult.	OUTPUT REGISTER CONTENTS
	R66-R67 = 16-bit binary number A R76-R77 = 16-bit binary number B R54-R57 = 32-bit result A#B
1 2 3 4 5 6 7 9 11 12 13 14 15 16 17 9 21 22 23 24 25 26 27 9 31 32 33 34 35 36 37 9 41 42 43 44 45 46 47 9 51 52 53 54 55 56 57 9 61 62 63 64 65 66 67 9 71 72 73 74 75 76 77	NOTE: INTMUL does a SAD at entry and a PAD at exit and saves and restores all registers used except for RS4-RS7.
RIARIDO E ISTIPTRIPTRE	

INTMUL MATH

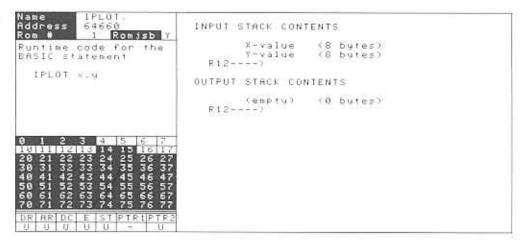
#### INTORL MATH

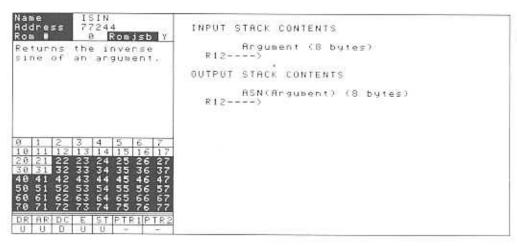


#### IP5 MATH

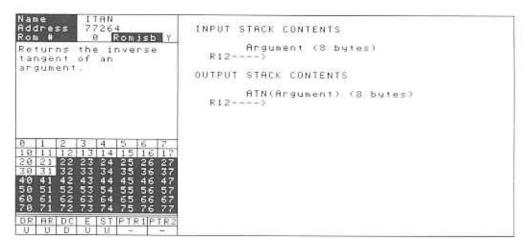


#### IPLOT. CRT





ISIN MATH

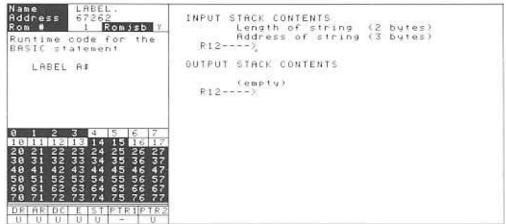


ITAN MATH

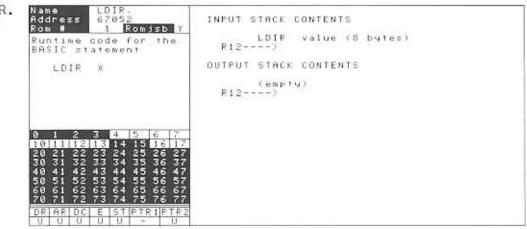
Name KEYLA, Address 13360 Rom # 0 Romisb N This routine does the same as the KEY LABEL key.					on doe	s t	he.	INPUT CONDITIONS  If R66 = 140 then the CRLC mode key labels will be displayed, else the RUN mode key labels will be display
							7	
3	1	2	3	4	5	6		
0	1 1 1	12	13	14	15	16	12	
8	1 1 2 1	12	3 13 23	4 14 24	15 25	16	27	
8	1 11 21	_	3 13 23 38	4 14 24 34	15 25 35	16	17 27 37	
0 0	1 11 21 31	12 22 32 42	3 13 23 68 43	4 14 24 34 44	15 25 35 45	16 26 36 46	17 27 37 47	
8 8	1 21 31 41 51	12 22 32 42 52	13 23 38 43 53	14 24 34 44 54	15 25 35 45 55	16 26 36 46 56	17 27 37 47 57	
8		12 22 32 42 52 62	13 23 38 43 53	4 24 34 44 54	15 25 35 45	16 26 46 56	17 27 37 47 57	
8 8 8 8 8		12 22 32 42 52	13 23 38 43 53	14 24 34 44 54	15 25 35 45 55	16 26 36 46 56	17 27 37 47 57	
1 8 2 8 3 9 3 9 5 9 7 9	61 71	12 22 32 42 52 62 72	13 23 43 53 63 73	14 24 34 44 54	15 35 45 55 65 75	16 26 36 46 56 66 76	17 27 37 47 57	

KEYLA. MISC.

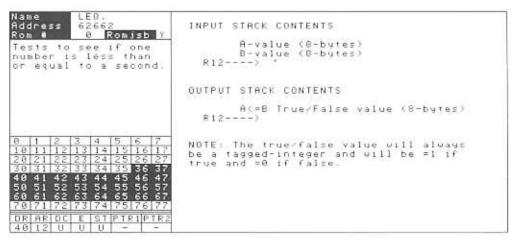




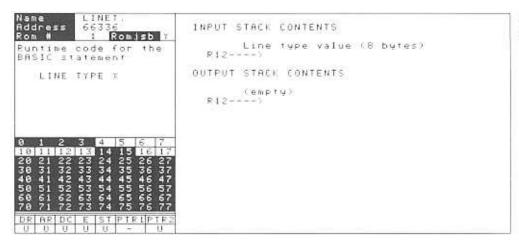
#### LDIR. CRT



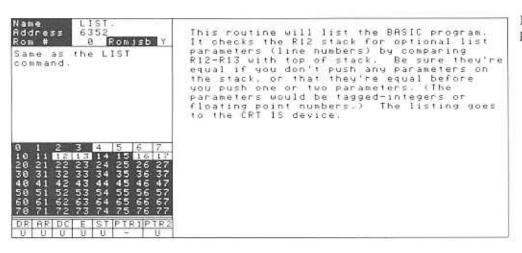
## LEQ\$. MATH



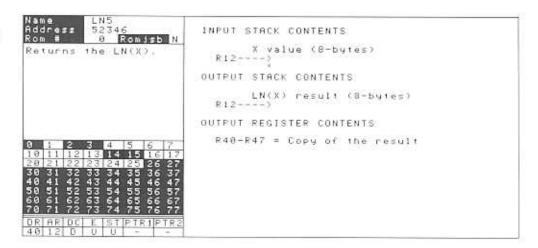
LEQ. MATH



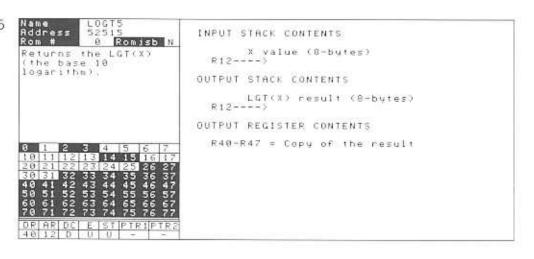
LINET. CRT



LIST. PRINT LN5 MATH

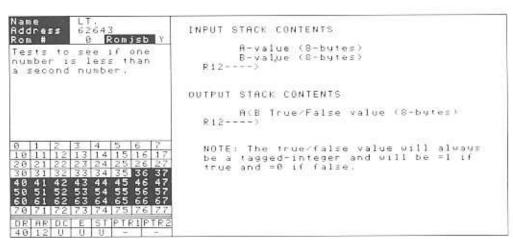


LOGT5 MATH



LT\$. MATH

Name LIS. Address 3635 Rom # B Romjsb N Checks to see if one string is less than a second string.	INPUT STACK CONTENTS  Length of string 'A' (2-bytes) Address of string 'A' (3-bytes) Length of string 'B' (2-bytes) Address of string 'B' (3-bytes) R12)
	OUTPUT STACK CONTENTS
	A(B True/False value (8-bytes)
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17	OUTPUT REGISTER CONTENTS
20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37	R70-R77 = Copy of true/false value.
40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 50 51 52 63 64 65 66 67 70 71 72 73 74 75 76 77	NOTE: The true/false value is =0 if false, =1 if true and is in floating-point format.
DR AR DC E ST PTRIPTR2 70 12 D U U	



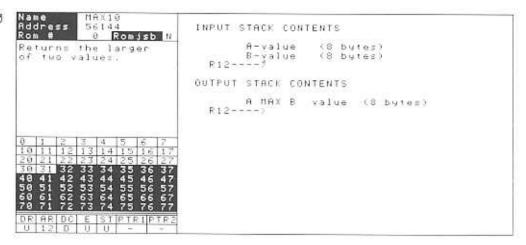
LT. MATH

Name LTCUR Address 13623 Rom # 8 R INPUT CONDITIONS 8 Romasb M The CPU must be in BIN mode at entry. Moves the cursor left one space on the ALPHA display. CRIBYT must contain the current byte. (Checks to see if it goes off of the current page of CRT memory and wraps it around if it does.) The cursor must be off at entry (a call to DECUR2 will do that). 4 5 t OUTPUT CONDITIONS CRTBYT and CRTBAD will be pointing to the new cursor address. 30 68 61 62 63 64 65 66 67 78 71 72 73 74 75 76 77 The cursor will still be off. DR AR DC E ST PTRIPTR2 R34-R35 = The new cursor address.

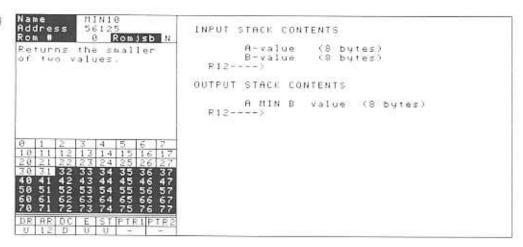
LTCUR. CRT

Name Address Rom # INPUT CONDITIONS 13757 8 Romisb N The CPU must be in BIN mode at entry. Moves the cursor address left one space in ALPHA memory (Doesn't check to see if it goes off of the current page of CPT memory.) CRIBYT must contain the current byte address. The cursor must be off at entry (a call to DECUR2 will do that). OUTPUT CONDITIONS CRIBYT and CRIBAD will be pointing to the new oursor address. 14 15 16 17 30 The cursor will still be off. 40 41 50 51 R34-R35 = The new cursor address. 64 65 66 67 74 75 **76 77** OR AR DC E ST PTRIPTR

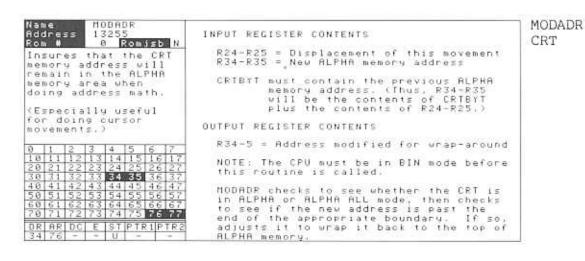
LTCURS CRT MAX10 MATH

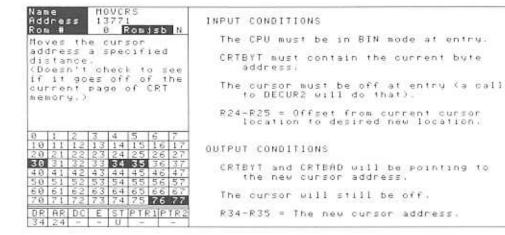


MIN10 MATH



MOD10 MATH



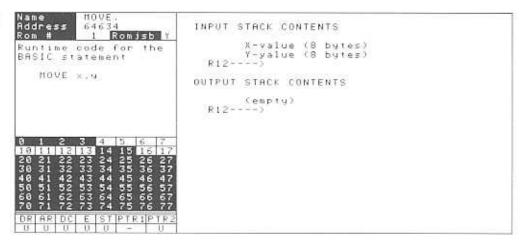


MOVCRS CRT

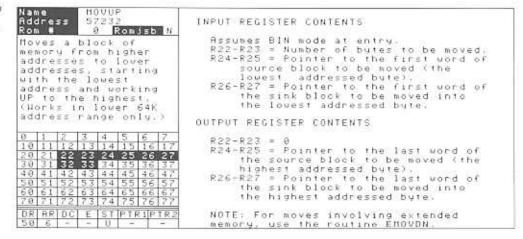
	/DN 172 Romisb N	INPUT REGISTER CONTENTS
Hoves a bloomemory from addresses to addresses, such the hiddess and bound to the (Works in le	lower to higher starting ghest working lowest.	Assumes BIN mode at entry.  R22-R23 = Number of bytes to be moved.  R24-R25 = Pointer to the first word of source block to be moved (the highest addressed byte).  R26-R27 = Pointer to the first word of the sink block to be moved into (the highest addressed byte).
iddress ran	ge only.)	OUTPUT REGISTER CONTENTS
20 21 22 23 30 31 32 33 40 41 42 43 50 51 52 53 50 61 62 63	54 55 56 57	R22-R23 = 0 R24-R25 = Pointer to the last word of the source block to be moved (the lowest addressed byte). R26-R27 = Pointer to the last word of the sink block to be moved into (the lowest addressed byte).
	ST PTR1PTR2	NOTE: For moves involving extended memory, use the routine EMOVUP.

MOVDN MISC.

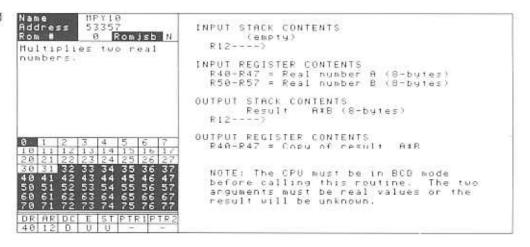
#### MOVE. CRT

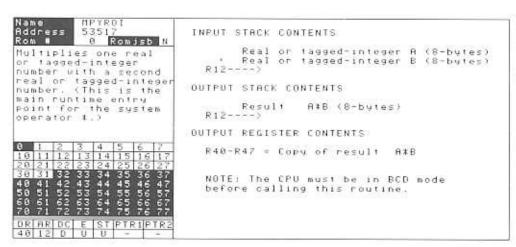


#### MOVUP MISC.

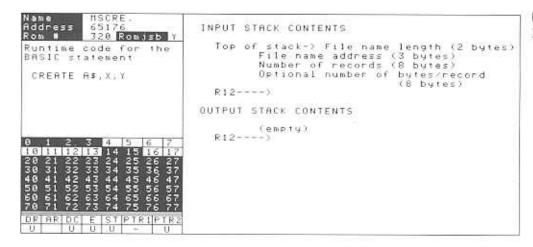


#### MPY1Ø MATH





MPYROI MATH

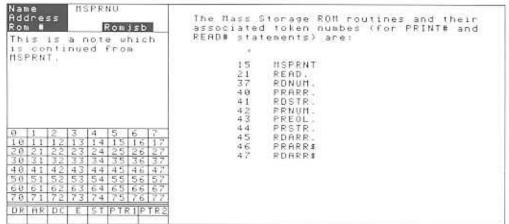


MSCRE. DISC

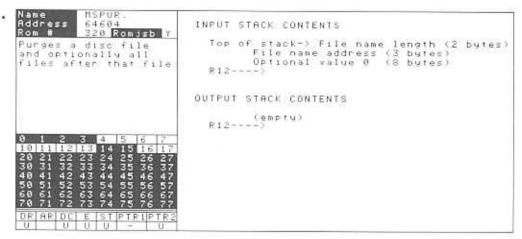
Name Address 66221 Rom # 320 Romisb Y Sets the file print pointers to the appropriate file buffer. Part of the statement	INPUT STACK CONTENTS  Top of stack-> Buffer number (8 bytes) Optional record # (8 bytes) R12>
PRINT# 1	OUTPUT STACK CONTENTS R12
8 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTRIPTR2 U U U U — U	NOTE: For figuring out what routines to call and in which order when reading from or printing to disc data files, first write it as a BASIC statement first line of a program. Using the HEM command, look into memory (the line will be at FHCURR-40) to see what the token form is. Refer to HSPRNU for a list of routines and token numbers.

MSPRNT DISC

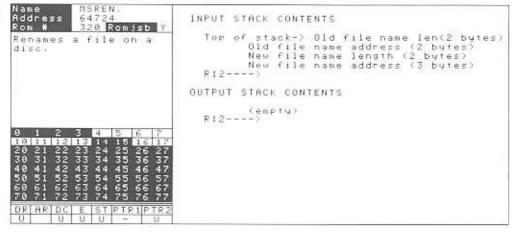
#### MSPRNU DISC

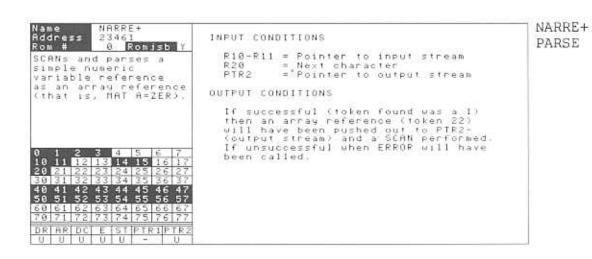


#### MSPUR. DISC



#### MSREN. DISC

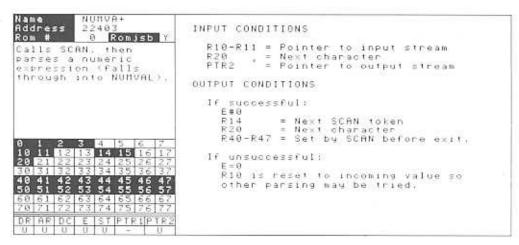




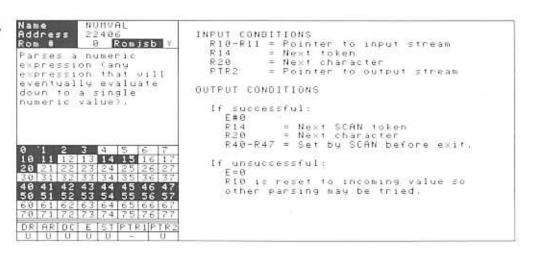
NARREF NARREE Address Rom # INPUT CONDITIONS 23465 8 Romisb Y PARSE R10-R11 = Pointer to input stream R14 = Next token R20 = Next character This routine parses a simple numeric = Pointer to output stream variable reference as an array reference (that is, MAT A=ZER). DUTPUT CONDITIONS If successful (token found was a 1) then an array reference (token 22) will have been pushed out to PTR2-(output stream) and a SCRN performed. If unsuccessful then ERROR will have been called. 10 11 12 13 14 15 16 1 20 21 22 23 24 25 26 2 30 31 32 33 34 35 36 3 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77

NUMCON PARSE

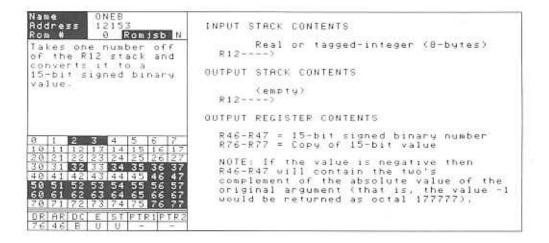
#### NUMVA+ PARSE

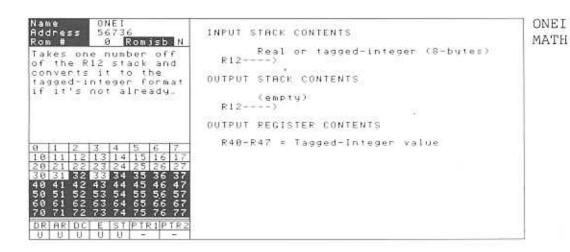


#### NUMVAL PARSE



#### ONEB MATH



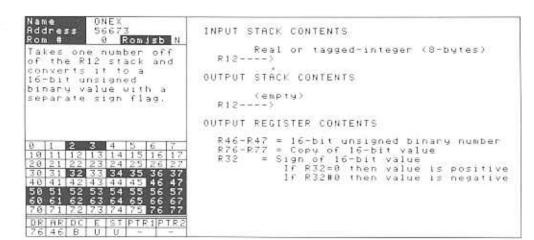


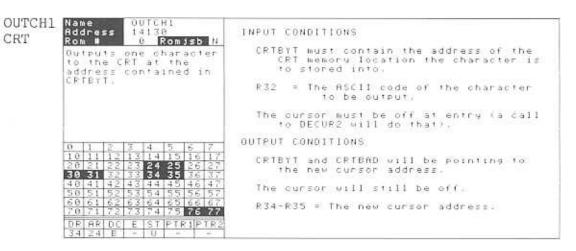
ONER

MATH

Name ONER Address 56777 Rom # 0 Romjsb N Takes one number off of the R12 stack and, if it's not in the	INPUT STACK CONTENTS  Real or tagged-integer (8-bytes) R12)
real (floating point) format, converts it to that format.	OUTPUT STACK CONTENTS  (empty) P12>
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTR1PTR2 50 40 0 0 0	OUTPUT REGISTER CONTENTS  R40-R47 = Real value from R12 stack R60-R67 = Copy of real value from stack

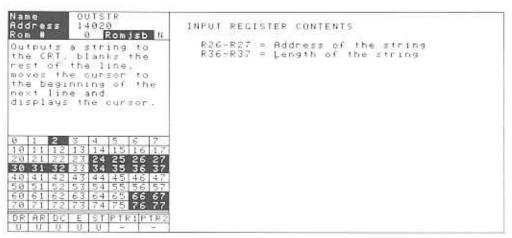
ONEROI MATH ONEX MATH





OUTCHR CRT

Name	INPUT CONDITIONS
Outputs one character to the CRT at the address contained in	CRTBYI must contain the address of the CRT memory location the character is to stored into.
CRTBYT and scrolls display up if the cursor position moves	R32 = The ASCII code of the character to be output.
off of the bottom- right corner.	The cursor must be off at entry (a call to DECUR2 will do that).
0 1 2 3 4 5 6 7	OUTPUT CONDITIONS
10 11 12 12 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37	CRIBIT and CRIBHU will be pointing to the new cursor address.
40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57	The cursor will still be off.
60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77	R34-R35 = The new cursor address.
DR AR DC E ST PTR1PTR2	



OUTSTR CRT

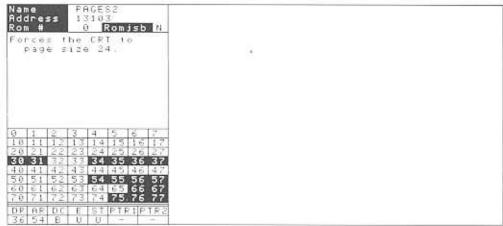
Name PAGES. Address 12756 Rom # Romisb N Sets the CRT to page size 16 or to page size 24 (same as the PAGESIZE statement).	INPUT STRCK CONTENTS  Page size (16 or 24 decimal)(8 bytes)  R12)  OUTPUT STRCK CONTENTS  (empty)  R12)
9 1 2 3 4 5 6 7 18 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 48 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 73 71 72 73 74 75 76 77 DR AR DC E ST PTRIPTR2	NOTE: Gives a warning message if the parameter is not equal to 16 or 24.

PAGES. CRT

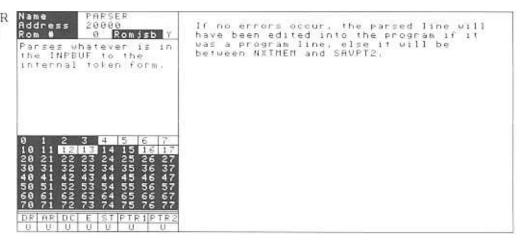
Nar Ade Ror	ne dre	8 5		GE: 00		j≨b	11
For	rce pag	s t	he izq	CR	5		
0	1	2	3	4	5	6	7
28	21	12	13	24	25	26	27
33	251	32	33	EO	35	36	37
40	41	42	43	44	4.5	4.6	4.7
50 60	61	62	53	LEI	135	56	57 67
7.0	71	72	73	74	65	76	77
DR	AR	DC	E	ST	PTF	-	
36			1.1	111			_

PAGES1 CRT

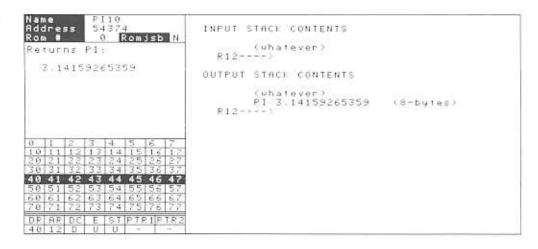
#### PAGES 2 CRT

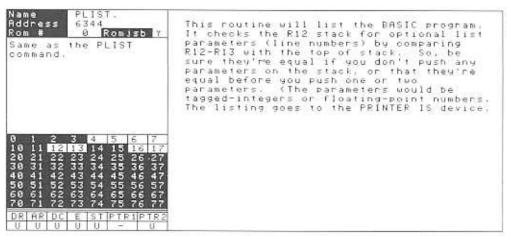


#### PARSER PARSE

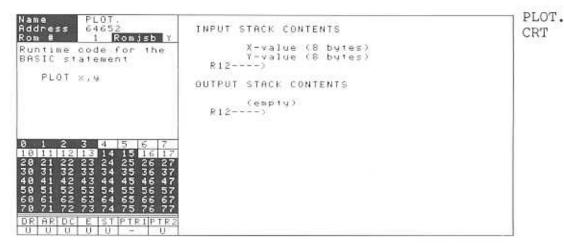


### PI10 MATH





PLIST. PRINT

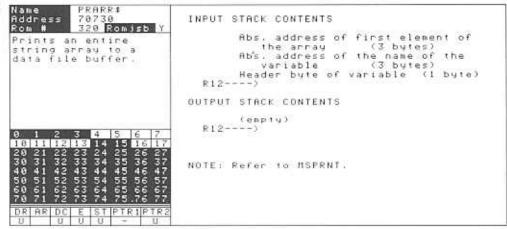


CRT

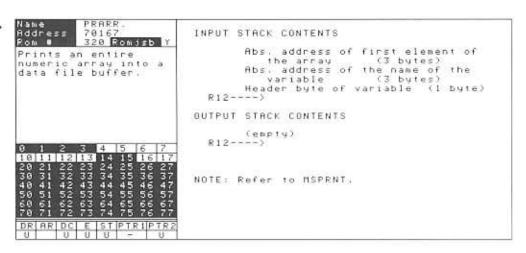
Name Address 4227 Rom # 0 Romisb N Returns a value which is the position in one string of a second string. This is the runtime code for the system function POS.	INPUT STACK CONTENTS  Length of arg string 'A' (2 bytes) Address of arg string 'A' (3 bytes) Length of arg string 'B' (2 bytes) Address of arg string 'B' (3 bytes) R12> OUTPUT STACK CONTENTS  Position of string B in string B
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTRIPTE2	R12) (8 bytes)  NOTE: Position value will be 0 if string B does not exist in string A.

POS. MISC.

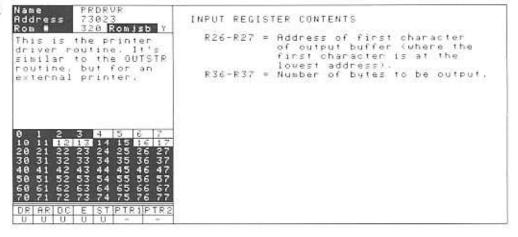
#### PRARR\$ DISC

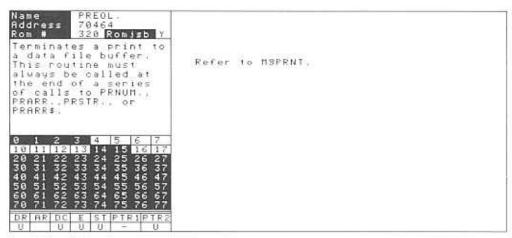


#### PRARR. DISC

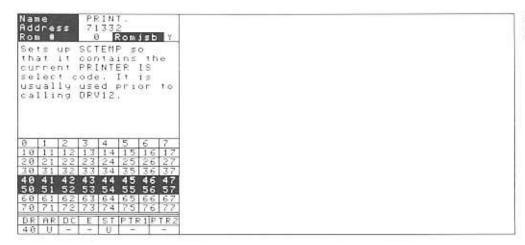


#### PRDRVR PRINT

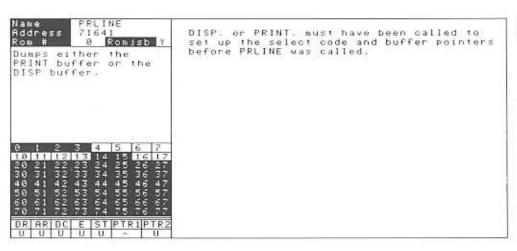




PREOL. DISC



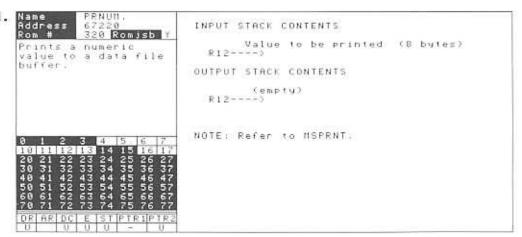
PRINT. PRINT



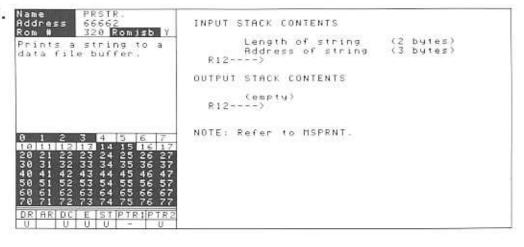
PRLINE PRINT

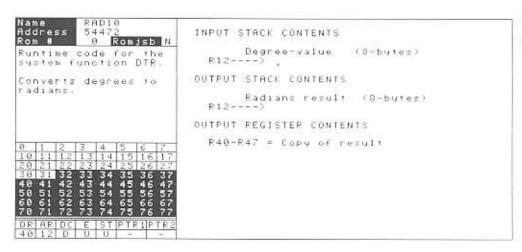




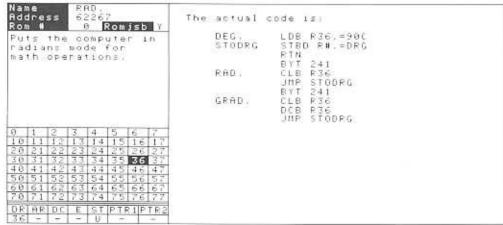








RAD1Ø MATH



RAD.

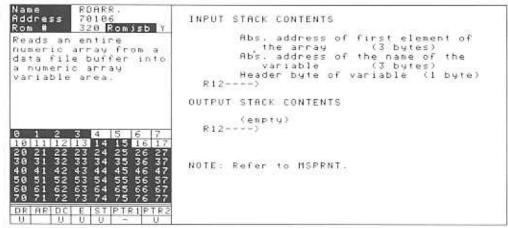
MATH

```
Name RDARR$
Address 70312
Rom # 320 Romisb V
Reads an entire string array from a data file buffer.into a string array variable area.

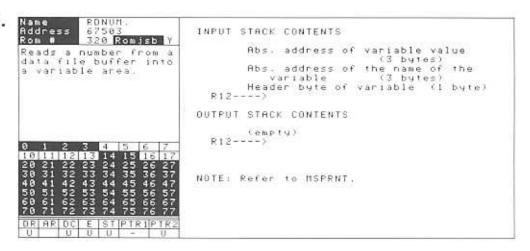
O 1 2 3 4 5 6 7
IO 11 12 13 14 15 16 17
20 21 22 23 24 25 26 27
30 31 32 33 34 34 34 45 46 47
50 51 52 53 54 55 56 57
60 61 62 63 64 65 66 67
70 71 72 73 74 75 76 77
DR AR DC E STPTRIPTR2
U U U U U - U
```

RDARR\$ DISC

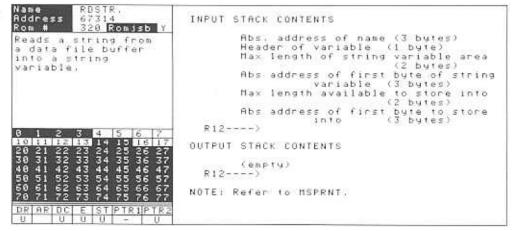
#### RDARR. DISC

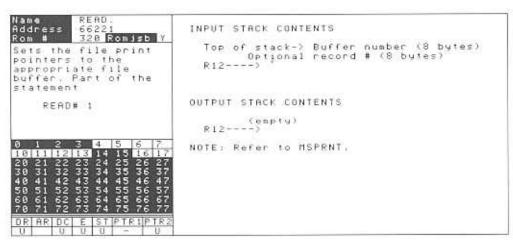


#### RDNUM. DISC

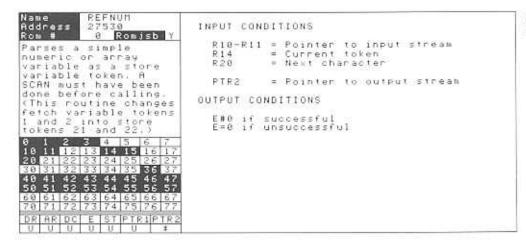


#### RDSTR. DISC





READ. DISC

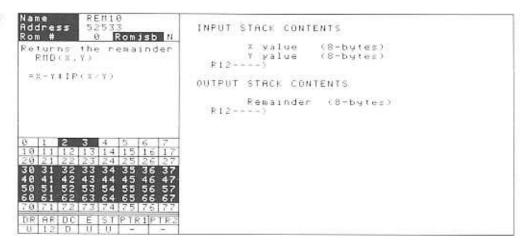


REFNUM PARSE

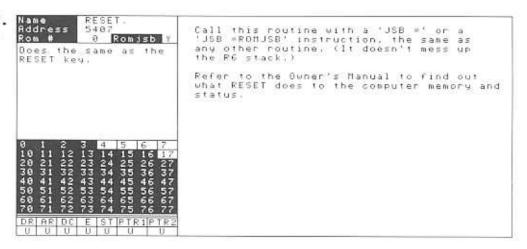
0 3	2 1 3 1 4 1 5 1	22 42 42 62	23 43 53 63	24 34 44 54 64	25 45 55	26 36 46 56	27 37 47 57	
9 1	1 1	2	3	4 1.4	5	6	7	
ele emo	eas on c	se d t ved	all hat	t e u :	vil espe es ell	ora		of memory currently reserved.

RELMEM MISC.

#### REM10 MATH

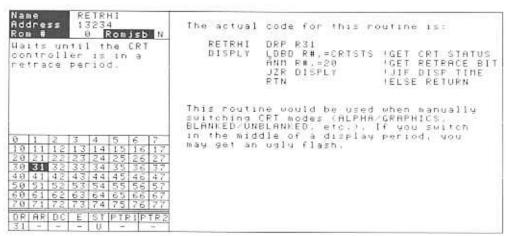


#### RESET. MISC.

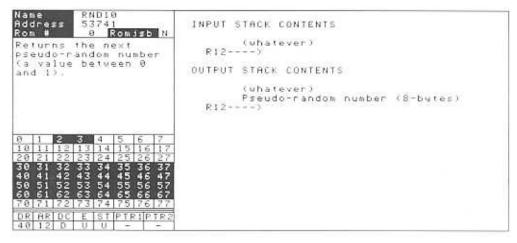


RESMEM MISC.

Name RESHEH Address 31741 Rom # 0 Romisb	INPUT REGISTER CONTENTS
Reserves temporary soratchpad memory. (It gets released at the end of each line of a BASIC program and at each @ sign (concatenation of statements).)	R55-R57 = Number of bytes to be reserved OUTPUT REGISTER CONTENTS  R55-R57 = Number of bytes reserved (same as input) R65-R67 = Address of highest byte + 1
40 41 42 43 44 45 46 4	such that the following code will store a bute into the highest addressed location of the block reserved:  STMD R65,=PTR2



RETRHI CRT



RND10 MATH

Name RNDIZ. Address 55115 Rom # @ Romisb Y Runtime code for the RANDONIZE statement.	INPUT STACK CONTENTS  Top of stack-> Optional RANDOMIZE value (8 bytes)  DUTPUT STACK CONTENTS  (empty)  R12>
9 1 2 3 4 5 6 7 18 11 12 13 14 15 16 17 28 21 22 23 24 25 26 27 39 31 32 33 34 35 36 37 48 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTR PTR2	NOTE: If no parameter is passed to this routine then the contents of R12 and the contents of the top of stack must be equal. If a parameter is passed then R12 must have been stored into the top of stack before the parameter was pushed onto the stack. In other words, the top of stack must be pointing to the first byte of the parameter.

RNDIZ. MATH



#### INPUT CONDITIONS

ROMFL = Reason for the call:

\* 0 Power on Reset Scratch 3 Loadbin 4 Run, Init 5 Load 6 Stop,Pause 7 Chain / Uhain 18 Allocate class >56 11 De-allocate class >56 12 De-compile class >56 13 Program halt on error

NOTE: ROMINI falls through into BPINI. Binary programs must insure that R0 does not get destroyed during their INIT routine as R0 is used as a counter by BPINI,

#### ROMJSB MISC.



INPUT CONDITIONS

INPUT CONDITIONS

Calls to ROMJSB must be like this:

JSR =ROMJSB

DEF routine name

BYT rom# of destination routine

ROMJSB will use the RTN address (on the

R6 stack from the 'JSB =ROMJSB') to fetch

the address and rom# you want to call.

When control returns, it will be to the

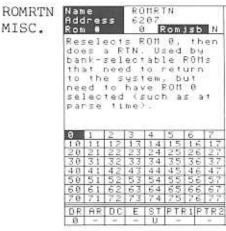
next intstruction after the 'BYT rom#'.

OUTPUT CONDITIONS

The first four butes of ERTEMP are

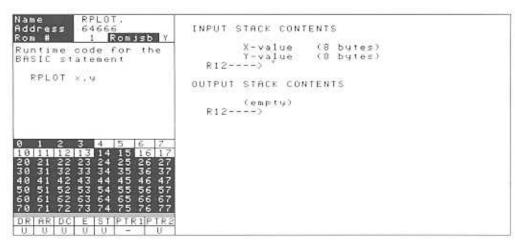
JTPUT CONDITIONS
The first four bytes of ERTEMP are destroyed by ROMJSB. The DRP=65 and the ARP=8 when the destination routine is reached. When control returns from ROMJSB to your calling routine, DRP, ARP, E. status, DCM, and the EMC PTRs are set according to the routine that was called. RO-R1 are saved on the R6 stack along with the number of the ROM that was selected when the call was initiated. They are restored before ROMJSB returns. Other registers are destroyed according to the routine that was called.

# MISC.

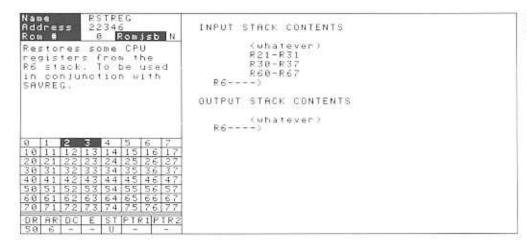


The actual code for ROMRIN is:

ROMRTN CLB RO STBD R0,=RSELEC



RPLOT. CRT



RSTREG MISC.

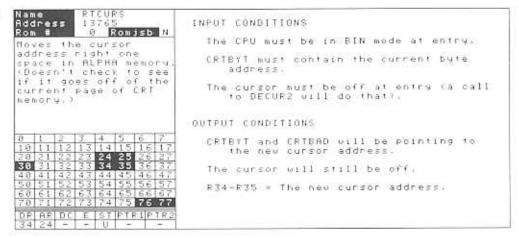
Name RSUMSK Rddress 37670 Rom # 0 Romisb N Used by ROMs perform a checksum on themselves to insure that they haven't gone bad.	INPUT CONDITIONS  The CPU must be in BIN mode at entry. The last two bytes of the ROM  (addresses 777776 and 77777) must be the checksum.  R32-R33 = Base address of rom (this will be 60000 for bank-selectable roms).  OUTPUT CONDITIONS  Upon exit, the zero flag is set if the checksum was good, else it is cleared. The actual code for RSUMSK is: PSUMSK LDM R34,=377,017 t 8K/2 + 1
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 48 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 0R HR DC E ST FTRIPTE2 46 32 *	CLM P40 POND P36,+R32 ADM R44,R36 DCM R34 JNZ RSUM ADM R46,R44 NCM R46 CMMD R46,R32 RTN

RSUM8K MISC.

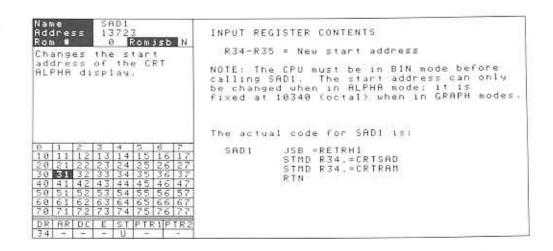
# RTCUR. CRT

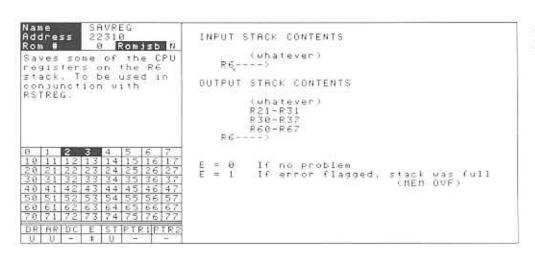
Name ATCUR Address 13651 Rom # 0 Romjsb N  Noves the cursor right one space on the ALPHA display. (Checks to see if it goes off of the current page of CRT memory and wraps it around if it does.)	INPUT CONDITIONS  The CPU must be in BIN mode at entry.  CRIBYT must contain the current byte address.  The cursor must be off at entry (a call to DECUR2 will do that).
8 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 26 21 22 23 24 25 26 27 38 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR RP DC E ST PTRIPTR2	OUTPUT CONDITIONS  CRIBIT and CRIBAD will be pointing to the new cursor address.  The cursor will still be off.

## RTCURS CRT



# SAD1 CRT





SAVREG MISC.

```
INPUT CONDITIONS

R20 = Next char, from input stream

R10-R11 = Pointer to input stream

OUTPUT CONDITIONS
Address
Rom #
                21110
0 Romisb Y
Gets the next token
from the input stream
                                                 R10-F11 =
                                                                   Pointer to input stream
                                                R41-R42 = R(N# (14 R42=0)

or binary program base address

(if R42#0)

R43 = R0M or binary program token #

or Type if variable

R44-R46 = If variable, R44-R45 = pointer

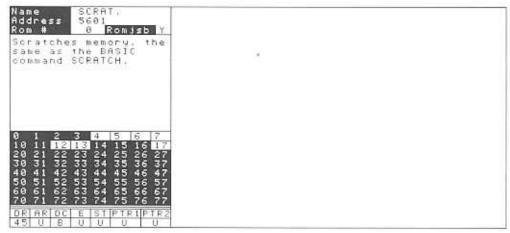
to name and R46 = length

of name
                13 14 15 1
                   3 24 25 26
3 34 35 36
                                                                 or integer value
                                                                or secondary attributes for
functions
 60 61 62 63 64 65 66 67
70 71 72 73 74 75 76 77
                                                 R47
                                                              = Class (primary attribute)
DR AR DC E ST PTR 1 PTR 2
```

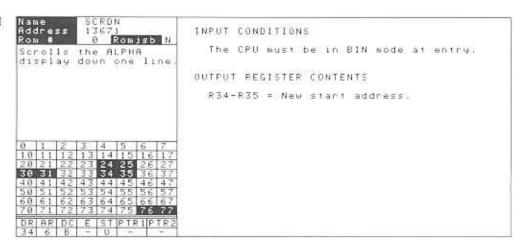
SCAN PARSE

SCAN+ PARSE

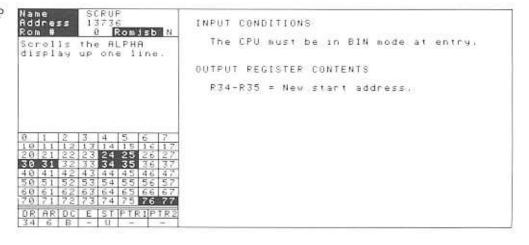
## SCRAT. MISC.

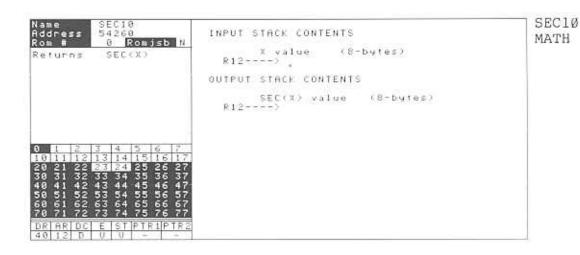


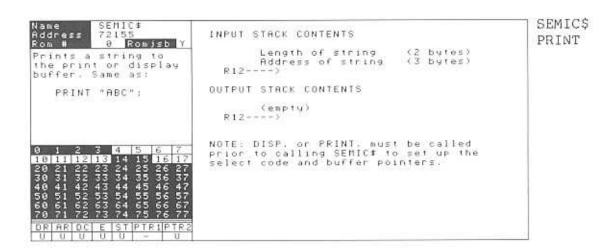
# SCRDN

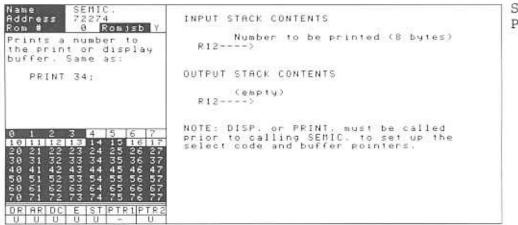


# SCRUP CRT



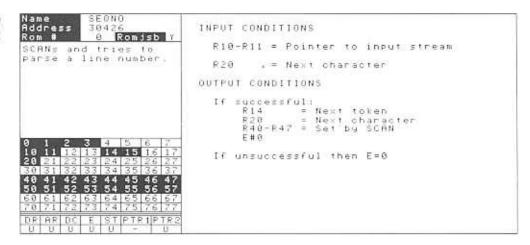




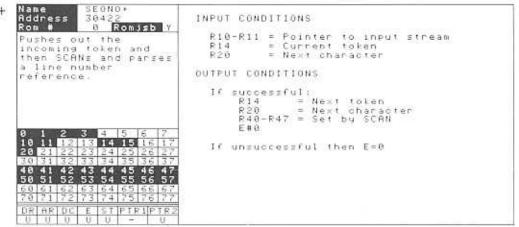


SEMIC. PRINT

### SEQNO PARSE



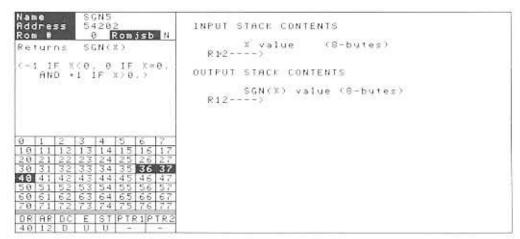
#### SEQNO+ PARSE



# SET240 MISC.

Name Addr Rom	ess	2.1	T2-		jsb	N	The	actual	co	de is:	
Sets bits		and	a t		nea n P			SET24	0	PUBB R36,+R6 LDB R36,=240 ORB R17,R36 POBD R36,-R6 RTN	
0 1	. 2	3	4	5	6	7.					
20 3	1 2 2	-	24	15	16	17					
30 3		33	34	3.5	36	37					
40 4				45	46	4.7					
	1 52		54		56	57					
	1 6	63	64								
70 7	1 6 4	16.5	1.4	75		77					
	R D	_			21 P						
36	6	-	U	10.		-					

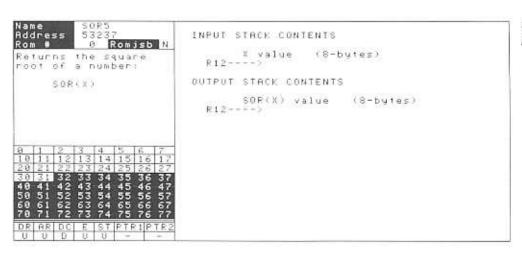
Section 8: Reference Material



SGN5 MATH

Name SIN16 Address 54343 Rom # @ Romisb N Returns the SIN(X)	INPUT STACK CONTENTS  X value (8-bytes)  R12)
	OUTPUT STACK CONTENTS  SIN(X) value (8-bytes) R12>
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PTR1 PTR2 48 12 D U U	

SIN10



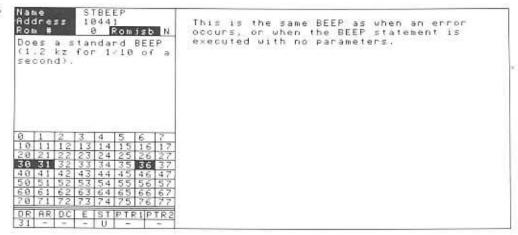
SQR5 MATH

## Section 8: Reference Material

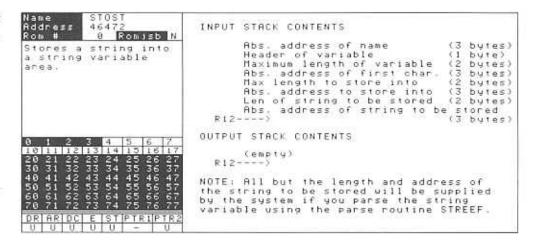
# ST240+ MISC.

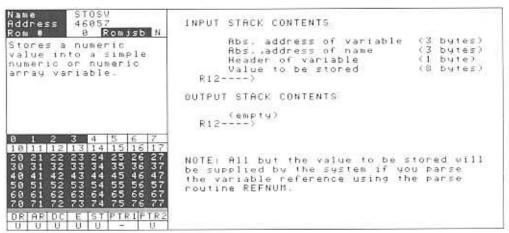
ST240+ SET240	CLB R16 PUBD R36,+R6 LDB R36,=246 ORB R17.R36 POBD R36,-R6 RTN	

# STBEEP MISC.

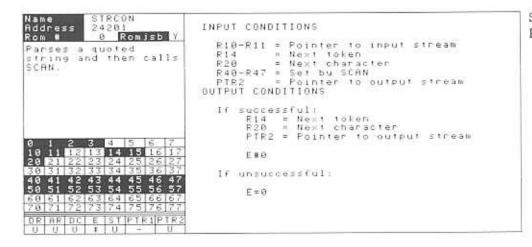


# STOST MISC.





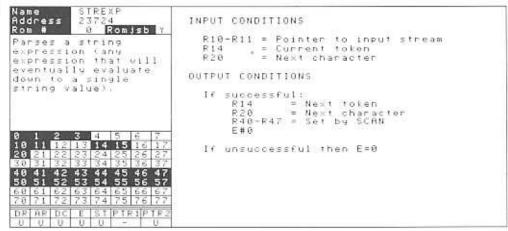
STOSV MISC.



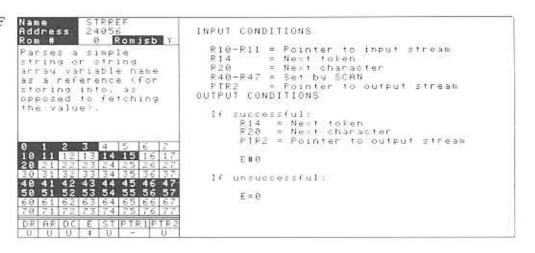
STRCON PARSE

STREX+ PARSE

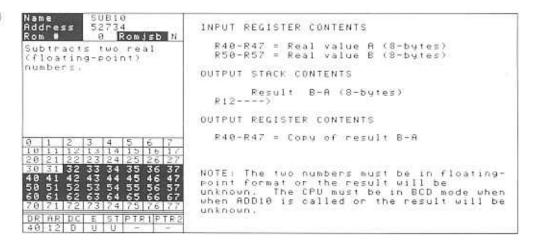
# STREXP

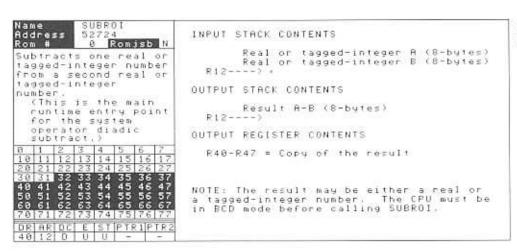


#### STRREF PARSE

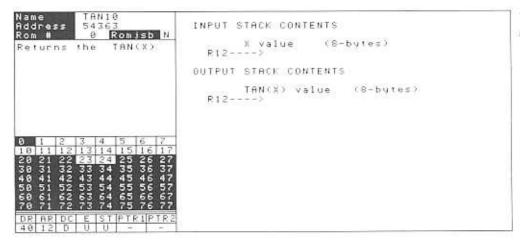


# SUB10 MATH





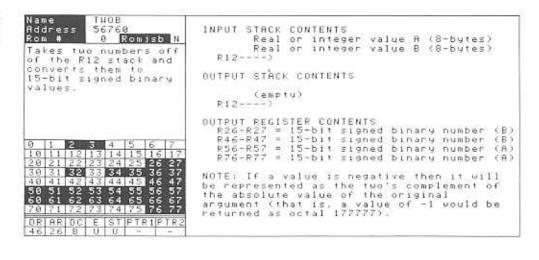
SUBROI MATH



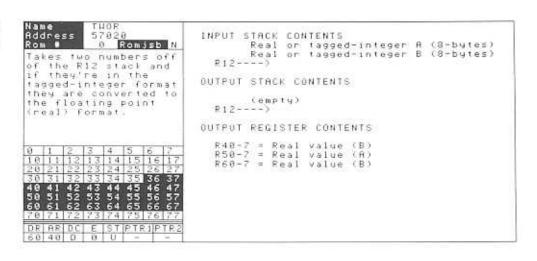
TAN10

lame TIME. Address 66211 Com # 8 Romisb V	INPUT STACK CONTENTS	
Runtime code for the system function TIME.	R12)	
	OUTPUT STACK CONTENTS  (uhstever) Time (8-bytes) R12> OUTPUT REGISTER CONTENTS	
3 1 2 3 4 5 6 7 10 11 12 12 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 50 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77	R48-R47 = Copy of time	
DR AR DC E ST PTR1 PTR2 40 12 D U U		

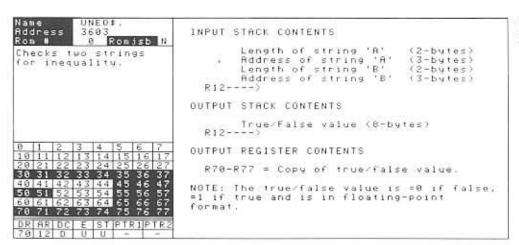
TIME. MISC. TWOB MATH



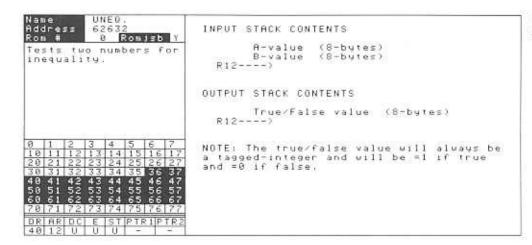
TWOR MATH



TWORO1 MATH



UNEQ\$.



UNEQ. MATH

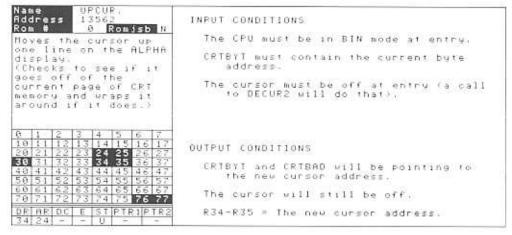
```
UNQUOT
                                                INPUT CONDITIONS
Address 24366
Rom # 8 Romisb Y
Parses an unquoted
string and then calls
SCAN, Unquoted
                                                    R10-R11 = Pointer to input stream
                                                    RIG-RII = Pointer to input stream
RI4 = Next token
R20 = Next character
R40-R47 = Set by SCAN
PIR2 = Pointer to output stream
strings are
terminated by a comma
(as in a DATA state-
ment).
                                                OUTPUT CONDITIONS
                                                     If successful:
                                                            R14 = Next token
R20 = Next character
PTR2 = Pointer to output stream
0 1 2 3 4 5 6
10 11 12 13 14 15 16
26 21 22 23 24 25 26
30 31 32 33 34 35 36
                                                           五井府
                                                    If unsuccessful:
           42 43 44
52 53 54
                                                           F = 0
60 61 62 63 64 65 66
70 71 72 73 74 75 76
DR AR DC E ST PTR1PTR2
```

UNQUOT PARSE

# UPC\$. MISC.

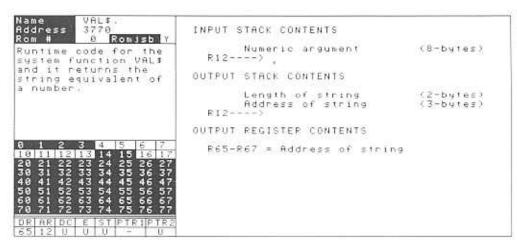
Name Address 4142 Rom # 0 Romish N Runtime code for the system function UPCs and it forces all alpha characters in a string to upper case.	INPUT STACK CONTENTS  Length of argument string (2 bytes) Address of argument string (3 bytes) R12> OUTPUT STACK CONTENTS  Length of result string (2 bytes) Address of result string (3 bytes) R12>
0 1 2 3 4 5 6 7 10 11 12 13 14 15 16 17 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40 41 42 43 44 45 46 47 50 51 52 53 54 55 56 57 60 61 62 63 64 65 66 67 70 71 72 73 74 75 76 77 DR AR DC E ST PIRIPIEZ	NOTE: The length of the result string will be the same as the length of the argument string but the addresses will be different.

# UPCUR. CRT

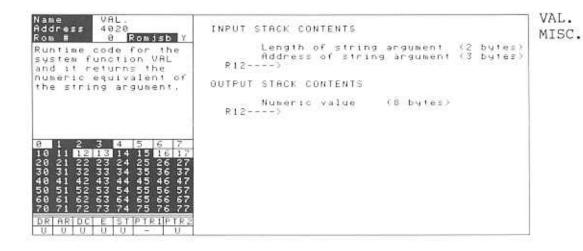


# UPCURS CRT

	ne dre n #	55		700 174		sb	N N	INPUT CONDITIONS
ad in (D if	dre AL oes it	## PHP n'1	up ch es pag	on mo nec of	sor e 1 ry. k : f of	o s f t	ee he	The CPU must be in BIN mode at entry.  CRTBYT must contain the current byte address.  The cursor must be off at entry (a call to DECUR2 will do that).
								OUTPUT CONDITIONS
0 10 20	1 11 21	2 12 22	3 13 23	14	5 15 25	6 16 26	7 17 27	CRIRYT and CRIBAD will be pointing to the new cursor address.
0	31	32 43	33 43	44	35 45	36 46	37 47	The cursor will still be off.
50 50 70	51 61 71	52 62 72	53 63 73	54 64 74	65 75	56 66 76		R34-R35 = The new cursor address,
	8R 24		E	SY	PTR	1 P	TR2	

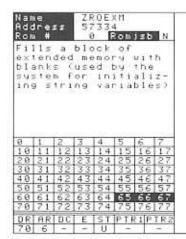


VAL\$. MISC.



YTX5 MATH

#### ZROEXM MISC.



#### INPUT REGISTER CONTENTS

R65-R67 = Number of bytes to filled.
PTR2 = Address of first word +1 of
the area to be filled with
blanks (the highest address;
this routine stores to PTR2-,
filling from highest address
to lowest).

#### OUTPUT REGISTER CONTENTS

R65-R67 = 0PTR2 = Address of first word +1.

NOTE: For filling a block of memory (in the lower 64K address space only) with blanks or zeroes, refer to ZROMEM.

# ZROMEM MISC.



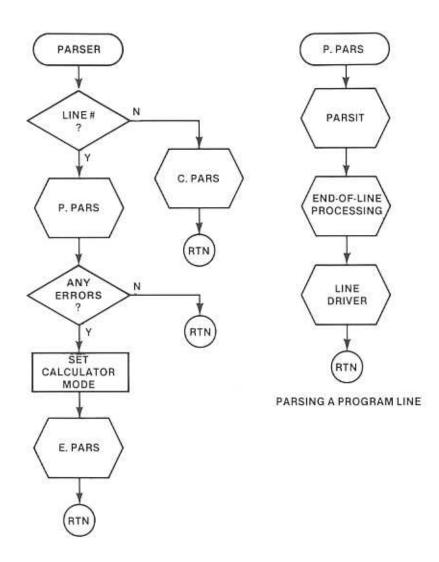
#### INPUT CONDITIONS

The CPU must be in BIN mode at entry. If R23=3 then ZROMEM will blank-fill, else it will zero-fill. R56-R57 = Number of bytes to be filled. R36-R37 = First byte to be filled (the lowest addressed byte).

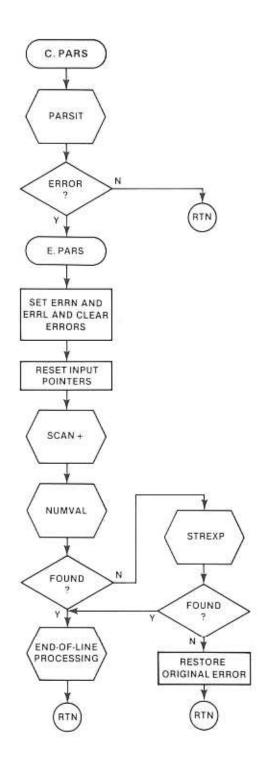
NOTE: This routine will only work in the lower 64K address space. There is another routine called ZROEXM that will blank-fill blocks of extended memory but it will not zero-fill.

# 8.4 Parsing Flow Diagrams

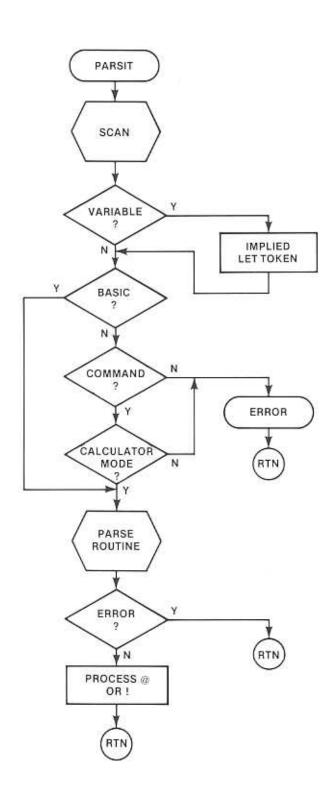
Main Parse Loop



Parsing a Calculator Mode Statement

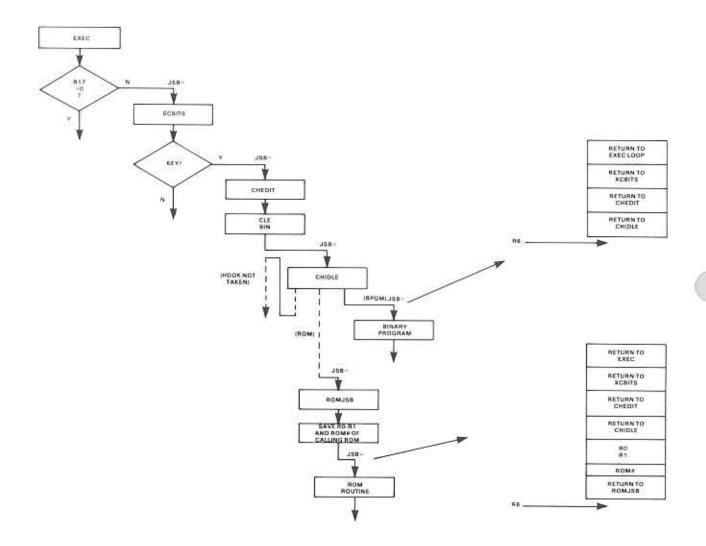


Parsit Routine

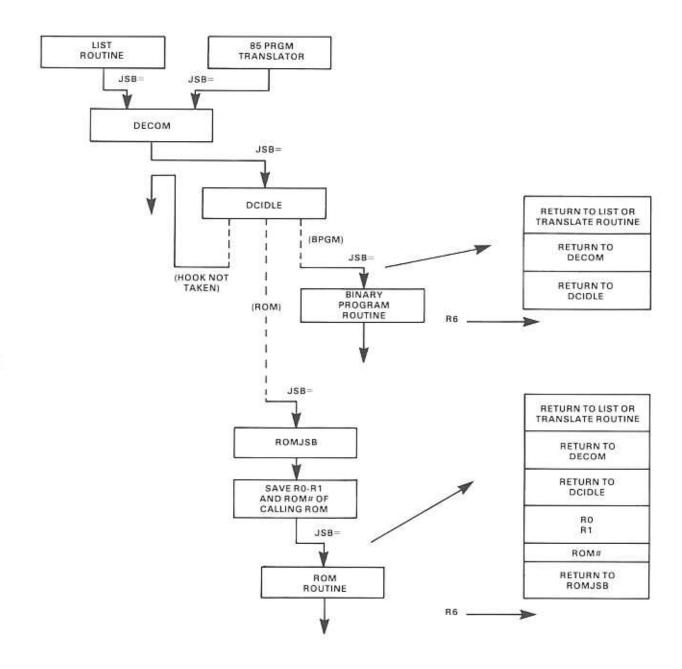


# 8.5 Hook Flowcharts

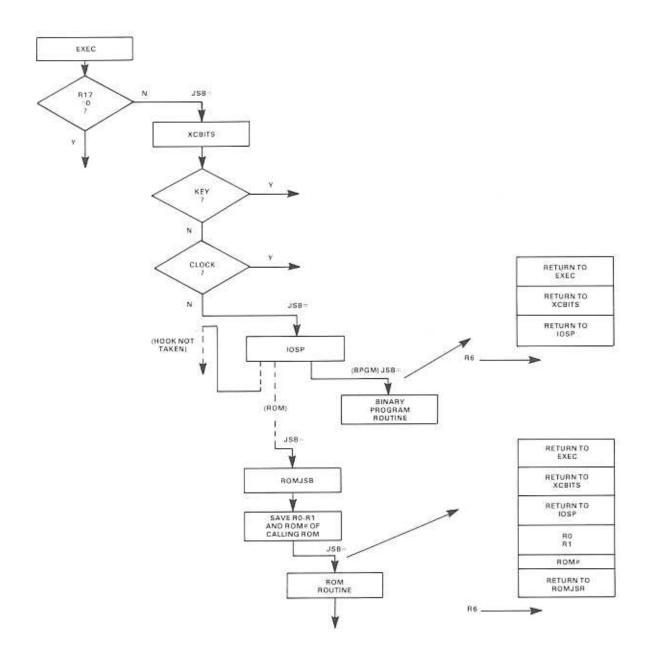
#### CHIDLE



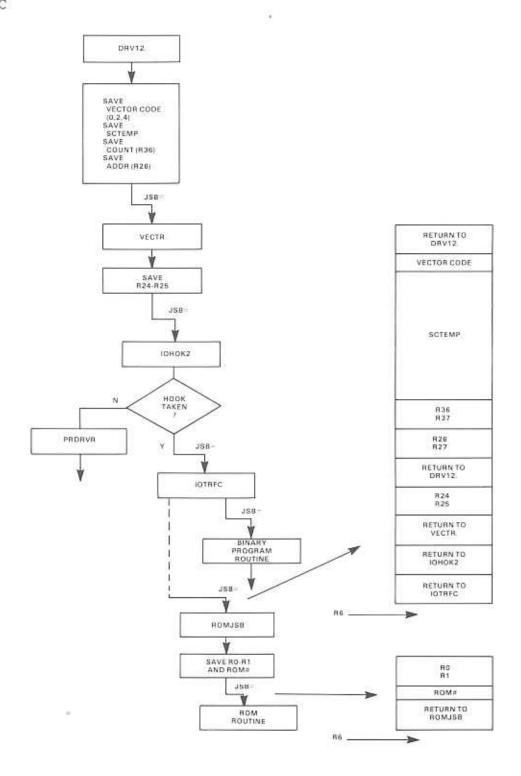
# DCIDLE



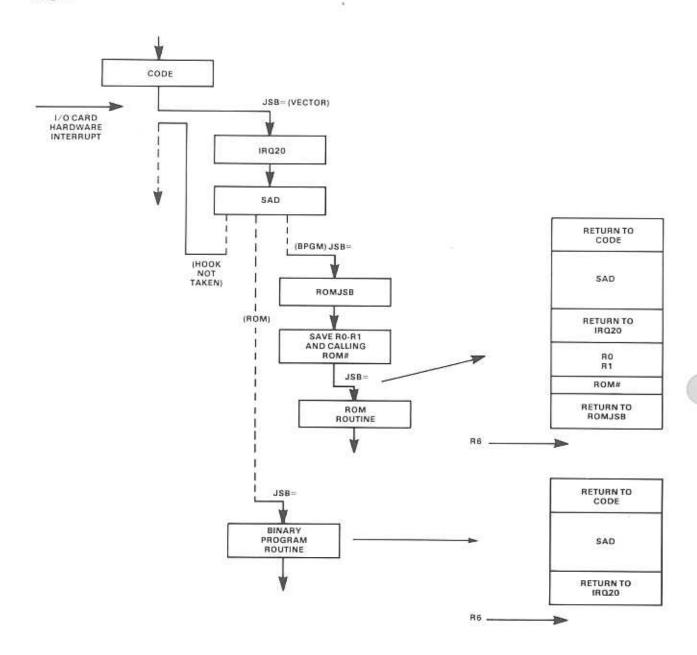
IOSP



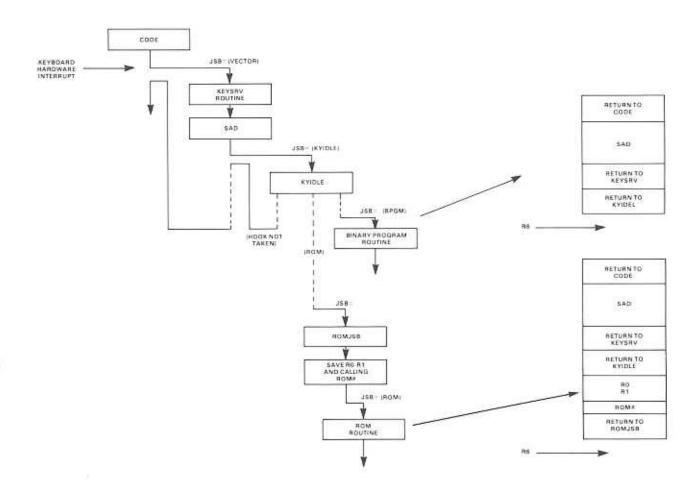
# IOTRFC



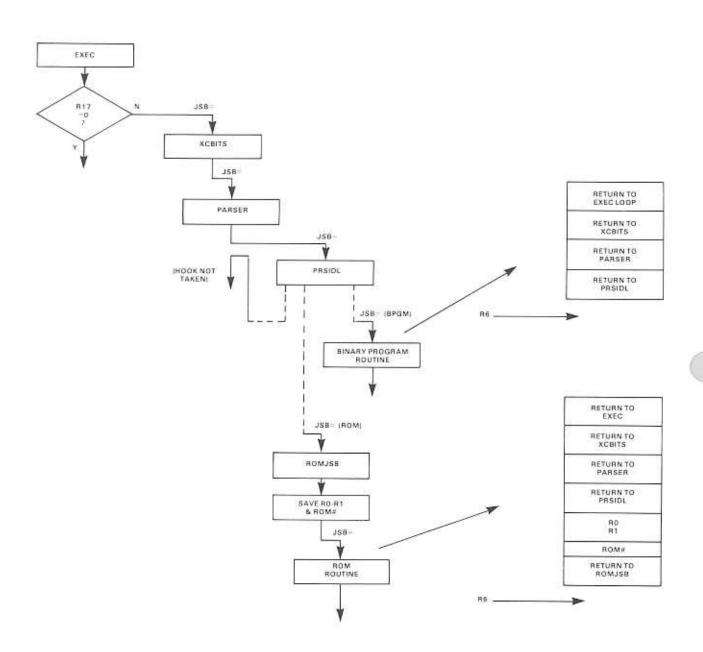
IRQ20



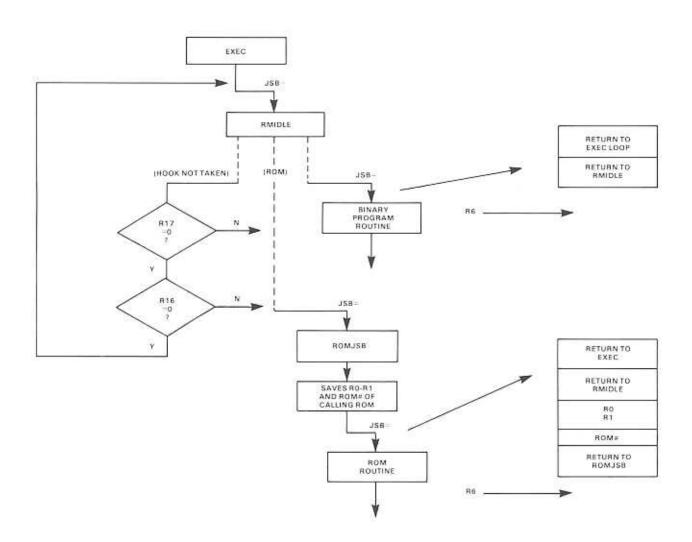
# KYIDLE



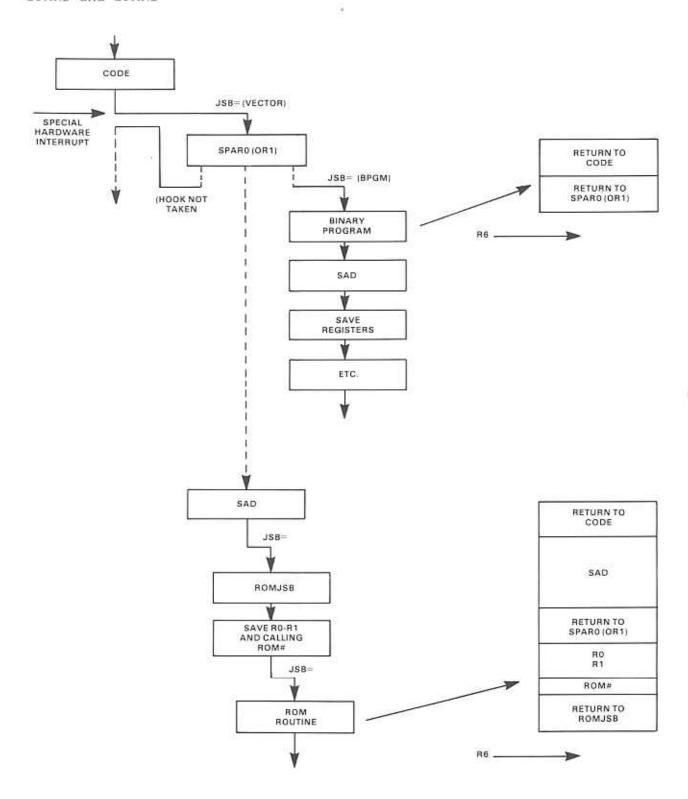
#### PRSIDL



#### RMIDLE



# SPARØ and SPAR1



# SYSTEM RUNTIME TABLE

	ROUTINE	NAME -	TOKEN	ATTRIBUTES
BTAB.R	DEF ERRORX	ERROR	0	0,44
	DEF FTSVL	SNV	1	0,1
	DEF SVADR	SAV	2	0,1
	DEF FTSTL	STRVAR	3	0,1
	DEF ICONST	REAL CONST	4	0.4
	DEF SCONST	"QUOTED STR	5	0.5
	DEF SCONST	UNQUOT STR	6	0,5
		STO. STRING	7	0,31
	DEF STOSY	STORE SV	10	0,31
	DEF AVADR1	1-DIM ADR	1.1	0.32
	DEF AVADR2	2-DIM ADR	12	0,32
	DEF AVVAL1	1-DIM VALUE	13	0,32
	DEF AVVAL2	2-DIM VALUE	14	0.32
		CARRIAGE RTN	15	0.44
	DEF GORTN	ENDSTMT	16	0.0
	DEF ERRORX	DUMMY	17	0.44
	DEF ERRORX	DUMMY	20	0.44
	DEF FTADR	SNV ADR	21	0,3
	DEF SVADR+		22	0.3
	DEF FTSTLS		23	0,3
	DEF STOSVM		24	0,43
	DEF STOSTM	MULTI STOS	25	0.43
	DEF FNCAL	FUNCTION CL	26	0,6
	DEF FNCALS	STR FUNC CL	27	0,6
	DEF JTRUE#		30	0.7
	DEF ERRORE		31	0,44
	DEF INTCON		32	0,2
	DEF JFALSR	JMP FALSE	33	0,11
	DEF JMPREL	JMP REL	34	0,11
	DEF SUBST1		35	0,34
	DEF SUBST2	2 DIM SUBST	36	0,34
	DEF EJMP#		37	0,34
	DEF FTSTA		40	0,25
	DEF JMPLB	THEN LABEL	41	
	DEF P#ARAY			0,207
	DEF EJMPLB	Array PRINT#	42	0,36
		ELSE LABEL	43	0,225
	DEF R#ARAY	Array READ#	44	0,44
	DEF ERRORX	a concer	45	0,44
	DEF CONCA.	& CONCAT	46	7,53
	DEF NOP47.	V	47	0,42
	DEF ERRORX	(	50	0,44
	DEF ERRORX	)	51	0,44
	DEF MPYROI	\$	52	12,51
	DEF ADDROI	+	53	7,51
	DEF ERRORX	. Divinio	54	0.44
	DEF SUBROI	- DIADIC	55	7,51
	DEF ERRORX	<i>t</i> 0	56	0.44
	DEF DIV2	/	57	12,51
	DEF YTX5	^	60	14,51
	DEF UNEQS.	#	61	6,53
	DEF LEQ\$.	<=	62	6,53
	DEF GEQ\$.	>=	63	6,53

ROUTINE	NAME	TOKEN	ATTRIBUTES
DEF UNEQ\$.	< ,	64	6,53
DEF EQ\$.	= '	65	6,53
DEF GRS.	>	66	6,53
DEF LTS.	<	67	6,53
DEF CHSROI	- MONADIC	70	7,50
DEF UNEQ.	#	71	6,51
DEF LEQ.	<=	72	6,51
DEF GEQ.	>=	73	6,51
DEF UNEQ.	$\Diamond$	74	6,51
DEF EQ.	=	75	6.51
DEF GR.	>	76	6.51
DEF LT.	<	77	6,51
DEF ATSIGN	@	100	0.42
DEF ONERR	ON ERROR	101	0.241
DEF OFFER.	OFF ERROR	102	0.241
	ON KEY#	103	0.241
	OFF KEY#	104	0.241
DEF AUTO.	AUTO	105	0.141
DEF BEEP.	BEEP	106	0.241
1911 001 de la company de la c	CLEAR	107	0.241
DEF CONTI.	CONT	110	0.141
DEF ONTIM.	ON TIMER#	111	0.241
DEF INIT.	INIT	112	0.141
DEF LIST.	LIST	113	0.241
DEF BPLOT.		114	0,241
DEF STIME.	BPLOT SETTIME	115	0,241
DEF CHAIN.	CHAIN	116	0.241
DEF SECUR.	SECTION	117	0.241
	READ#	120	0.241
DEF RENAM.	RENAME	121	0.241
DEF ALPHA	ALPHA	122	0.241
DEF CRT.	CRT IS	123	0.241
DEF RUN.	RUN	124	0.141
DEF DEG.	DEG	125	0,741
	37-7-7	126	0,241
DEF GCLR.	DISP	127	0,241
DEF SCRAT.	GCLEAR SCRATCH	130	0.141
DEF BEFA	DECAULT ON		100000000000000000000000000000000000000
DEF DEFAT.	DEFAULT UN	131	0.241
DEF JMPLN#		132	0,210
DEF JMPSUB		133	0,210
	PRINT #	134	0.241
DEF GRAD.	GRAD	135	0,241
DEF GRAPH.	GRAPH	136	0.241
DEF INPUT.	INPUT	137	0,241
DEF IDRAW.	IDRAW	140	0,241
DEF FNLET.	LET FN	141	0.217
DEF NOP.	LET DOINT ALL	142	0.241
DEF PRALL.	PRINT ALL	143	0,241
DEF CAT.	CAT	144	0,241
DEF DRAW	DRAW	145	0.241
DEF ON.	ON	146	0,230
DEF LABEL	LABEL	147	0.241
DEF WAIT,	WAIT	150	0.241

ROUTINE	NAME	TOKEN	ATTRIBUTES	
DEF PLOT.	PLOT .	151	0,241	
DEF PRNTR.	PRINTER IS	152	0.241	
DEF PRINT.	PRINT	153	0,241	
DEF RAD.	RAD	154	0,241	
DEF RNDIZ.	RANDOMIZE	155	0,241	
DEF READ.	READ	156	0.241	
DEF STORE	STORE BIN	157	0.241	
	RESTORE	160	0.241	
DEF RETRN.	RETURN	161	0,241	
DEF OFTIM.	OFF TIMER#	162	0.241	
DEF MOVE.	MOVE	163	0,241	
DEF FLIP.	FLIP	164	0,241	
DEF STOP.	STOP	165		
DEF STORE		166	142500000	
DEF PENUP.	PENUP	167	0.241	
DEF TROVB.	TRACE VRBL	170	0,241	
DEF TRCAL	TRACE ALL	171	0,241	
DEF XAXIS.	XAXIS	172	0.241	
DEF YAXIS.	YAXIS	173		
DEF COPY.	COPY	174		
DEF NORMA.		175		
DEF ERAST.	ERASE TAPE	176	0,241	
DEF INTEG.	INTEGER	177	0.323	
DEF SHORT.	SHORT	200	0.322	
	DELETE	201	0.141	
DEF SCALE	SCALE	202	(2.5 in 12 i	
DEF SKIP!	BEMARK	203	0.241	
DEF OPTIO.	REMARK OPTION BASE	204	0.315	
DEF COM.	COM	205	0.324	
DEF SKIPEM		206	0.320	
DEF DEFFN.		207	0.312	
DEF DIM.	DIM	210	0.321	
DEF KEYLA.	KEY LABEL	211	0,241	
DEF STOP	END	212	0.241	
DEF ENRTN	FN END	213	0.313	
DEF FOR.	FOR	214	0.341	
DEF ERRORT	1000000	215	0.344	
DEF SKIPIT		216	0.341	
DEF NEXT.	NEXT	217	0,341	
DEF UNSEC.	UNSECURE	220	0.141	
DEF ERRORT	LET (IMPLY)	221	0.244	
DEF ASIGN.	ASSIGN	222	0,244	
DEF CREAT.	CREATE	223	0,241	
DEF PURGE.	PURGE	224	0,241	
DEF REWIN.	REWIND	225	0,241	
DEF LOADB.	LOADBIN	226	0,241	
DEF PAUSE.	PAUSE	227	0,241	
DEF LOAD.	LOAD	230	0,141	
DEF REAL.	REAL	231	0.141	
DEF RENUM.	REN	232	0,141	
DEF SKIPI	1	233	0,141	
DEF DEFA	DEFAULT OFF	234	0,241	
DEF PEN.	PEN	235	0,241	

ROUTINE	NAME	TOKEN	ATTRIBUTES	
DEF PLIST.	PLIST	236	0,241	
DEF LOIR.	LDIR "	237	0,241	
DEF IMOVE.	IMOVE	240	0,241	
DEF FNLET.	FN ILET	241	0,217	
DEF CTAPE.	CTAPE	242	0,241	
DEF TRACE.	TRACE	243	0.241	
DEF TO.	TO	244	0,41	
DEF OR.	OR	245	2,51	
DEF MAX10	MAX	246	40,55	
DEF TIME.	TIME	247	0.55	
DEF DATE	DATE	250	0.55	
DEF FP5	FP	251	20,55	
DEF IP5	IP	252	20,55	
DEF EPS10	EPSILON	253	0,55	
DEF REM10	RMD	254	40.55	
DEF CEIL10	CEIL	255	20,55	
DEF ATN2.	ATN(X/Y)	256	40,55	
DEF SKPLBL	STMT LABEL	257	0.3	
DEF SQR5	SOR	260	20,55	
DEF MIN10	MIN	261	40,55	
DEF GTOLBL	GOTO LABEL	262	0,210	
DEF ABS5	ABS	263	20,55	
DEF ICOS	ACS	264	20,55	
DEF ISIN	ASN	265	20,55	
DEF ITAN	ATN	266	20,55	
DEF SGN5	SGN	267	20,55	
DEF GSUB.	GOSUB LABEL	270		
			0,210	
DEF COTTO	COT	271	20,55	
DEF CSEC10		272	20,55	
DEF FTADR3	1-D ST ARAY	273	0,1	
DEF EXP5	EXP	274	20,55	
DEF INTS	INT	275	20,55	
DEF LOGTS	LGT (10)	276	20,55	
DEF LN5	LOG (E)	277	20,55	
DEF FTADR4	2-D ST ARAY	300	0,1	
DEF SECTO	SEC	301	20,55	
DEF CHR\$.	CHR\$	302	20,56	
DEF VALS.	VALS	303	20,56	
DEF LEN.	LEN	304	30,55	
DEF NUM.	NUM	305	30,55	
DEF VAL.	VAL	306	30,55	
DEF INF10	INF	307	0,55	
DEF RND10	RND	310	0,55	
DEF PI10	PI	311	0,55	
DEF UPC\$.	UPC\$	312	30,56	
DEF USING.	USING	313	0,341	
DEF ERRORX	THEN	314	0.44	
DEF TAB.	TAB	315	20,45	
DEF STEP.	STEP	316	0,41	
DEF EXOR.	EXOR	317	2,51	
DEF NOT.	NOT	320	7,50	
DEF INTDIV	DIV (\)	321	12,51	
DEF ERNUM.	ERRN	322	0,55	

ROUTINE	NAME	TOKEN	ATTRIBUTES
DEF ERRL.	ERAL	323	0,55
DEF RESET.	RESET "	324	0,44
DEF AND	AND	325	4,51
DEF MOD10	MOD	326	12,51
DEF ERRORX	ELSE	327	0,44
DEF SIN10	SIN	330	20,55
DEF COSTO	cos	331	20,55
DEF TAN10	TAN	332	20.55
DEF NOP2.	TO (ASSIGN)	333	77,51
DEF RSTO	RESTORE LN	334	0.227
DEF RESTL.	RESTORE LBL	335	0,227
DEF ERRORX	Parameter Control	336	0,44
DEF ERRORX	j	337	0.44
DEF INTDIV	i .	340	12,51
DEF POS.	POS	341	52,55
DEF DEG10	BTD	342	20.55
DEF RAD10	DTR	343	20,55
DEF INTS	FLOOR	344	20,55
DEF USINL	USING LABEL	345	0.327
DEF READN.	READ (NUM)	346	0,44
DEF ULIN#.	USING LINE #	347	0.327
DEF INPUN.	INP NUMERIC	350	0,33
DEF INPU\$.	INP STRING	351	0,33
DEF FNRET.	LET FN(::=)	352	0.16
DEF READS.	READS	353	0.44
DEF PRLINE	PRINT END	354	0,35
DEF SEMIC.	PRINT;	355	0.36
DEF COMMA.	PRINT,	356	0.36
DEF SEMICS	PRINT;\$	357	0,36
DEF COMMAS	PRINT, \$	360	0,36
DEF ERRORX	DUMMY	361	0,241
DEF STEPK.	STEP KEY	362	0,241
DEF FTADR1	1-D NUM ARY	363	0.1
DEF FTADR2	2-D NUM ARY	364	0,1
DEF TEST.	TEST KEY	365	0,341
DEF ERRORX	DUMMY	366	0,44
DEF INDEN.	INDENTATION	367	0.2
DEF ROM:GO	EXTERNAL ROM	370	0.214
DEF BP:GO	BINARY PROG	371	0,214
DEF ERRORX	DUMMY	372	0.44
DEF ERRORX	DUMMY	373	0,44
DEF ERRORX	DUMMY	374	0.44
DEF ERRORX	DUMMY	375	0,44
DEF ERRORX	DUMMY	376	0,44
DEF ERRORX	DUMMY	377	0,44

Runtime Table/Tokens and Attributes for Graphics ROM #1

	ROUTINE	NAME	TOKEN	ATTRIBUTES
RUNTAB	DEF INIT	DUMMY # 0	0	
	DEF PLOTR.	PLOTTER IS	1	241
	DEF PRNTR.	PRINTER IS	2	241
	DEF CRT.	CRT IS	3	241
	DEF LIMIT.	LIMIT	4	241
	DEF GCLR.	GCLEAR	5	241
	DEF LOCAT.	LOCATE	6	241
	DEF BPLOT.	BPLOT	7	241
		SCALE	10	241
	DEF SHOW.		11	241
		MSCALE	12	241
	DEF CLIP.	CLIP	13	241
	DEF UNCLI.	UNCLIP	14	241
	DEF SETGU. DEF SETUU.	SETGU	15	241
			16	241
	DEF PENUP.	PENUP	17	241
	DEF GREAD.	BREAD	20	241
	DEF PEN.	PEN	21	241
	DEF LINET.	LINETYPE	22	241
	DEF PLOT.	PLOT	23	241
	DEF IPLOT.	IPLOT	24	241
	DEF MOVE.	MOVE	25	241
	DEF IMOVE.	IMOVE	26	241
	DEF DRAW.	DRAW	27	241
	DEF IDRAW.	IDRAW	30	241
	DEF RPLOT.	RPLOT	31	241
	DEF POIR.	PDIR	32	241
	DEF BLOFF	NOBLINK	33	241
	DEF AXES.	AXES	34	241
	DEF LAXES.	LAXES	35	241 -
	DEF GRID.	GRID	36	241
	DEF FRAME.	FRAME	37	241
	DEF LABEL	LABEL	40	241
	DEF BLINK	BLINK	41	241
	DEF LORG.	LORG	42	241
	DEF LDIR.	LDIR	43	241
	DEF CSIZE.	CSIZE	44	241
	DEF WHERE	WHERE	45	241
			16.20	
	DEF CONTR.	CONTROL	46	241
	DEF CURSR.	CURSOR	47	241
	DEF DIGIT.	DIGITIZE	50	241
	DEF DUMMY	TRANSLATE	51	241
	DEF LGRID.	LGRID	52	241
	DEF GRAPH.	GRAPHICS	53	241
	DEF XAXIS.	XAXIS	54	241
	DEF YAXIS.	YAXIS	55	241
	DEF FXD.	FXD	56	241
	DEF ERRSC.	ERRSC	57	0,55
	DEF ERROM.	ERROM	60	0,55
	DEF RATIO.	RATIO	61	0.55
	DEF TAB,	TAB	62	20,45
	DEF LABEOL	LABEL EOLINE	63	35
	DEF PAGES.	PAGE SIZE	64	241
	DEF ALFAL.	ALPHA ALL	65	241
	DEF GRAFA.	GRAPH ALL	66	241
	DEF FRE.	FREE MEMORY	67	0.55

Runtime Table/Tokens and Attributes for Mass Storage ROM #320

	ROUTINE	NAME	TOKEN	ATTRIBUTES
RUNTIM	DEF INITIT	DUMMY # 0	0	241
	DEF ASSIG.	ASSIGN	1	241
	DEF MSCAT.	CAT	2	241
	DEF CHKOF.	CHECK READ OFF	3	241
	DEF CHECK.	CHECK READ	4	241
	DEF ERRORX	DUMMY ROUTINE	5	44
	DEF MSCPY.	COPY	6	241
	DEF MSCRE.	CREATE	7	241
	DEF INITI.	INITIALIZE	10	241
	DEF MSCHA.	CHAIN	11	241
	DEF MSLDB.	LOADBIN	12	241
	DEF MSLOD.	LOAD	13	141
	DEF MASSS.	MASS STORAGE IS	14	241
	DEF MSPRNT	PRINT#	15	241
	DEF ERRORX	DUMMY ROUTINE	16	44
	DEF ERRORX	DUMMY ROUTINE	17	44
	DEF MSPUR.	PURGE	20	241
	DEF READ.	READ#	21	241
	DEF MSREN.	RENAME	22	241
	DEF MSSTB.	STOREBIN	23	141
	DEF MSSTO.	STORE	24	141
	DEF PACK.	PACK	25	241
	DEF VOLUM.	VOLUME	26	241
	DEF GLOAD.	GLOAD	27	241
	DEF GSTOR.	GSTORE	30	241
	DEF ERROM.	ERROM	31	0,55
	DEF ERRSC.	ERRSC	32	0,55
	DEF TYP.	TYP	33	20,55
	DEF IS.	(VOLUME) IS	34	1.51
	DEF ERRORX	DUMMY ROUTINE	35	44
	DEF TO.	(RENAME) TO	36	1,51
	DEF RONUM.	READ# NUMERIC	37	44
	DEF PRARR.		40	36
	DEF ROSTR.		41	44
	DEF PRNUM.	PRINT# NUMERIC	42	36
	DEF PREOL	PRINT# END OF LINE	43	35
	DEF PRSTR.	PRINT# STRING	44	36
	DEF RDARR.	READ# NUM ARRAY	45	44
	DEF PRARRS	PRINT# STRING ARRAY	46	36
	DEF RDARRS	READ# STRING ARRAY	47	44

# 8.7 Error Messages

Following is a list of the error messages provided by the Assembler ROM and the system monitor. For other errors refer to the owner's manual or to the manuals for other peripherals that may be attached to the HP-87.

# Assembler System Errors

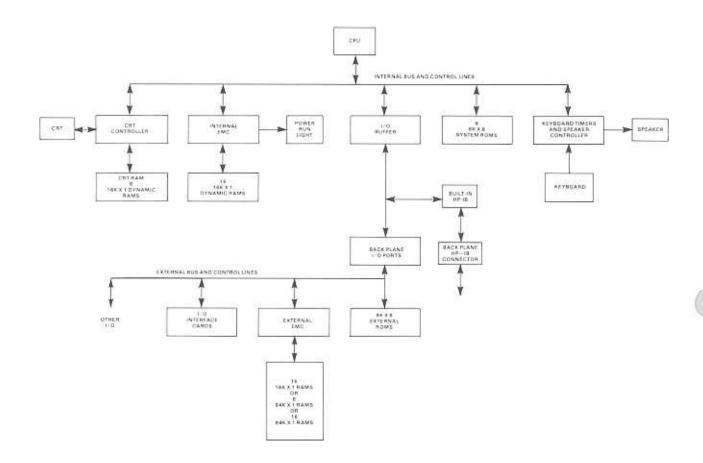
ERROR	109:	ILL MODE	A command has been executed in the wrong operating mode (that is, ASSEMBLER has been typed when the computer is already in assembler mode).
ERROR	110:	LBL	An invalid label has been seen; may have been longer than six characters or started with a digit.
ERROR	111:	OPCO	The opcode is not recognized; may have been misspelled, no space was typed between the label and the opcode, or because the opcode was entered in the first or second column after the line number.
ERROR	112:	ARP-DRP	Invalid ARP or DRP; ARPs and DRPs must be between $\emptyset$ and 77 inclusive, and cannot be 1.
ERROR	113:	OPER	Bad operand; that is, LDM R34,=3,remark. Because a number follows the equal sign in this example, the assembler expects another number after the comma. Also, each literal value must be specified with two digits if a BCD quantity.
ERROR	114:	FIN-LNK	Missing FIN or LNK statement. If the file name or file type is wrong in the LNK statement, then a "FILE NAME" or "FILE TYPE" error will be generated.
ERROR	115:	ASSM ROM	At power-on, this means the ROM had a checksum error. At a breakpoint, any errors generated give this message.

### Section 8: Reference Material

## Assembly Errors

ILL NAM	A NAM statement has already been executed, or an ABS ROM has been executed.
AIF UND	The specified conditional assembly flag has not yet been defined as set or cleared.
ILL ABS	An ABS or NAM statement has already been encountered.
JMP FROM	The jump from the spcified line is out of range.
JMP TO	The jump to the specified line is out of range.
UND LAB	After assembly was completed, this label had not been defined in the program or in the global file.
ILL GLO	The GLO statement occurs after a NAM statement, ABS statement, or another GLO statement.

### 8.8 System Hardware Diagram



#### 8.9 Assembler Instruction Set

On the following pages is a list of all CPU instructions available on the Assembler ROM.

### Legend

EA

DR Data register. Can be register number (that is, R32),  $R^*$ , or  $R^\#$ .

AR Address register. Can be register number, R\*, or R#.

Literal Literal value, up to 10 octal bytes in length. Can be BCD constant (that is, 99C), octal constant (that is, 12), or decimal constant (that is, 20D). Can also be specified by a label, where the literal quantity is a one- or two-byte value or address assigned to the label.

Label Address of literal quantity. Label name must begin with an alphabetic character, can use any combination of alphanumeric characters, and can be 1-6 characters in length.

Clock Cycle 1.6 sec.

B Number of bytes.

T Add one clock cycle if true (that is, the jump occurs).

R(x) CPU register addressed by (x).

M(x) Memory location addressed by (x) where (x) is a 16-bit address.

PC Program counter stored in CPU registers R4 and R5. Used to address the instruction being executed.

SP Subroutine stack pointer stored in CPU registers R6 and R7. Used to point to the next available location on the subroutine return address stack.

Effective address. The location from which data is read for load-type instructions or the location where data is

placed for store-type instructions.

ADR	Address. The two-byte quantity directly following an instruction that uses the literal direct, literal indirect, index direct, or index indirect addressing mode. This quantity is always an address.
n	Literal value.
-	Is transferred to.
( )	Contents of.
( <del></del>	Complement (that is, $x$ is complement of $x$ ). This is one's complement if DCM=0 and nine's complement if DCM=1.
•	Logical "and."
v	Inclusive "or."
$\oplus$	Exclusive "or."
JIF	Jump if.
1	Status bit is set.
0	Status bit is cleared.
Х	Status bit is affected.
127	Status bit is not affected.
Y	This option is available to this instruction.

Section 8: Reference Material

- Lorenzo - Const			10.1		200 - 17-20-2					Sta	tus						Binar
Instruction Format	Description	Addressing Mode	Opcode	Clock	Operation			ROZ			- 2	CM:	90 <b>7</b> -0		CM:		BCE Optio
Later Control	220227280	2500000	22000	-		LSB	MSB	LDZ	Z	DCM	E	CY	OVF	E	CY	OVF	100
ADB DR, AR	Add byte	Reg. imm.	302	5	DR + RD → RO	Х	×	×	X		=	×	Х	7	Х	0	Y
ADB DR. =	Add byte	Lit. imm.	312	5	DR + DR + MIPC + 1)	Х	Х	X	X	<i>5</i> 7.0	-	х	X		х	D	Ä
ADBD DR. AR	Add byte	Reg. dir.	332	6	DR + DR + M(AR)	Х	Х	Х	Х	-	-	Х	X	-	Х	0	
ADBD DR. = label	Add byte	Lit. dir.	322	5	DR - DR + MIADRI	Х	Х	Х	X	-	-	Х	X	-	Х	D	ुभ
ADM DR. AR	Add multi- byte	Reg. imm.	303	4 + 8	DR + DR + AR	Х	х	x	X	550	577	х	X		X	0	Y
ADM DR. =	Add multi- byte	Lit. imm.	313	4+8	DR + DR + M(PC + 1)	Х	X	X	Х	-	-	Ж	X	-	Х	0	y
ADMD <i>DR, AR</i>	Add multi- byte	Reg. dir.	333	5 + 8	OR + M(AR)	Х	×	X	X	700	-	х	X	-	X	ŋ	Y
ADMO DR; =	Add multi- byte	Lit. dir.	323	4+8	OR + DR + M(ADR)	×	Ж	X	Х	-	-	X	X	-	X	0	y
ANM DR, AR	Logical AND (multi-byte)	Reg. imm.	307	4 + B	DR + DR · AR	Х	х	X	X	-	_	D	0	_	0	0	
ANM DR. =	Logical AND (multi-byte)	Lit. imm.	317:	4 + B	DR - DR · M(PC + 1)	×	ж	X	х	-	-	П	0	-	0	0	
ANMO DR. AR	Logical AND (multi-byte)	Rey Dir	337	5 + B	DR + DR - M(AR)	х	Х	х	X	7		D	0	_	0	0	
ANMO DR =	Logical AND (multi-byte)	Lit. dir.	327	5 + B	DR - OR - MIADRI	Х	х	X	X	=	-	0	0	-	ū	0	
ARP AR	Load ARP		000-077 (≠001)	2	ARP - n	=	45	_	_	3	-	_	-	_	_	-	
ARP *	Load ARP with contents of R0		001	3	ARP - RO		=	=	-	-	=	÷.	5 =	-51	=	275	٠
BCD	Set BCD mode		231	4	DCM + 1	-	-	-	-	1	-			-	-		
BIN	Set binary mode		230	4	DCM - B	-	-	=	=	0	=	=	-	=	-	=	
CLB DR	Clear byte	Reg imm	222	5	DR - 0	X	X	X	Χ	_	-	0	0	_	0	0	
CLM DR	Clear multi- byte	Reg. imm.	223	4 + B	DR ← 0	х	х	×	X	=	-	0	Ø	-	0	0	
CLE	Clear E		235	2	E + 0	_	-			-	а	-	-	0	_	-	
CMB DR, AR	Compare byte	Reg imm	300	5	OR + AR + 1	X	×	x	x	_	-	×	X	-	X	D	Y
CMB DR =	Compare byte	Lit. imm	310	5	DR + M(PC + 1) + 1	x	x	×	x	=		х	×	-	×	0	Y
literal	Conservation 1	Reg. dir.	330		DR + M(AR) + 1	· v	OV.	w.	· v			240	v			on o	\ Y
CMBD DR. AR	Compare byte		330	fi e	105	X	×	X	X	-		X	X	-	X	0	- SY
CMBD BR. = label	Compare byte	Lit. dir.	320	6	DR + MJADRJ + 1	Х	Х	Æ.	·A	-		٨	*		A .	5.M.S.S.	
CMM DR. AR	Compare multi-byte	Reg. imm.	301	4 + B	OR + AR + 1	х	Х	Х	χ			×	×	7	X	0	Y
CMM DR. = /itera/	Compare multi-byte	Lit. imm.	311	4 + B	DR + M(PC(+ 1) + 1	х	Х	X	Х	-	_	÷χ	X	-	X	0	ुर
CMMD DR AR	Compare multi-byte	Reg dir.	331	5 + 8	DR + M(AR) + 1	х	Х	X	X	77	-	X	х	-	Х	0	Y
CMMD BR = label	Compare multi-byte	Lit. dir.	321	5 + 8	DR + MJADRJ + 1	х	×	X:	Х	-	-	х	X	-	X +	Û	Y
DCB DA	Decrement byte	Reg. imm.	212	5	0R - 0R - 1	х	х	X	X	=	2	X	×	100	x	0	Y
DCM <i>DR</i>	Decrement multi-byte	Reg imm.	213	4 + 8	DR ← DR — 1	х	х	X	X	-	-	х	х	-	Х	0	Ŷ
DCE	Decrement E		233	2	€ - € - 1	=	-	-	_	-	Х	_	<u> </u>	χ	_	_	
DRP <i>DR</i>	Load DRP		100-177 (≠101)	2	DRP - n	=3	_		-	-	-	-	1000		-	-	
DRP 1	Load DRP with contents of RO		101	3	DAP - NO	77.0	77		-	=	-	770	-	=		=	

Instruction	/autovarozer	Addressing	1200010	Clock	Sugaritori con M				Sta							Binar	
Format	Description	Mode	Opcode	Cycles	Operation	LSB	MSB	RDZ LDZ	z	DCM		CY	OVF		CM =	= 1 OVF	BCD Optio
ELB OR	Extended left byte	Reg. imm.	200	5	Circulate DR =	x	Х	X	X	-	_	X	X	X	0	0	¥
ELM DR	Extended left multi-byte	Reg. imm.	201	4 + B	Circulate DR left once	x	х	X	х	-	-	X	x	х	0	0	Y
ERB DR	Extended right byte	Reg. imm	202	5	Circulate DR right once	X	Х	Χ	X	-	-	X	0	X	0	0	Y
BM 08	Extended right multi-byte	Reg imm	203	4 + 8	Circulate DR right once	X	X	Х	Х	-	-	X	0	X	0	0	Y
CB DA	Increment byte	Reg. (mm.	210	5	DR + DR + 1	X	X	Х	X	_	-	X	X		Х	0	Y
CM <i>DR</i>	Increment multi-byte	Reg. imm.	211	4 + 5	DR - DR + 1	х	х	X	X	-	-	χ	X	-	Х	0	Y
CE:	Increment E		234	2	E ≠ E + 1	: <u>-</u>	-	-	_	-	X	_	-	Х	-	-	
ICY label	Jump on carry		373	4 + T	JIF + CY = 1	0=	_	-	_	_	_	-	-	_	_	-	
IEN label	Jump on E non-zero		370	4+1	JIF E ≠ 0000	114	100	-	_	_	-	=	-	iven	-	_	
IEV /aba/	Jump on even		363	4 + T	JIF LSB = 0	-	-	-	_	-	-	-	-	-	-	040	
IEZ label	Jump on E zero		371	4 + T	JIF E = 0000	2=	tt	-	-	-	-	-	SS <del>10</del>	-	-	-	
ILN /abe/	Jump on left digit non-zero		375	4+1	JIF LDZ ≠ 1	-	#	_		_	_	-	-		-	_	
ILZ label	Jump on left digit zero		374	4 + T	JEF LDZ = 1	250	-	-	-	-	-	-	-	-	-	-	
MP (abel	Unconditional jump		360	4 + T	Jomp always	·=	77.5						-	7		77	
NC /ata/	Jump on no carry		372	4 + T	JIF CY = 0	:==	-	-	-	-	-	-	-	-	-	-	
ING label	Jump on negative		364	4 + T	JIF MSB ≠ 0VF	=	77.0	-	-	-	-		-	-		777	
INO /abe/	Jump on no overflow		361	4 + T	JIF OVF = 0	-	-	5 <del>=1</del>	-	=	-	-	-	-	-	-	
INZ label	Jump an non-zero		366	4+1	JIF Z ≠ 1	-	20		_	27				_	_		
IOD (abe)	Jump on odd		362	4 + T	JIF LSB = 1	-	-	-	-		-	-	-		-	_	
IPS label	Jump on positive		365	4+1	JIF MSB = GVF	-	-	-	-	-	-	-	_	-	-	-	
RN label	Jump en right digit non-zero		377	4+T	JIF RDZ ≠ 1	_	227	-		-	_	_	_		_	_	
RZ/ebel	Jump on right digit zero		376	4 + T	JIF #8Z = 1	-		in	-	-	-	-	_	_	-		
SB = label	Jump subrouting	Literal direct	316	g	Jump subroutine	-		-		===				=	-	-	
SB XA. label	Jump subroutine	Indexed	306	31	Jump subroutine indexed	-	= 3	-	_		-	_	_	_	-	-	
2 Riebal	Jump an zero		367	4+1	JIF Z = 1	-	-	-	-	-	_	-	_	-	-	-	
DB DR, AR	Load byte	Reg. imm.	240	5	DR + AR	X	X	X	X	-	-	D:	0	-	0	0	
DB DR = literal	Load byte	Lit. imm.	250	5	DR - M(PC + 1)	X	Х	X	Х	-	-	0	0	-	0	0	
.DBD <i>DR, AR</i>	Load byte	Reg. dir.	244	6	OR ← M(AR)	X	х	х	х	=	-	0	0	-	0	0	
DBD DR. = (abal	Load byte	Lit. dir.	260	6	OR - MIADRI	X	Х	х	Х	-	-	D	0	7	0	0	
DBD DR, XAR, label	Load byte	Index dir.	264	8	DR - M(ADR + AR)	X	Х	Х	X	_	-	0	0	-	0	0	
LDBI DR. AR	Load byte	Reg. indir.	254	8	DR + M(M(AR))	X	×	х	X	-		0	0	-	0	0	
LDBI <i>DR.</i> =  abel	Load byte	Lit. indir.	270	8	DR + M(M(ADR))	X	Х	Х	Х		-	0	0	77	0	0	

Section 8: Reference Material

Instruction		Addressing		et						Sta	tur						Binery
Format	Description	Mode	Opcode	Clock Cycles	Operation	Water.	n oreres	RDZ	10000			DCM			DCM		BCD Option
LDBI DR. XAR,	Load byte	Index indir.	274	10	DR - M(M(ADR +	LSB X	MS8	LDZ X	Z X	DCM	E -	CY	0VF	E	CY 0	OVF.	
label LDM DR. AR	Load	Reg. imm.	241	4 + B	AR)) DR + AR	X:	х	x	х	_	_	0	0		0	0	
IDM DR =	multi-byte Load	Lit. imm.	251	4 + B	OR - M(PC + 1)	×	x	X	x	_	_	0	α	_	0	0	
literel .DMD DR. AR	multi-byte Load multi-byte	Reg dir.	245	5 + B	DE - MIARI	XΞ	х	X	X	-	-	D	0		0	0	
DMD DR. =	Load multi-byte	Lrt. dir	261	5 + B	DR MIADRI	×	X	X	X		37.	D	0	77	D	a	
DMD DR, XAR,	Load multi-byte	Index dir.	265	7 + B	OR - MIADR + AR)	X	Х	Х	χ	-	-	0:	0	_	0	0	
.DMI DR. AR	Load multi-byle	Reg. indir.	255	7 + B	OR M(M(AR))	×	X	×	х	-	-	D	0	=	0	0	
DMI DR = label	Load multi-byte	Lit, indir.	271	7 + B	DR + M(M(ADR))	X:	X	X	X	-	-	0	0.0	-	:0	8	
OMI DR. XAR. (abel	Load multi-byte	Index indir.	275	9 + B	DR ← M(M(ADR + ARI)	×	×	X	X	-	_	Û	0		0	ú	
LLB <i>DR</i>	Logical left byte	Reg. imm.	204	5	Logical left shift DR	х	x	Х	X	=	-	X	X	X	0	0	Y
LM DR	Logical left multi-byte	Reg. imm.	205	4 + 8	Logical left shift DR	X	Х	Х	X	-	_	X	х	X	Ü	0	Υ
RB D#	Logical right byte	Reg imm:	206	5	Logical right shift DR	×	x	×	X	-	-	X	0	X	0	D	Y
RM <i>DR</i>	Logical right multi-byte	Reg imm	207	4 + B	Logical right shift DR	X	X	Х	Х	-	-	X	0	х	0	0	Y
NCB <i>DR</i>	Nine's (or one's) complement byte	Reg. imm.	216	5	DR DR	X	×	x	х		=	X	х		X	0	Y
NCM DR	Nine's (or one's) complement multi-byte	Reg. imm.	217	4 + B	DR - DR	X	х	X	X	-	-	X	X		Х	:0	Y
DAB DR. AR	Or byte inclusive	Reg imm	224	5	DR + DR Y AR	Х	X	χ	ЭХ	-	-	0	0	=	0	0	
IRM DR, AR	Ör multi-byte inclusive	Reg. imm.	225	4 + B	DR - DR V AR	Х	x	x	X		-	0	0		0	D	
PAD	Pop ARP, DRP and status from stack		237	8	Status - M(SP)	⊗:	X	×	ЭХ	X	-	X	X	-	X	х	
P080 DR +AR	Pop byte with post-	Stk. dir.	340	6	DR + M(AR). AR + AR + 1	х	х	х	х	-	-	0	0	-	0	D	
POBO DRAR	Pop byte with pre- decrement	Stk dir	342	6	DR - M(AR), AR - AR - 1	х	X	x	х	-	-	0	D	-	0	Ū	
POBL <i>DR</i> , +AR	Pop byte with post-	Stk. indif.	350	В	DR - M(M(AR)). AR - AR + 2	x	x	x	х			0	0	-	0	0	
POBLOR —AR	Pop byte with pre- decrement	Stk. indir.	352	8	DR + M(MJARJ). AR + AR - 2	×	X	x	Х	T	_	0	0	-	0	D	
POMD DR, +AR	Pop multi-byte with post- increment	Stk. indir.	341	5 + B	DR ← M(AR) AR ← AR + M	×	x	X	X	_	-	D	0	-	0	a	
OMO <i>DR</i> , —AR	Pop multi-byte with pre- decrament	Stk, dir.	343	5 + 8	DR MIARI, AR AR M	x	X	x	Х	-	-	0	0	-	0	0	
POMI BR. +AR	Pop multi-byte with post- increment	Stk indir	351	7 + 8	DR M(M(AR)), AR AR +- 2	x	X	X	х	-	-	0	0		0	0	

Section 8: Reference Material

Instruction	1200 CON 1000	Addressing	20000	Clock	// <u>2</u> 8/2002/20					Sta							Binary
Format	Description	Mode	Opcode	Cycles	Operation	LSB	MSB	RDZ	Z	DCM		CY	= 0 OVF		CY	= 1 OVF	9CD Option
POMI <i>OR</i> , —AR	Pop multi-byte with pre- decrement	Stk, indir.	353	7 + 8	DR + M(M(AR)), AR + AR - 2	Х	х	x	x			0	0	=	0	0	
PUBD DR. +AR	Push byte with post- increment	Stk. dir.	344	6	M(AR) = OR, AR = AR + 1	х	х	х	X	277	-	0	0	=	0	Û	
PUBD BR AR	Push byte with pre- decrement	Stk. dir.	346	6	AR + AR - 1. M(AR) + DR	Х	X	X	X	20	_	0	0		0	0	
PUBI <i>DR</i> . +AR	Posh byte with post- increment	Stk. indir.	354	8	M(M(AR)) + DR. AR + AR + 2	Х	x	Х	х	34	_	0	0	=	0	0	
PUBI <i>DR</i> , —AR	Push byte with pre- decrement	Stk. indir.	356	8	AR - AR - 2_ M(M(AR)) - DR	X	X:	X	ЭХ	: <del></del>	-	Ø:	0	-	0	0	
PUMO <i>DR</i> +AR	Push multi- byte with post- increment	Stk. dir.	345	5 + B	M(AR) - DR. AR - AR + M	х	x	x	Х	-	-	0	0	-	0	0	
PUMO <i>dr. —Ar</i> i	Push multi- byte with pre- decrement	Stk. dir.	347	5 + B	AR ← AR − M, M(AR) ← DR	Х	X	Χ	X	=	_	0	0	-	0	0	
PUMI <i>or</i> , + <i>AR</i>	Push multi- byte with post- increment	Stk. indir.	355	7 + B	M(M(AR)) + DR. AR + AR + 2	х	x	x	х	_	_	0	0	-	D	0	
PUMI <i>DR</i> ; —AR	Push multi- byle with pre- decrement	Stk. indir.	357	7 + 8	AR AR 2, M(M(ARI) DR	х	x	х	х		-	a	D	=	0	0	
RIN	Subroutine return		236	5	SP ← SP — 2, PC ← M(SP)	-	-	-	-	-	_	-	-	-		_	
SAD	Save ARP, DRP and status on stack	99	232	8	M(SP) - Status		~	_	7	-	_	_	-	_	_	-	
SBB DR, AR	Subtract byte	Reg imm	304	5	DR ← DR + AR + 1	x	X	X	X		_	X	X	_	Х	0	Y
SBB DR. =	Subtract byte	Lit. imm.	314	5	DR - DR + M(PC + 1) + 1	х	X	Х	X	=	-	x	х	-	Х	а	Y
S880 <i>DR. AR</i>	Subtract byte	Reg dir.	334	6	OR ← OR + M(AR) + 1	X	X1	X	X	-	-	X	X	_	$\mathbb{Z} X$	0	.∵Y
SBBD DR, = label	Subtract byte	Lit. dir.	324	6	DR ← DR + M(ABR) + 1	X	X	Х	Х	_	-	Х	X	-	X	.0	Y
SBM DR, AR	Subtract multi-byte	Reg. imm.	305	4 + B	DR - OR + AR + 1	х	X	Х	X	=	-	Х	X	-	Х	0	Y
SBM DR; = literal	Subtract multi-byte	Lit. imm.	315	4 + B	DR - DR + M(PC + 1) + 1	X	Х	Х	Х	-	-	Х	Х	_	X	0	Y
SEMD DR. AR	Subtract multi-byte	Reg. dir.	335	5 + B	DR ← DR + M(AR) + 1	Х	Х	Х	Х	-		X	х		Х	0	Y
SBMD DA; = L'ODEL	Subtract multi-byte	Lit. dir.	325	5 + B	DR - DR + M(ADR) + 1	X	X	Х	X	-	-	Х	Х	-	) X	0	- 53
STB DR, AR	Store byte	Reg. imm.	242	5	DR - AR	X	×	×	×	85	-	Đ	D	-	D	0	
STB DR = literal	Store byte	Lit. imm.	252	5	DR → M(PC + 1)	X	X	х	×	ि	7	a	D	=	0	0	
STBD DR, AR	Store byte	Reg. dir.	246	6	DR - M(AR)	X	X	X	Х	-	-	0	0	-	0	0	
STBD DR. = lebel	Store byte	Lit. dir.	262	6	DR → M(ADR)	X	X	X	Х	-	-	0	0	-	0	0	
STBD DR, XAR, label	Store byte	Index dir.	266	8	DR -+ M(ADR + AB)	Х	X	Х	X	22		0	D	_	0	0	154
STBI DR. AR	Store byte	Reg. indir.	256	8	DR - M(M(AR))	X	X	X	X	_	_	0	Đ	-	0	0	

Section 8: Reference Material

20000000000		r <del>a can</del> nonneonn								Sta	tus						Binary
Instruction Format	Description	Addressing Mode	Opcode	Clock Cycles	Operation			RDZ			E	)CM	= 0	Ľ	CM	= 1	BCD Option
						LSB	MSB	LDZ	Z	DCM	E	CY	OVF	£	CY	OVF	574188
STBI DR. = label	Store byte	Lit. indir.	272	8	DR → M(M(ADR))	Х	X	X	X	***	-	0	0	-	0	0	
STBI DR, XAR, label	Store byte	Index indir	276	10	DR → M(M(ADR + AR))	х	X	х	Х	-	_	0	ū	_	0	0	
STM DR. AR	Store multi- byte	Reg. imm.	243	4 ± B	DR AR	∃X	Х	х	X	+	-	0	0	-	0	0	
STM DR. = literal	Store multi- byte	Lit. imm.	253	4 ± 8	DR → M(PC + 1)	×	×	х	X	-	_	0	0	-	D	0	
STMO DR. AR	Store multi- byte complement byte	Reg. dir.	247	5 + B	DR M(AR)	x	Х	х	X	=.	-	0	0	-	0	0	
NCM DR	Nine's (or one's) complement multi-byte	Reg. imm.	217	4 + B	DR DR	x	X	X	Х	-	-	X	x	==	X	Ü	
GRB DR AR	Or byte inclusive	Reg imm	224	5	DR - DR Y AR	X	χ	×	X	-	-	0	0	-	0	0	
GRM DR. AR	Or multi-byte inclusive	Reg. (mm.	225	4 + B	DR + DR Y AR	х	X	x	X	<u> </u>		0	D	Ξ	0	0	
PAD	Pop ARP, DRP and status from stack		237	38	Status + M(SP)	х	х	X	X	X	-	×	Х		X	Χ	
P080 DR. +AR	Pop byte with post- increment	Stk. dir.	340	6	OR M(AR). AR AR + 1	X	Х	Х	X	-	-	0	D	-	0	0	
POBO <i>DRAR</i>	Pap byte with pre- decrement	Stk. dir.	342	6	DR - M(AR). AR - AR - 1	Х	х	X	х	-		О	0	=	0	0	
POBI <i>DR</i> , +AR	Pop byte with post- increment	Stk. indir.	350	8	DR M(M(AR)), AR AR + 2	х	х	х	×		777	Ü	0		0	0	
POBI <i>DR, —AR</i>	Pop byte with pre- decrement	Stk. indit.	352	В	DR ← M(M(A8)), AR ← AR — 2	х	х	Х	X	-		0	0		0	0	
POMO DR. +AR	Pop multi-byte with post- increment	Stk. indir.	341	5+B	DR ← M(AR), AR ← AR + M	x	Х	×	X	_	_	0	0	-	0	0	
STMD DR. =	Store multi- byte	Lit. dir.	263	5 + B	DR - M(ADR)	X	X	х	X		7	U	0	=	0	0	
STMD DR, XAR, label	Store multi- byte	Index dir.	267	7 + B	DR → M(ABR + AR)	х	X	х	X	-	-	0	0	-	0	0	
STMI <i>DR, AR</i>	Store multi- byte	Reg. indir.	257	7 + B	DR → M(M(AR))	×	x	Х	X	-	===	D	0	T	Ü	0	
STMLDR; = Jabel	Store multi- byte	Lit. indir.	273	7 + 8	DR - M(M(ADR))	Х	χ	х	X	-	-	0	0	-	0	0	
STMI DR. XAR. (abal	Store multi- byte	Index indir.	277	9 + 8	DR - M(M(ADR + ARI)	X	X	Х	X	-		0	0	5	0	a	
TCB DA	Ten's (or two's) complement byte	Reg. imm.	214	5	DR ← DR + 1	ж	Х	Х	X	_	-	0	:0	-	0	0	SYS.
TCM DR	Ten's (or two's) complement multi-byte	Reg. imm	215	4 + B	DR ← DR + 1	x	х	х	X	_	=	0	0	-	0	0	¥
TSB DR	Test byte	Reg. imm.	220	5	Test DR	Х	X	Х	X	_	_	X	Х	_	χ	0	Y
TSM DR	Test multi- byte	Reg. imm.	221	4 + B	Test DR	X	х	Х	X	-	-	X	х	1	X	0	Y
XRB OR, AR	Or byte exclusive	Reg. imm.	225	5	DR - DR + AR	х	х	Х	X	-		0	0	-	0	0	
XRM DR, AR	Or multi-byte exclusive	Reg. imm,	227	4 + 8	DR - DR + AR	X	x	X	X	_	23	0	0	-	0	0	

# 8.10 Assembler Instruction Coding '

7	б	5	4	3	2		1	0
0	DRP/ ARP	≠000001 =000001	Load w Load w	ith literal ith RØ				
1	0	0	0	Ö	Logical/ Extended	R	ight/Left	M/B
1	0	0	0	1	0		ecrement/ Increment	М/В
1	0	0	0	ñ	ä	Nine' Ten'	s Complement/ s Complement	M/B
1	0	Ð	î	0	Ü	C	lear/Test	M/B
1	0	Ü	1	0	1		XOR/OR	M/B
1	C	0	).	3	000 001 010 011 100 101 110		BIN BCD SAD DCE ICE CLE RIN PAD	
1	0	1	000 001 010 011 100 101 110	REG REG LIT REG LIT INX LIT INX	IMM DIR IMM IMD DIR DIR IND	Š	tore/Load	M/B
1	4	0	00 01 10	REG IMM LIT IMM LIT DIR	01 10		CMP ADD SUB	M/B
	do.		11	REG DIR	0 1 1000		AND	1
1	1	۵	00	INX	11		JSB	0
1	i	1	0	IND, DIR	/ PU:	SH/ OP	-ADR/ +ADR	M/B
1	1	1	ä		000 001 010 011 100 101 110	P R		JNO/JMP JEY/JOD JPS/JNG JZP/JNZ JEZ/JEN JCY/JNC JLN/JLZ JRN/JRZ

X/Y = 1/0

## 8.11 Keycode Table

	KEYO	CODE		KEYC	ODE
)EC	OCT	KEY	DEC	OCT	KE\
0	0	ctrl@	48	60	0
1	1	ctrl A	49	61	1
2	2	ctrl B	50	62	
3	3	ctrl C	51	63	2
4	4	ctrl D	52	64	4
5	5	ctrl E	53	65	5
6	6	ctrl F	54	66	6
7	7	ctrl G	55	67	7
8	10	ctrl H	56	70	8
9	11	ctrl I	57	71	9
10	12	ctrl J	58	72	20
11	13	ctrl K	59	73	194
12	14	ctrl L	60	74	<
13	15	ctrl M	61	75	
14	16	ctrl N	62	76	= >
15	17	ctrl O	63	77	?
16	20	ctrl P	64	100	@
17	21	ctrl 0	65	101	A
1 B	22	ctrl R	66	102	В
19	23	ctrl S	67	103	C
20	24	ctrl T	68	104	D
21	25	ctrl U	69	105	E
22	26	etrl V	70	106	F
23	27	ctrl W	71	107	G
24	30	ctrl X	72	110	н
25	31	ctrl Y	73	111	4
26	32	ctrl Z	74	112	J
27	33	ctrl [	75	113	K
28	34	ctrl \	76	114	1
29	35	ctrl ]	77	115	M
30	36	ctrl ^	78	116	N
31	37	ctrl	79	117	0
32	40	SPACE	80	120	P
33	41	1	81	121	ū
14	42	40	82	122	R
35	41	H	83	123	S
6	44	\$	84	124	T
37	45	%	85	125	u
38	46	&	86	126	v
19	47		87	127	W
0	50	0	88	130	X
11	51	1	89	131	Ÿ
12	52	\$2 *2	90	132	z
3	53	4	91	133	Ĩ
44	54	+ >	92	134	1
45	55		93	135	
46	56		94	136	-80
17	57	1	95	137	

### 8.12 Programming Hints

If execution of certain advanced programming ROM statements is attempted in assembler mode, unpredictable results can occur. These statements are:

- X REF L
- X REF V
- REPLACE VAR

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